



Al Albayt University

Prince Hussein Bin Abdullah College for Information Technology

Department of Computer Science

**VELOCITY AND CONGESTION-AWARE ROUTING PROTOCOL
FOR MOBILE AD-HOC NETWORKS**

**بروتوكول لتحديد المسارات في الشبكات الخاصة المتحركة اعتماداً على السرعة
والازدحام**

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تفويض

أنا أحمد محمد عقلة المدان، أفوض جامعة آل البيت بتزويد نسخ من رسالتي للمكتبات أو المؤسسات أو الهيئات أو الأشخاص عند طلبهم حسب التعليمات النافذة في الجامعة.

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أعلن بأنني قد التزمت بقوانين جامعة آل البيت وأنظمتها وتعليماتها وقراراتها السارية المفعول المتعلقة بإعداد رسائل الماجستير والدكتوراه عندما قمت شخصياً بإعداد رسالتي بعنوان :

Velocity and Congestion-aware Routing Protocol For Mobile Ad-hoc Networks

وذلك بما ينسجم مع الأمانة العلمية المتعارف عليها في كتابة الرسائل والأطاريح العلمية. كما أنني أعلن بأن رسالتي هذه غير منقولة أو مستله من رسائل أو أطاريح أو كتب أو أبحاث أو أي منشورات علمية تم نشرها أو تخزينها في أية وسيلة إعلامية، وتأسيساً على ما تقدم فأنتني أتحمّل المسؤولية بأنواعها كافة فيما لو تبين غير ذلك بما فيه حق مجلس العمداء في جامعة آل البيت بإلغاء قرار منحي الدرجة العلمية التي حصلت عليها وسحب شهادة التخرج مني بعد صدورها دون أن يكون لي حق في التظلم أو الاعتراض أو الطعن بأي صورة كانت في القرار الصادر عن مجلس العمداء بهذا الصدد.

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COMMITTEE DECISION

This Thesis (VELOCITY AND CONGESTION-AWARE ROUTING PROTOCOL FOR MOBILE AD-HOC NETWORKS) was Successfully Defended and Approved on 3/1/2017

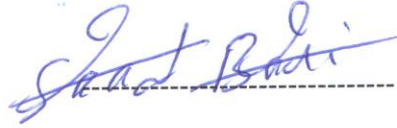
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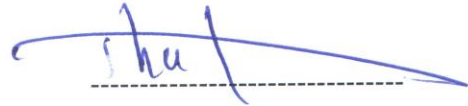
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DEDICATION

I dedicate this thesis to my great father Mohammed Al-Maddan, my lovely mother Sokout, you have successfully made me the person I am becoming.

To my lovely wife Duha Malw-Elain, and my beautiful daughter Mira, you are always the most important part of my success.

To all my family, for their support.

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LIST OF ABBREVIATIONS

Abbreviation	Meaning
AODV	Ad hoc On-Demand Distance Vector
ABR	Associativity-Based Routing
CRP	Congestion adaptive Routing Protocol
CARM	Congestion Aware Routing protocol for Mobile ad hoc networks
CF	Congestion Factor
CMF	Congestion-Movement Factor
CBR	Constant Bit Rate
DSDV	Destination-Sequence Distance-Vector
DLAR	Dynamic Load-Aware Routing
DSR	Dynamic Source Routing
GPS	Global Positioning System
LWR	Load aWare Routing
LAR	Location-Aided Routing
MANET	Mobile Ad hoc NETwork.
OTCL	Object oriented extension Tool Command Language
PT	Pause Time
RERR	Route Error
RREQ	Route Request
RREP	Route Reply
SSA	Signal Stability-based Adaptive
TORA	Temporally-Ordered Routing Algorithm
VCAR	Velocity and Congestion-Aware Routing
WCD	Weighted Channel Delay
WRP	Wireless Routing Protocol
ZRP	Zone Routing Protocol

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ABSTRACT

Mobile Ad-hoc NETWORKS (MANETs) are wireless infrastructure-less networks built from mobile nodes. Nodes in MANETs need routes between them when they attempt to communicate. They use a routing algorithm to find and maintain such routes. Most routing protocols proposed for MANETs are designed to find a valid route from source to destination without considering network traffic load and mobility of nodes. There are some algorithms that do consider only congestion when selecting the path from source to destination. The problem with such algorithms is selecting unstable routes, and increases the number of link failures in the network, and consequently increases the packets loss. While there are other algorithms based only on nodes velocity. This type of algorithms selects congested routes for transmitting, and consequently increasing the end-to-end delay of the packets.

In this thesis, we propose a Velocity and Congestion-Aware Routing (VCAR) protocol that aims to select the most stable and least congested intermediate nodes. The VCAR protocol uses two metrics when selecting the best route between source and destination: the congestion level and velocity of intermediate nodes.

To estimate the congestion level at intermediate nodes, we use two metrics: the interface queue length: Q_{len} , and the number of routes that the node participates in: R . These metrics are measured by the nodes themselves.

We use two approaches to select the route between the source and the destination. The first approach selects the route with the least accumulative value for all intermediate nodes in the path. While the second approach selects the route with the least maximum value among all intermediate nodes in the path.

We implement four versions of VCAR. The first version selects the best route using the least accumulative value of Qlen and velocity for all intermediate nodes. The second version selects the best route using the least maximum value of Qlen and velocity among all intermediate nodes. The third version selects the best route using the least accumulative value of R and velocity for all intermediate nodes, and the fourth version selects the best route using the least maximum value of R and velocity among all intermediate nodes.

The NS2 simulator was used to implement the VCAR protocols, and extensive simulations were conducted to analyze the performance of VCAR variants against the Ad Hoc on-Demand Vector (AODV) protocol. The results show that VCAR variants always outperform AODV in terms of packet delivery ratio, average end-to-end delay, and energy consumption percentage, but the routing overhead was higher.

CHAPTER ONE: INTRODUCTION.

Wireless communication has been widely used during the last decades, because this type of communication can offer services to users anywhere and at any time. Wireless communication can either have infrastructure or be infrastructure-less. The first class needs an infrastructure, and nodes can only communicate through a central device. While the second one does not need any infrastructure or central administration and is commonly known as ad hoc [15, 23, 26, 28]. MANETs are common infrastructure-less wireless networks [15, 24, 25, 28].

A MANET is usually defined as a set of mobile nodes without any infrastructure or central administration. The nodes are free to move around in any direction. Each node in a MANET has a limited transmission range and can operate as a host and as a router [8, 13, 14, 15, 28]. The node can directly send a message to all other nodes within its transmission range, but for the nodes outside the transmission range, the message should be forwarded to the destination via other intermediate nodes in the MANET [8, 23]. To route a messages in the MANET, a routing algorithm is needed to allow the node to discover and maintain a route to the destination.

1-1MANET Characteristics

MANETs have several characteristics that are very important to be understood when studying the nature of this type of networks [15, 28]. The following summarizes these characteristics:

1. Each node in a MANET has a limited transmission range [4]. A route is needed to reach nodes outside this range. As shown in Figure 1-1, node B lies within the transmission range of node A, so a direct message can be sent from node A to node B,

while a route is needed between node A and node C because node C lies outside the transmission range of node A.

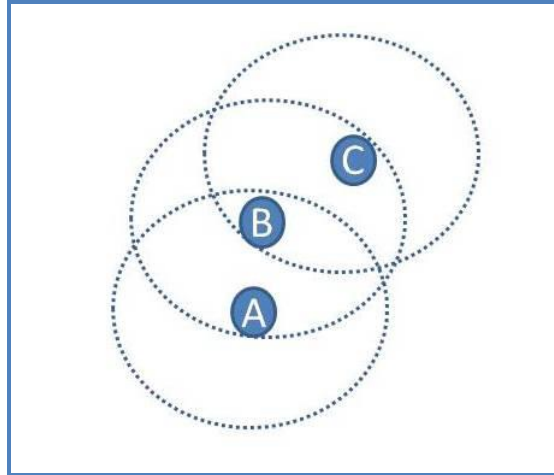


Figure 1-1: Node Transmission Range

2. Each node in a MANET is battery-powered and the power is quickly consumed because each node acts as a host and as a router [4].

3. The nodes in a MANET are free to move around in any direction and this makes the topology change frequently. This characteristic should be taken into account when designing the routing algorithm. Figure 1-2 shows the effect of mobility on the communication of nodes. In Figure 1-2(b) node B has moved so that it is now outside the transmission range of node A. When node A needs to communicate with node B, it must send the message to node C, which sends it to node B.

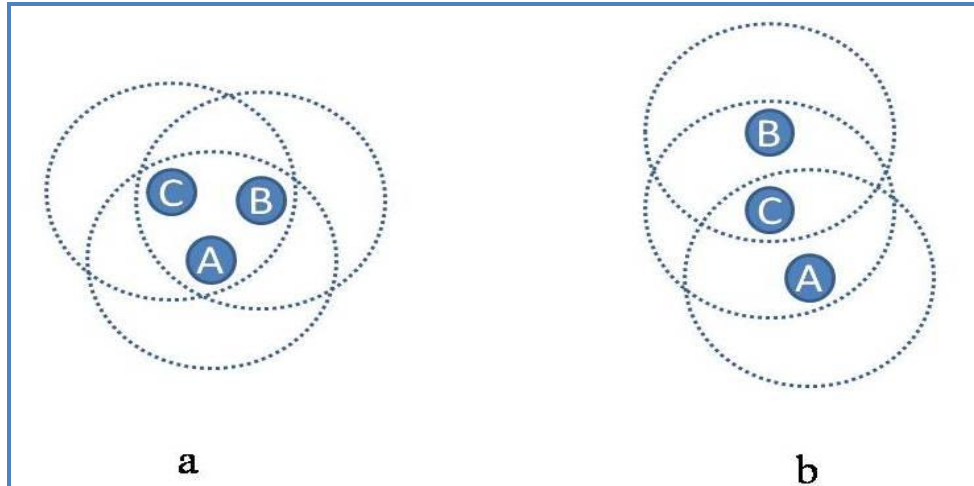


Figure 1-2: Node Mobility Causes Topology Change in MANET

1-2MANET Applications

MANETs can be used when building a fixed network is expensive or hard to do [25].

MANETs have been proposed for use in many industrial and commercial fields [24]:

1. MANETs have been used in wireless sensor network applications. These applications are used mainly in environmental fields, such as data tracking and remote sensing for weather forecasting.

2. MANETs have been proposed for some civilian applications, such as search and rescue operations, and disaster relief efforts in areas where there is no network infrastructure.

3. MANETs have been used in military applications because of their rapid deployment capability.

1-3 Routing in MANETs

When studying MANETs, one of the most important questions is how mobile nodes can communicate with each other without any infra-structure using wireless media. The answer is that each node needs to maintain a routing path to nodes with which it needs to communicate. Many routing algorithms were proposed for MANETs.

One of the most important issues that must be considered when designing any routing algorithm is selecting a route that can adapt well to topology changes [13, 29].

Routing influences the overall performance of the MANET. A more efficient routing algorithm leads to a better MANET performance. Routing in MANETs should take into account the mobility of the nodes because nodes are free to move around. The routing information should be updated dynamically. There are mainly three types of routing protocols in MANETs: proactive, reactive, and hybrid [5, 8, 10, 15, 23, 28, 30].

1-3-1 Proactive Routing Protocols.

In proactive routing, each node stores routing information in a table, so this type is also called table-driven routing [5, 23, 30]. When a node needs a route to another node, it finds the route by searching in its routing table. The routing table information is propagated periodically through the network, so that each node has an updated routing table. This type of routing has low latency because the routes are always available in the table, but it suffers high overhead because of periodic routing table updates. Examples of such type of routing include Destination-Sequence Distance-Vector Routing (DSDV) [20], and Wireless Routing Protocol (WRP) [17].

1-3-2 Reactive Routing Protocols.

In reactive routing, the route is determined only when needed, so that it is also called source-initiated on-demand routing [5, 23, 30]. When a node needs to send a message to another node for which it does not have a known route, it first finds a route to that node by flooding the network with a Route Request (RREQ) packet, and then it uses the discovered route to send the message. This type of routing has high latency. However, it has low overhead because routes are determined only when needed. Examples of routing protocols of this type include Ad hoc On-Demand Distance Vector (AODV) protocol [21], and Dynamic Source Routing (DSR) protocol [9].

1-3-3 Hybrid Routing Protocols.

This type of routing combines the advantages of the above two types. In hybrid routing, the routes between nodes are initially discovered proactively, then they are maintained reactively as a result of link failures or topology changes [5, 23, 30]. Examples of this type of routing protocols include Zone Routing Protocol (ZRP) [27], and Temporally-Ordered Routing Algorithm (TORA) [19].

1-3-4 Ad-hoc On-demand Distance Vector (AODV) Routing Protocol

AODV is a reactive routing protocol proposed in [21]. AODV uses a route discovery process to find the shortest path between the source and the destination. AODV consists of the following phases.

- **Route Discovery Phase.**

When the source does not know a valid route to the destination it broadcasts a route request (RREQ) packets to all of its neighbors. The RREQ packet has a unique sequence number to

detect duplicate packets. It also includes some other information such as: destination identifier, source identifier, and time to live. Each intermediate node that has been received the RREQ uses the sequence number field to determine if the received RREQ has been previously received or not. Then it will check if it has a valid route to the destination, if so it will send a Route Reply (RREP) packet to the source, otherwise it will rebroadcasts the RREQ packet to all of its neighbors, and the time to live field will decremented by one to prevent packet looping. The RREQ packet continue traveling until it reaches to the destination. When the destination receives the RREQ, it replies with a Route Reply Route (RREP) packet to the intermediate node from which it receives the RREQ. The RREP travels back until it received by the source. [21].

- **Route Maintenance Phase.**

The route maintenance phase is responsible for detecting any link failure. The node detects any link failure by listening to hello messages from its neighbors. If any link failure is detected, a node broadcasts a Route Error (RERR) packet to notify the source that it detects a link break, so that the source will initiate a new RREQ packet to search for a valid route to the destination [25].

1-4 the Problem and Motivation.

Most of reactive protocols proposed for MANETs have been designed to choose the shortest path from source to destination [14]. The shortest path is the path with the smallest number of hops between the source and destination. Using this single metric can lead to a situation in which the nodes that participate in the shortest path are congested while other nodes are idle or lightly-loaded [14, 16]. Figure 1-3 shows such situation in which node D lies in the shortest path in most of the routes because of its location in the middle section of

the MANET, and this increases the probability of congestion at node D. For example, the shortest path from A to F is A->D->F, while other routes are available such as A->C->E->F and A->B->G->F. To avoid such problems, other metrics should be used when selecting the best path.

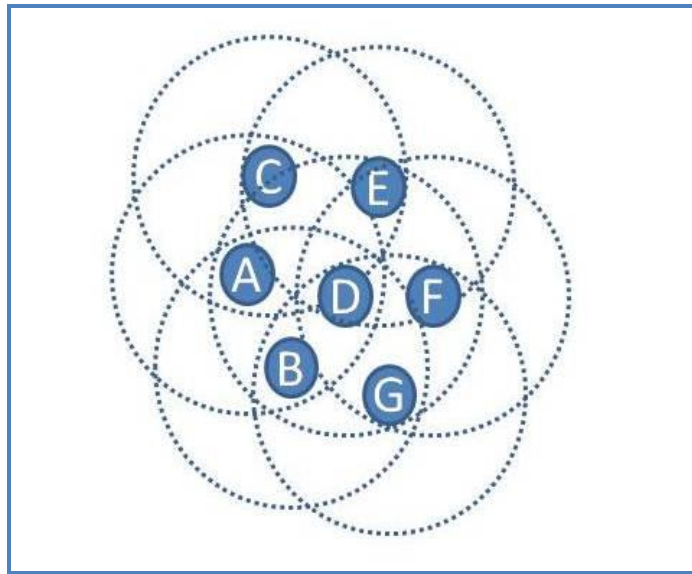


Figure 1-3: Shortest Path Leads to Congested Nodes

Another problem with such protocols is that nodes in MANET may often move around. The movement of nodes results in dynamic changes in routes, as shown previously in Figure 1-2, and this requires finding new routes. Continuous unpredictable changes in MANET topology increase the overhead in the route maintenance mechanism [13, 29].

To increase the performance of the routing protocol, the route discovery mechanism should take into account both the congestion state and movement of the nodes. The routes in a MANET should be as stable as possible. The metric used in route selection should combine both the congestion level and the movement of nodes in the path selected. Any node participating in a route should have low congestion and low speed.

1-5 The Objectives of The Study

The purpose of this study is to propose a new routing protocol for MANETs called Velocity and Congestion-Aware Routing (VCAR) protocol. VCAR selects the routes between nodes according to both the congestion and velocity of intermediate nodes. VCAR aims to achieve the following goals:

- Reduce the congestion in the network.
- Increase the packet delivery ratio in the network.
- Decrease the average end-to-end delay of packets.
- Decrease the energy consumption of the nodes.

CHAPTER TWO: RELATED WORKS

Routes determined by reactive routing protocols can be congested. Many algorithms have been proposed to solve this congestion problem. Most of these algorithms use some metric to measure node congestion, and select the least congested node or route [8, 16, 26].

Routes can be unstable because of node movement in MANETs. Therefore, some routing algorithms use the mobility metric to enhance network performance. These algorithms try to choose route nodes with low mobility so as to discover relatively stable paths. A more stable path decreases the path failure rate and increases the throughput of the network [5].

2-1 Associativity-Based Routing Protocol

The Associativity-Based Routing (ABR) [29] protocol uses node stability as the main metric in selecting the best path. ABR is a reactive protocol that searches for a route only when a source needs to find a route to a destination. ABR uses an associativity-based scheme in which a route is constructed from nodes that have an associativity state with their neighbors that guarantees stability. Thus, ABR selects routes that are likely to be long-lived. ABR measures node stability by the node's association with its neighbors. ABR finds all the possible routes from source to destination, and then selects the best path according to the selection criteria. The disadvantage of ABR is that it does not consider node congestion in the route selecting process; this may lead to using a route with congested nodes.

2-2 Signal Stability-Based Adaptive Routing Protocol.

The Signal Stability-Based Adaptive (SSA) [3] routing protocol is another on-demand protocol that proposed to perform route discovery by selecting the longest-lived path between source and destination. SSA selects the best route by including signal strength and node location stability. This protocol ranks the channels as strong and weak according to their average signal strength between the two ends of the channel. The node location stability is used by SSA so as to choose the longest-lived route. By considering these two criteria, SSA always chooses strong channels that have existed for a period of time that is greater than some threshold value identified by the protocol. In SSA, the source broadcasts the RREQ packet to all of its neighbors. When any intermediate node receives the request packet, it rebroadcasts it only if it was received over a strong channel. The destination chooses the route of the first arriving request because it is probably shorter and less congested, and sends a reply message that contains the best route. The main disadvantage of SSA is that it works fine only when there are a significant number of strong routes between source and destination.

2-3 Location-Aided Routing Protocol.

The Location-Aided Routing (LAR) [13] protocol was proposed to improve the performance of routing discovery by utilizing the location information of the nodes that are obtained using a Global Positioning System (GPS). LAR uses location information to limit the search of a new route to a small area in the MANET. This small area is called request zone. LAR aims to reduce the overhead by decreasing the number of control packets in the discovery process. Once the request zone is identified, the source searches only within this zone, without flooding all the MANET with RREQ packets. The request zone is

rectangular in shape and is identified by the source. When the source wants to discover a route to the destination, it must identify the request zone of the destination by using some information that includes the previous location and the average speed of the destination. After identifying the request zone, the source broadcasts the RREQ packet only to the nodes that are located within the request zone boundaries. LAR performance depends on both the availability and accuracy of GPS used, and the need to use GPS is the major disadvantage of LAR.

2-4 Dynamic Load-Aware Routing protocol

Lee and Gerla proposed the Dynamic Load-Aware Routing (DLAR) [14] protocol in 2001. DLAR is a reactive protocol that uses the routing loads of nodes as the main metric to select the best path. In this protocol, the highly loaded nodes are prevented from participating in the routes. The load of a node is measured by the number of packets buffered in its interface queue. This protocol selects always the more stable path, and reduces end-to-end delay. DLAR also controls the congestion states of the active routes and reconstructs them when any node in the path reaches its maximum queue capacity. In DLAR, the source broadcasts a request packet to find a route to the destination, and each intermediate node appends its load and broadcasts the packet again. The receiver can select the best path according to the loads of the intermediate nodes. It selects the path with the lowest aggregate load. The main disadvantage of DLAR is that it has a large overhead caused by the request flooding process in which each intermediate node appends extra information before rebroadcasting the request.

2-5 Load aWare Routing Protocol.

In 2003, a new protocol was proposed by Yi and Gerla. It is called the Load aWare Routing (LWR) [31] protocol. LWR was built using the same idea as DLAR and was proposed to solve the problem of large overhead. LWR is different from DLAR in that each intermediate node has a routing selection scheme with a congestion control mechanism at the same time. In LWR, intermediate nodes drop any request packet when they are overloaded. Another difference is that LWR concentrates on the route discovery phase to prevent any unnecessary broadcasting of requests, while DLAR does this in the route reply phase. LWR decreases the overhead and improves the routing performance.

2-6 Congestion Adaptive Routing Protocol.

The Congestion adaptive Routing Protocol (CRP) [30] is an adaptive routing protocol that tries to avoid congestion rather than dealing with it reactively. The main idea in CRP is that each node in any route should notify its previous node when it is about to be congested. To prevent congestion, CRP introduces a secondary route called the bypass route. CRP uses the bypass route to bypass the congestion area to the first non-congested node on the primary route. The traffic is split over these two routes. CRP has a significant overhead when there is a bypass route for each primary one. Another disadvantage of CRP is that there is a small loss rate because of dividing the traffic into two routes.

2-7 Congestion Aware Routing Protocol.

The Congestion Aware Routing protocol for Mobile ad hoc networks (CARM) [2] is another congestion-aware protocol that introduces a new parameter to measure the congestion level. This parameter is called the Weighted Channel Delay (WCD). WCD is

calculated using two weighted parameters: the total time spent at the MAC layer and the number of buffered packets in the node, as shown in Equation 2-1.

$$WCD = \alpha Q + bT \dots\dots\dots \text{Equation 2-1}$$

Where Q is the number of buffered packets, T is the total time spent at the MAC

layer, and α and b are constants between 0 and 1.

In CARM, the source node broadcasts a RREQ packet. Each intermediate node calculates its WCD, appends it in the request packet, and rebroadcasts the packet. When the destination receives the request, it responds with a reply message that contains the best route by choosing the smallest aggregate value of WCD of all intermediate nodes. CARM improves the adaptability to congestion by using the WCD parameter, which is used for selecting routes with high throughput and low congestion. On the other hand, CARM causes overhead at intermediate nodes.

2-8 Congestion-Aware Routing Protocols for Ad Hoc networks.

Another congestion-aware routing protocol was proposed by Seetan, Ibabneh, and Dala'ah in their paper [25] called Min_Total_CA protocol. In this protocol the congestion is measured using the number of routes that the node participates in. Min_Total_CA selects the route with the least total value of congestion among all available routes between the

source and the destination. Each route has two congestion value: the forward congestion value, and the backward congestion value. The first value is calculated through the route discovery process, and the second value is calculated through the route reply process. When the source does not have a valid route to the destination, it broadcasts a RREQ packet to all of its neighbors. Each intermediate node receives a RREQ packet, will reply with a RREP packet if it has a valid route to the destination. Otherwise, it adds its forward congestion value to the RREQ packet. The RREQ packet travels through the network until it reaches to the destination. The destination checks if it has a valid route to the source, if not it replies with a RREP packet. But if there is a valid route, then the destination will send RREP packet only if the received forward congestion value is smaller than that in the currently used route. The RREP travels back to the source, and each node in the reverse path adds its backward congestion value to the backward congestion value in the RREP. When the source receives a RREP packet from the destination, it checks if it's the first RREP, if so it starts using this route. Otherwise, it checks the freshness of the RREP by checking if the new reply has a destination sequence number that larger than that in the route currently in use, if so it starts using the new route. If the new and the current routes have the same freshness, the source checks if the backward congestion value of the received RREP is lower than that in the current route, if so it starts using the new route.

CHAPTER THREE: THE PROPOSED STUDY

3-1 Velocity and Congestion-Aware Routing Protocol

In this section, we present the design of VCAR protocol and illustrate the network model, congestion metrics, and node movement metrics.

3-1-1 Network model

We consider a MANET that contains mobile nodes with bidirectional wireless links connecting them. All nodes have the same transmission range, initial battery life, processing capacity, mobility pattern, transmission bandwidth, and responsibilities (each node can route packets). Each node has its own movement speed, and movement direction.

3-1-2 Congestion Metrics.

To address the congestion level at each node, we propose to use two metrics: Interface Queue Length (Qlen), and number of Routes (R) that the node participates in. These two metrics can be measured by the node itself so as to measure the congestion level.

- Queue Length (Qlen).

Each node has an interface that is used to communicate with other nodes. Packets are queued in the interface until they have been transmitted. At any given time, there are a number of waiting packets in the queue (Qlen). Lower Qlen value indicates lower congestion level, and vice versa.

- Number of Routes (R).

All nodes can act as routers, so any node could participate in a number of routes. To measure the congestion level at a node, it is important to count how many routes each

node participates in. A lower value of R means lower node congestion level, and vice versa.

VCAR uses either Qlen or R to measure the congestion level at each node. One of our contributions is to evaluate the use of these factors along with the node movement factor (which will be illustrated in the following subsection). For the rest of this thesis we will use the term Congestion Factor (CF) to indicate either Qlen or R.

3-1-3 Node Movement Metric.

In MANET routes, nodes that do not move frequently will provide better performance than other nodes because they reduce the overhead in construction and reconstruction of routes. The routes that are composed of slow nodes are more stable than other routes. Our proposed algorithm attempts to use slow nodes in constructing routes between nodes. We use node Velocity (V) as an indicator of node stability level. A lower value of V indicates higher node stability, and vice versa.

3-1-4 VCAR Design

VCAR reactively discovers routes from source to destination and selects the best route according to the values of CF and V. The main metric for selecting the best route is called the Congestion-Movement Factor (CMF). CMF is calculated as per Equation 3-1.

$$CMF = \alpha * CF + (1 - \alpha) * \frac{V}{V_{max}} \dots\dots\dots \text{Equation 3-1.}$$

Where, V_{max} is the maximum velocity of a node.

The value of (α) is between 0 and 1. This value should be tuned to achieve good performance.

There are two approaches to select the best route between the source and destination. The first approach selects the route with the least accumulative CMF value for all nodes in the path. While the second approach selects the route with the least maximum CMF value.

To differentiate between these two approaches, suppose we have the MANET shown in Figure 3-1. Each node is associated with its CMF value.

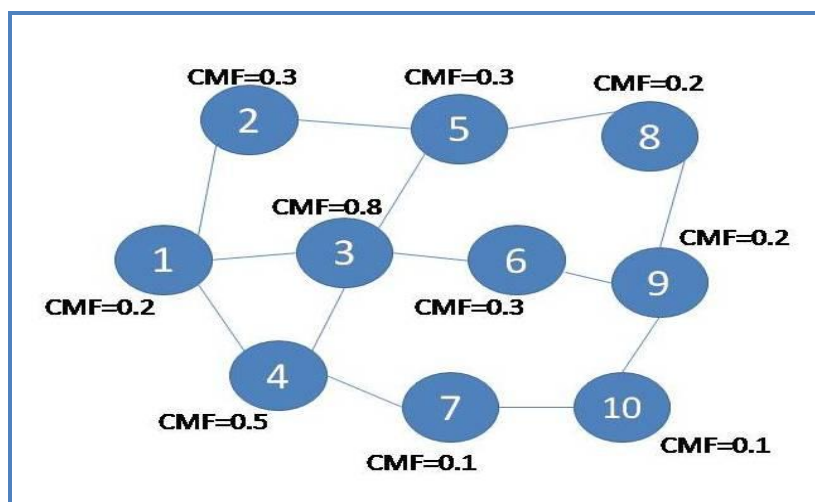


Figure 3-1: Each Node in a MANET Has its Own CMF Value

Suppose that node 1 needs to send data to node 9 and it has not a valid route to that node. Then it should discover a new route. There are many possible routes from node 1 to node 9. Each one of these possible routes has an accumulative CMF value for all intermediate nodes, and a maximum CMF among all the intermediate nodes, as shown in Table 3-1.

Table 3-1: Accumulative and Maximum CMF Values

Route	Accumulative CMF	Maximum CMF
1,2,5,8,9	0.8	0.3

1,2,5,3,6,9	1.7	0.8
1,2,5,3,4,7,10,9	2.1	0.8
1,3,5,8,9	1.3	0.8
1,3,6,9	1.1	0.8
1,3,4,7,10,9	1.5	0.8
1,4,3,5,8,9	1.8	0.8
1,4,7,10,9	0.7	0.5
1,4,3,6,9	1.6	0.8

If we use the least accumulative CMF approach, then the best route will be 1,4,7,10,9 because it has the least accumulative CMF value which is (0.7). On the other hand, If we use the least maximum CMF approach, then the best route will be 1,2,5,8,9 because it has the least maximum CMF value, which is (0.3).

VCAR consists of three phases: the route discovery phase, the route reply phase, and the route maintenance phase. In the first phase, the sender broadcasts a RREQ packet to discover a route to the designated destination. In the second phase, a Route Reply (RREP) packet is sent to the source to identify the selected route to the destination. Finally, the third one is responsible for sending a Route Error (RERR) packet when a route becomes invalid for any reason. These phases will be illustrated in details later in this section, but before that we will present some important information about routing control packets in the following sub-section.

3-1-4-1 VCAR Control Packets

There are three types of control packets: RREQ, RREP, and RERR. In this subsection, we will present the main fields of each one.

The RREQ packet contains the following main fields:

- Source Identification (SID): the address of the source node from which the route needs to be established.
- Destination Identification (DID): the address of the destination node, to which a packet or more are to be sent.
- Sequence Number (SEQ): a number that uniquely identifies each packet sent from a source. This number is used to detect duplicate RREQ packets.
- CMF: each node that receives a RREQ calculates its CMF as in Equation 3-1, and then updates the CMF field. The CMF field contains either the cumulative CMF values for all the nodes visited by the RREQ or the maximum CMF value among the CMF values of all the nodes visited by the RREQ.
- Number of Hops (HOP): number of hops traversed by the RREQ.
- Time To Live (TTL): represents the maximum number of hops that the RREQ can traverse before its lifetime expires. This value is used to prevent looping of packets. It is decremented by one at each node receiving the RREQ.

The RREP packet contains the following main fields:

- Source Identification (SID): the address of the destination of the RREQ.
- Destination Identification (DID): the address of the source node that initiated the RREQ.
- Sequence Number (SEQ): is a number that is uniquely assigned to a packet. This number is used to detect duplicate RREP packets.
- CMF: contains either the cumulative or the maximum CMF, as explained earlier.
- Number of Hops (HOP): number of hops traversed by the RREQ.

The RERR packet is the same as in AODV protocol [21].

3-1-4-2 Route Discovery Phase.

The route discovery phase is responsible for finding a route between two nodes: source and destination. The following steps describe this phase.

- 1.** If the source node does not have a valid route to the destination, it prepares a RREQ packet and broadcasts it to all neighbor nodes.
- 2.** If an intermediate node receives a RREQ packet, then it does one of the following:
 - 2.1** If TTL in RREQ is 0 or SEQ in RREQ had been processed previously, the node discards the RREQ.
 - 2.2** If the DID in RREQ is equal to the intermediate node's ID (i.e., the intermediate node is the destination node), the node replies with a RREP packet as will be explained in the route reply phase.

2.3 If the intermediate node is not the destination node, it does one of the following:

2.3.1 if the node has a valid route to the destination, the node replies with a RREP packet as will be explained in the route reply phase.

2.3.2 if the node does not have a valid route to the destination, the node takes the following actions:

2.3.2.A It calculates its CMF as expressed in Equation 3-1.

2.3.2.B It updates the value of CMF in the RREQ.

2.3.2.C The node rebroadcasts RREQ to all of its neighbors.

3-1-4-3 Route Reply Phase.

This phase is entered when an intermediate node has a valid route to the destination or the RREQ has arrived to the destination. The following steps describe this phase.

1. The responding node does the following:

1.1 It calculates CMF as expressed in equation 3-1.

1.2 It calculates the CMF value of the path, either accumulative or the maximum.

1.3 If it's the first RREP sent to the source, then go to step 1.5.

1.4 Otherwise, the node compares the calculated CMF with the last sent CMF value. If the current one is larger than the last one, then discard the packet. Else, go to step 1.5.

1.5 Node prepares a RREP packet that contains the CMF value.

1.6 The node sends its RREP to the node from which it received the RREQ.

2. The RREP message travels back to the source node that initiated the RREQ via all the nodes that previously rebroadcasted the RREQ packet.

3-1-4-4 Route Maintenance Phase.

The Route maintenance phase is responsible for detecting any link failure. Usually, link failure happens due to node movement. This phase is done as in AODV [21].

3-1-4-5 VCAR Versions

As we mentioned in section 3-3-2, we have two congestion metrics: queue length, and number of routes, and there are two approaches to select the best path: the least accumulative CMF value, and the least maximum CMF value. So that VCAR has four versions. Each version uses a congestion metric along with a measurement approach as shown in Table 3-2. We will use the names in the table to distinguish each one of the four protocols.

Table 3-2: VCAR Versions

Version	Congestion metric	Best Path Selection Approach	Protocol Name
1	queue length	least accumulative CMF	VCAR_Q_ACC
2	queue length	least maximum CMF	VCAR_Q_MAX
3	number of routes	least accumulative CMF	VCAR_R_ACC
4	number of routes	least maximum CMF	VCAR_R_MAX

CHAPTER FOUR: THE SIMULATION

Network simulators are used mainly to compare the performance of different routing protocols used in MANETs. NS-2, OPNET, QualNet, and GloMoSim are the most popular simulators used in this field. We used the NS-2 simulator to evaluate the performance of VCAR protocols against AODV. NS-2 is an accurate, event driven, and open source simulator. It contains the implantation of many existing routing protocols such as DSR and AODV [6].

4-1 NS-2 simulator Architecture

NS-2 is an object-oriented simulator, written in the C++ programming language, with an Object oriented extension Tool Command Language (OTCL) interpreter [18]. C++ is used mainly in the implementation of network protocols. While OTCL is used to write the simulation scenarios [6, 18].

The results from any simulation scenario are stored in a text files called trace files. These files contains all the events that take place in the network during the simulation. Another tool, called GAWK, was used to extract the needed information from these trace files.

4-2 Simulation environment

The NS-2 simulator uses a model for each layer [6]. We have used the models shown in Table 4-1.

Table 4-1: Models Used for Different Layers

Layer	Model
Application	CBR
Transport	UDP
Mac Layer	802.11

The simulation parameters include the physical channel specifications, mobility models, and network traffic. All simulations were performed using NS-2.35 version. Table 4-2 shows the simulation parameters used in the study, this environment is commonly used, and recommended in many studies such as [14, 25].

Table 4-2: The simulation parameters

Parameter	Value
Simulation time	300 seconds
Number of mobile nodes	50 nodes
Simulation area	1200 meters * 1200 meters
Node transmission range	250 meters
Maximum buffer queue	50 packets
Maximum speed of nodes	10 meters/second
Data flow type	CBR
Number of sending sources	5, 10, and 15 sources
Sending rate	1, 2, 4, and 6 packets/second
Routing protocol	AODV, VCAR_Q_ACC, VCAR_Q_MAX,

	VCAR_R_ACC, and VCAR_R_MAX
Weight factor used in VCAR protocols (α)	0.25, 0.50, and 0.75

The following explains each one of these parameters:

1. Simulation time: specifies the total simulation time in seconds. All the simulations lasted for 300 seconds.
2. Simulation area: specifies the dimensions of the simulation area. We used simulation area of 1200 meters * 1200 meters.
3. Number of nodes: specifies how many nodes were in the simulation area. We ran all the simulations with 50 nodes.
4. Mobility model: specifies the style of node mobility. NS-2 supports different mobility styles. The most widely used mobility style is the random-waypoint style. In this type of mobility, a node randomly chooses a destination in the simulation area, and moves toward this destination with a random velocity uniformly chosen between two values, generated as discussed below. When the node reaches its destination, it stays there for a period of time called Pause Time (PT). Then, it selects another destination and moves toward it. In this research we used random-waypoint style with PS values of: 0,100, 200, and 300 seconds. Node velocity is generated uniformly between 0 and 10 meters/second.
5. Traffic model: specifies the communication model used for transmitting packets. We used the Constant Bit Rate (CBR) model for sending packets from the source to the destination. Each packet is 512 bytes. In the CBR model a certain number of nodes

called communication sources send packets to other nodes at some rate. In this study, we use 5, 10, and 15 sources, with sending rates of 1, 2, 4, and 6 packets/second.

6. Transport layer protocol: specifies the transport protocol used in the simulation. We used UDP to ensure timely delivery of data packets with low network overhead.
7. Routing protocol: specifies the routing protocols used in the simulation. We ran the simulations using AODV, and VCAR protocols.

4-3 Experiments Design

For each experiment, we create two files: traffic connection file, and node movement file.

In the following two subsections, we illustrate how these two files were created.

4-3-1 Creating Traffic Connection File

The traffic connection file contains a number of CBR traffic connections generated using the command “ns cbrgen.tcl”. The general format of this command is as follows [6]:

```
ns cbrgen.tcl. [-type cbr] [-nn nodes] [-seed seed] [-mc connections] [-rate rate]
```

Where *type* specifies the traffic type, *nn* specifies the number of nodes in the simulation, *mc* specifies the number of traffic connections, and *rate* specifies the sending rate (packets/second).

As example, the following command creates 5 CBR traffic connections, each connection sends 1 packet/second.

```
ns cbrgen.tcl. -type cbr -nn 50 -seed 1 -mc 5 -rate 1
```

We created twelve traffic connection files to accommodate all the cases in the study. Each file was created using a combination of a number of communication sources and a sending rate as shown in Table 4-3.

Table 4-3: The Traffic Connection Files

File name	Number of communication sources	Sending rate (packets/second)
cbr-mc_5-rate_1	5	1
cbr-mc_5-rate_2	5	2
cbr-mc_5-rate_4	5	4
cbr-mc_5-rate_6	5	6
cbr-mc_10-rate_1	10	1
cbr-mc_10-rate_2	10	2
cbr-mc_10-rate_4	10	4
cbr-mc_10-rate_6	10	6
cbr-mc_15-rate_1	15	1
cbr-mc_15-rate_2	15	2
cbr-mc_15-rate_4	15	4
cbr-mc_15-rate_6	15	6

4-3-2 Creating nodes movement file

The nodes movement file contains the movement of nodes within the simulation area according to the random-waypoint mobility style. This file is generated using the command “*setdest*”. The general format of this command is as follows [6]:

setdest [-n num_of_nodes] [-p pausetime] [-m maxspeed] [-t simtime] [-x maxx] [-y maxy]

Where *n* specifies the number of nodes in the simulation, *p* specifies the pause time, *m* specifies the maximum speed, *t* specifies the simulation time, and *x* and *y* specify the dimension of the simulation area.

As example, the following command creates the node movement file for 50 nodes that move in an area of 500 meters* 500 meters, and follow the random waypoint style with a pause time=0 and a velocity generated randomly between 0 and 10 meters/second. The simulation lasts for 300 seconds.

setdest -n 50 -p 0.0 -m 10.0 -t 300 -x 500 -y 500

Because we have four pause times in the study: 0, 100, 200, and 300 seconds, then we create four corresponding files as shown in Table 4-4.

Table 4-4: Nodes Movement Files

Movement file	Pause time (seconds)
scen-0	0
scen-100	100
scen-200	200
scen-300	300

To cover all the cases, we create forty eight simulation scenarios. Each scenario exports one traffic file and one movement file.

4-4 Performance Comparison Metrics

Gawk tool was used to extract data from the trace files generated by the simulations. The performance analysis uses four metrics: packet delivery ratio, average end-to-end delay, routing overhead, and energy consumption percentage.

4-4-1 Packet Delivery Ratio

The packet delivery ratio is the ratio between the total number of received data packets to the total number of sent data packets [1, 7, 10, 11, 12]. As per equation 4-1.

.... Equation 4-1

4-4-2 Average End-to-End Delay

End-to-end delay is the delay that a packet suffers between the time it leaves the source application to the time it arrives at the destination application. The average end-to-end delay is the average value of such delays suffered by all the data packets received in the network [1, 7, 10, 11, 12]. It is computed as per Equation 4-2.

$$\text{Average End - to - End Delay} = \frac{\sum_{i=1}^P (RT_i - ST_i)}{P} \dots\dots\dots \text{Equation 4-2}$$

Where, P is the total number of data packets received in the network, RT_i is the time at which packet P_i was received, and ST_i is the time at which packet P_i was sent.

4-4-3 Routing overhead

Routing overhead is the ratio between the total number of control packets sent to the total number of data packets received [1, 7, 10, 11, 12]. It is computed as per Equation 4-3.

$$\text{Routing Overhead} = \frac{\text{total number of control packets sent}}{\text{total number of datapackets received}} \dots\dots\dots \text{Eq}$$

Equation 4-3

4-4-4 Energy Consumption Percentage

Energy consumption percentage is the ratio of the total energy consumed at all the nodes to the total initial energy of all the nodes. Energy consumption for each node is defined as the difference between its initial energy and the remaining energy at the end of the simulation.

It is computed as per Equation 4-4.

$$\text{Energy Consumption Percentage} = \frac{\sum_{i=1}^N (IE - RE_i)}{N * IE} * 100\% \dots\dots\dots \text{Equa}$$

Equation 4-4

Where N is the number of nodes in the network, IE is the initial energy of the node which is the same for all nodes, and RE_i is the remaining energy at node i .

CHAPTER FIVE: RESULTS AND DISCUSSION

In this chapter, we present and analyze the results obtained from the simulation scenarios and evaluate the performance of VCAR against other protocols. We ran the simulations using three values of (α): 0.25, 0.50, and 0.75. Empirically, we have found that $\alpha=0.50$ gives the best overall results. The following sections present and discuss the results obtained from the simulation scenarios when $\alpha=0.50$. To see the results when $\alpha=0.25$, and $\alpha=0.75$, please refer to appendix A, and appendix B, respectively. Each experiment was ran ten times, and the average value of these runs was considered.

We evaluate the performance of VCAR against a previous congestion-aware routing protocol called Min_Total_CA that was proposed by Seetan, Ababaneh, and Dalal'ah in [25]. We used the simulation environment that illustrated previously in section 4-2, and considered the average value for all the simulation scenarios. The results show that VCAR has achieved a better packet delivery ratio than Min_Total_CA. It has improved the packet delivery ratio by 6.50 percent. VCAR also outperforms Min_Total_CA by 6.00 percent in terms of average end-to-end delay. The results show that VCAR has a lower energy consumption than Min_Total_CA, with an improvement of 8.25 percent. The results also show that the routing overhead values for both VCAR and Min_Total_CA are fairly close to each other.

In the following sections we discuss the results of the simulations and evaluate VCAR performance against AODV performance.

5-1 Packet Delivery Ratio

Figures 5-1 ~ 5-12 show the packet delivery ratio for different transmission rate values and a different number of sources. The results show that all VCAR protocols outperform AODV in terms of packet delivery ratio for all mobility levels, all number of sources, and all transmission rates because they select less congested and more stable routes among all available routes, which reduces the number of dropped packets.

We noticed that all VCAR protocols outperform AODV more substantially when mobility is high and the load is heavy. On the other hand, they outperform AODV less when the mobility is low and the load is light. As example, when the number of sources is fifteen and each source sends six packets per second, VCAR_R_ACC outperforms AODV by 26.99 percent. On the other hand, when the number of sources is five and each source sends one packet per second, VCAR_R_ACC outperforms AODV by 13.94 percent.

The VCAR_R_ACC and VCAR_R_MAX protocols always have a better packet delivery ratio than VCAR_Q_ACC and VCAR_Q_MAX. The reason is that the last two versions (VCAR_Q_ACC and VCAR_Q_MAX) use buffer length to measure the congestion level at the intermediate nodes, and because the buffers at intermediate nodes are almost full (especially when the traffic is heavy) then these two protocols could not select the least congested route accurately.

The VCAR_R_ACC protocol always outperforms all other VCAR protocols in all the cases we have studied. On average, VCAR_R_ACC outperforms AODV by 15.63 percent for low mobility networks, and by 21.63 percent for high mobility networks.

Figure 5-1 shows the packet delivery ratio for all protocols when the number of sources is five and each source sends one packet per second. The figure shows that all VCAR protocols outperform AODV for all pause times. When the pause time is equal to zero (high mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 13.94, 9.27, 4.89, and 1.84 percent, respectively. When pause time is equal to 300 seconds (low mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 4.6, 4.44, 2.95, and .61 percent, respectively.

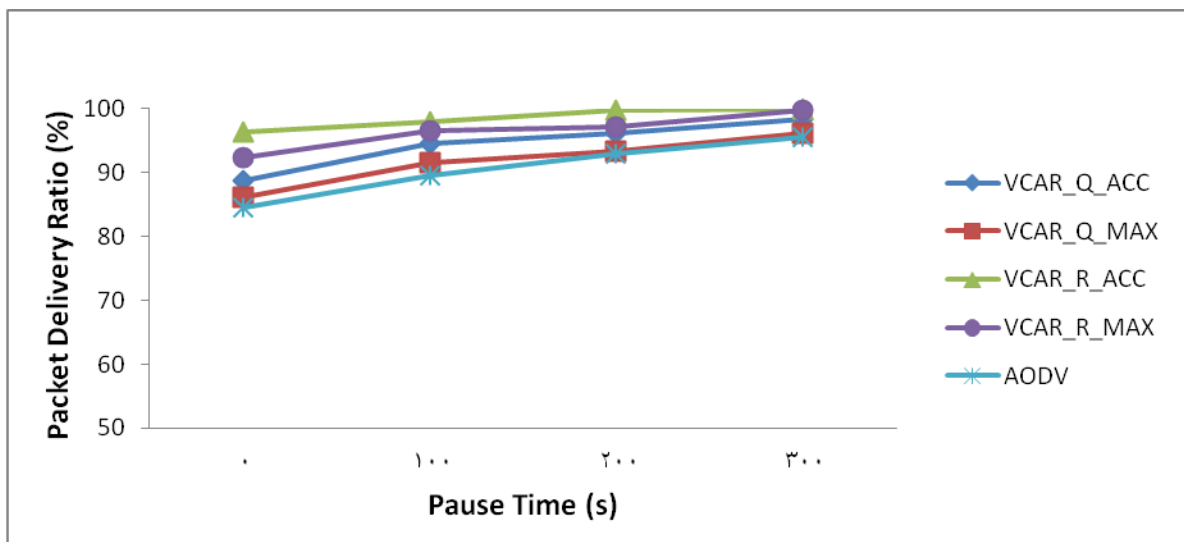


Figure 5-1: Packet Delivery Ratio of 5 sources each one sends 1 packet/s

Figure 5-2 shows the packet delivery ratio for all protocols when the number of sources is five and each source sends two packets per second. The figure shows that all VCAR protocols outperform AODV for all pause times. When the pause time is equal to zero (high mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 15.14, 10.43, 5.99, and 2.91 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), VCAR_R_ACC,

VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 10.67, 8.33, and 4.35 percent, respectively.

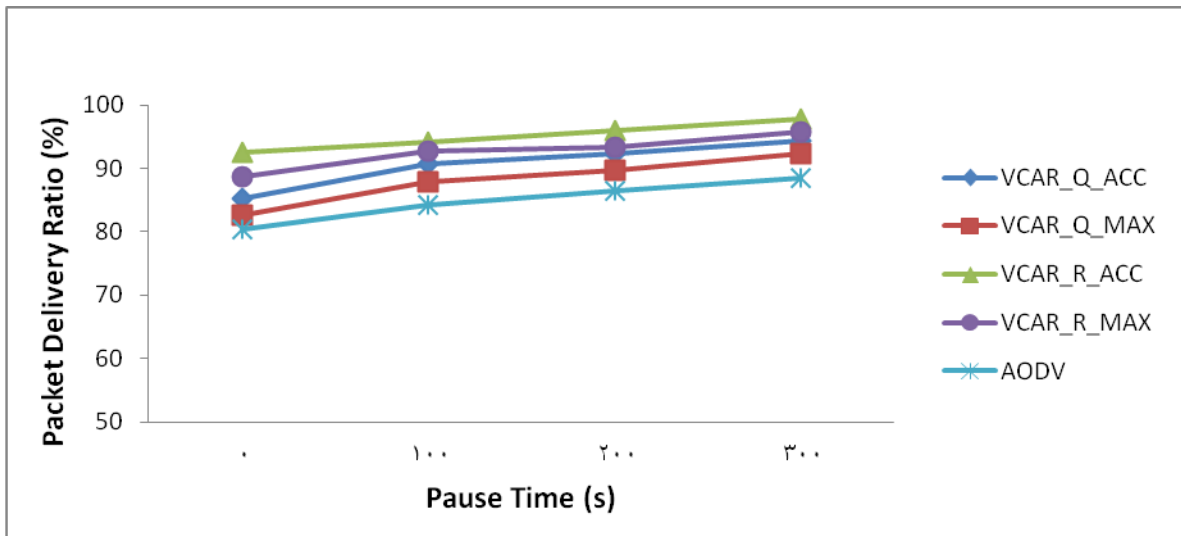


Figure 5-2: Packet Delivery Ratio of 5 sources each one sends 2 packets/s

Figure 5-3 shows the packet delivery ratio for all protocols when the number of sources is five and each source sends four packets per second. The figure shows that all VCAR protocols outperform AODV for all pause times. When the pause time is equal to zero (high mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 15.32, 10.60, 6.16, and 3.07 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 8.75, 6.45, 4.93, and 2.54 percent, respectively.

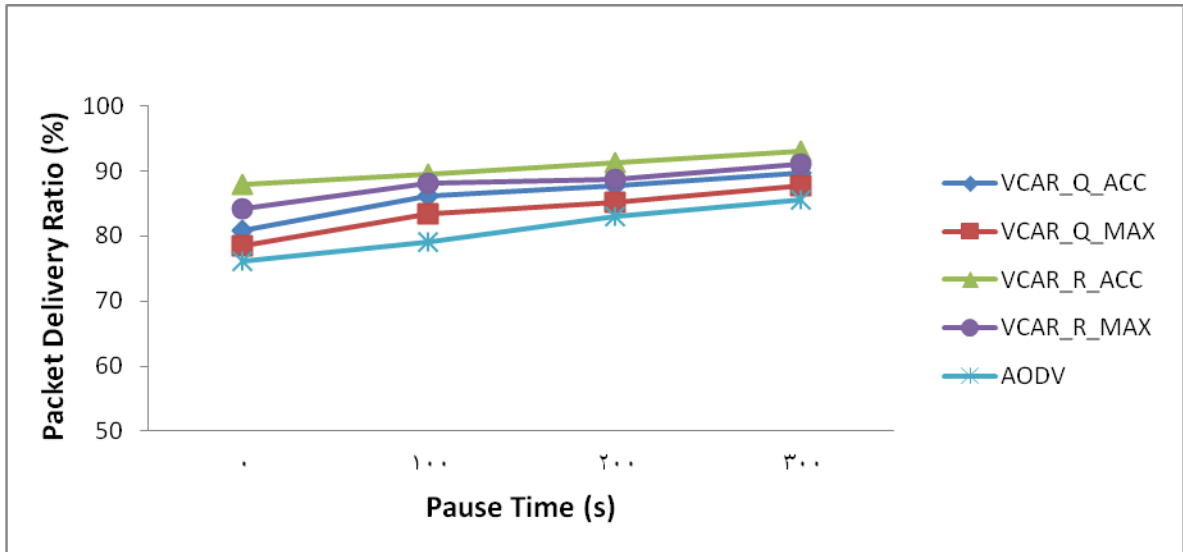


Figure 5-3: Packet Delivery Ratio of 5 sources each one sends 4 packets/s

Figure 5-4 shows the packet delivery ratio for all protocols when the number of sources is five and each source sends six packets per second. The figure shows that all VCAR protocols outperform AODV for all pause times. When the pause time is equal to zero (high mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 14.80, 10.19, 5.77, and 2.69 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 10.82, 8.47, 6.92, and 4.49 percent, respectively.

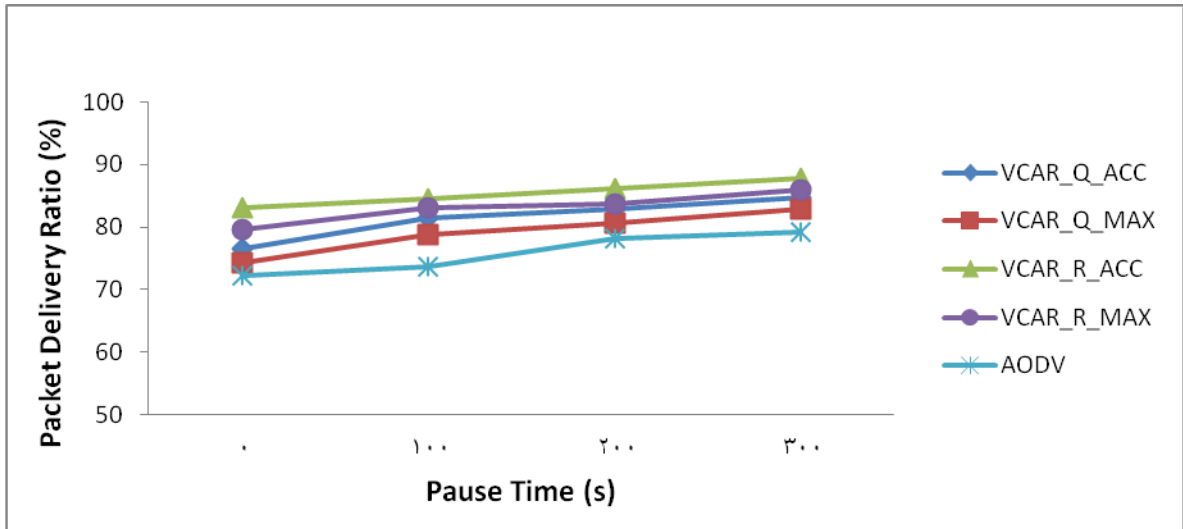


Figure 5-4: Packet Delivery Ratio of 5 sources each one sends 6 packets/s

Figure 5-5 shows the packet delivery ratio for all protocols when the number of sources is ten and each source sends one packet per second. The figure shows that all VCAR protocols outperform AODV for all pause times. When the pause time is equal to zero (high mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 14.73, 10.03, 5.62, and 2.54 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 6.14, 5.24, 3.74, and 1.37 percent, respectively.

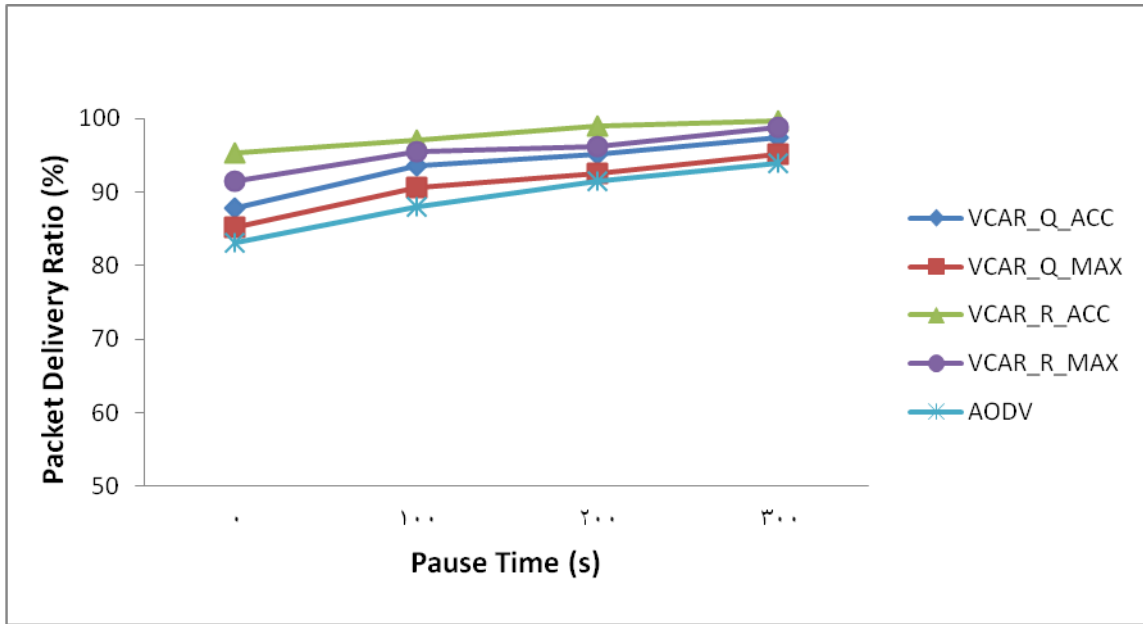


Figure 5-5: Packet Delivery Ratio of 10 sources each one sends 1 packet/s

Figure 5-6 shows the packet delivery ratio for all protocols when the number of sources is ten and each source sends two packets per second. The figure shows that all VCAR protocols outperform AODV for all pause times. When the pause time is equal to zero (high mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 22.43, 17.42, 12.71, and 4.29 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 10.24, 9.30, 7.74, and 5.28 percent, respectively.

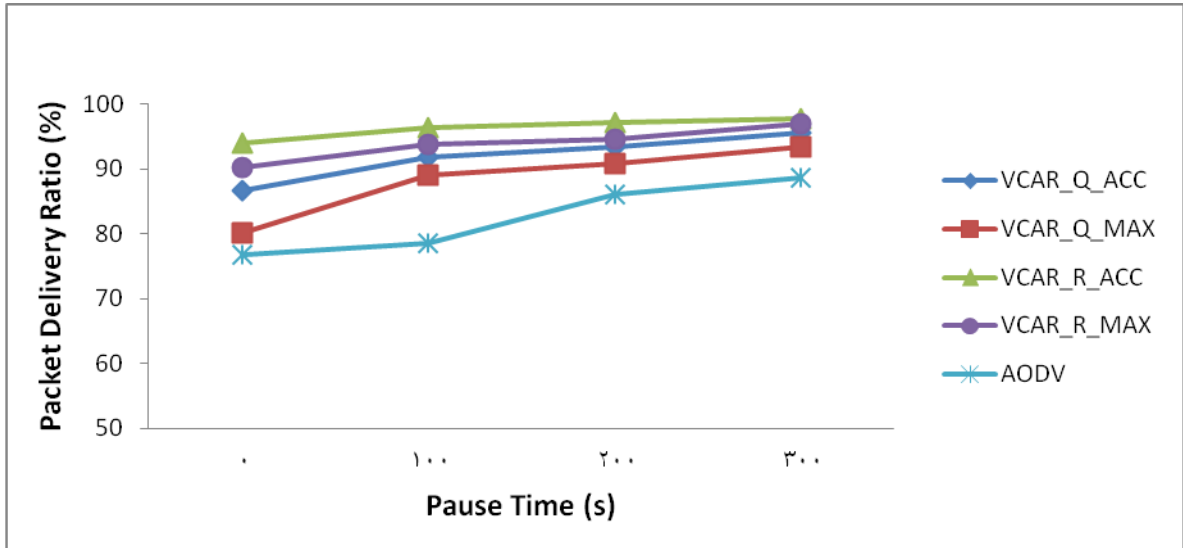


Figure 5-6: Packet Delivery Ratio of 10 sources each one sends 2 packets/s

Figure 5-7 shows the packet delivery ratio for all protocols when the number of sources is ten and each source sends four packets per second. The figure shows that all VCAR protocols outperform AODV for all pause times. When the pause time is equal to zero (high mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 28.26, 19.41, 14.66, and 3.19 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 23.19, 18.69, 15.79, and 11.12 percent, respectively.

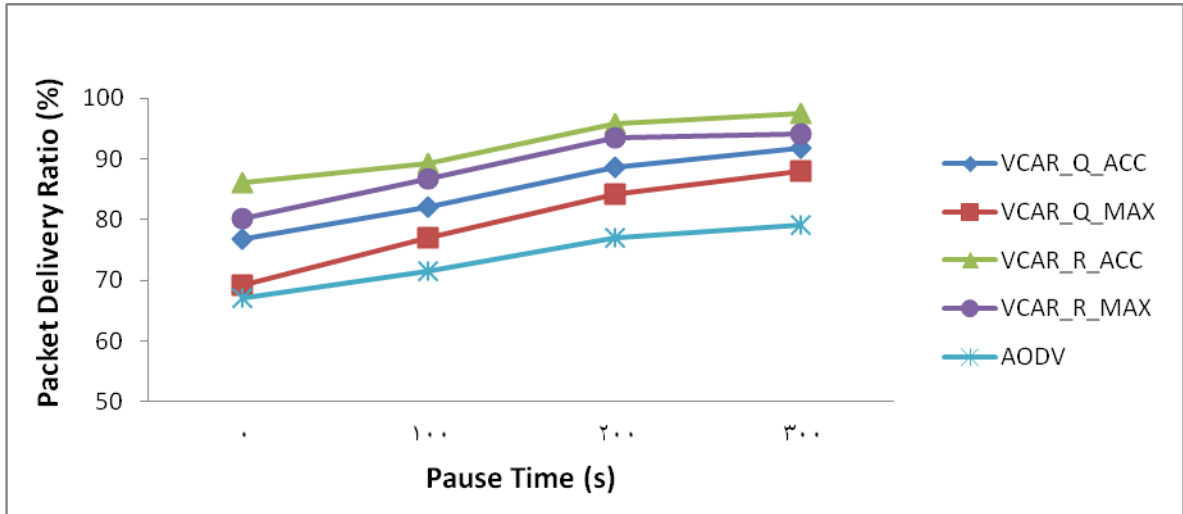


Figure 5-7: Packet Delivery Ratio of 10 sources each one sends 4 packets/s

Figure 5-8 shows the packet delivery ratio for all protocols when the number of sources is ten and each source sends six packets per second. The figure shows that all VCAR protocols outperform AODV for all pause times. When the pause time is equal to zero (high mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 31.35, 21.04, 20.03, and 5.46 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 26.69, 19.68, 16.39, and 11.11 percent, respectively.

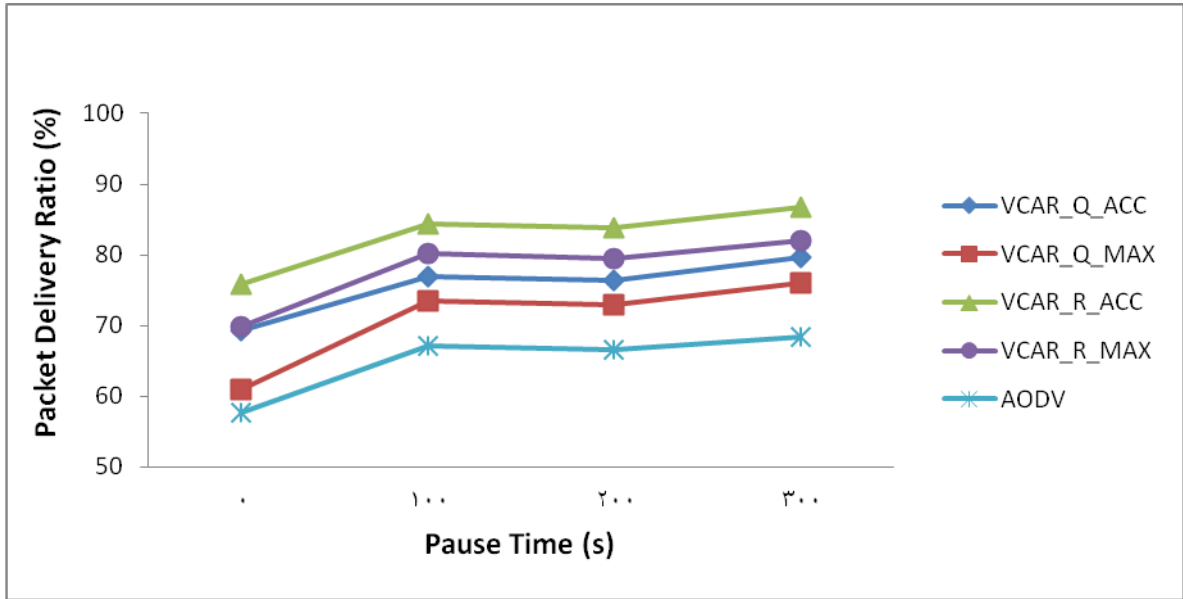


Figure 5-8: Packet Delivery Ratio of 10 sources each one sends 6 packets/s

Figure 5-9 shows the packet delivery ratio for all protocols when the number of sources is fifteen and each source sends one packet per second. The figure shows that all VCAR protocols outperform AODV for all pause times. When the pause time is equal to zero (high mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 22.36, 17.35, 12.64, and 4.23 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 18.84, 15.47, 14.72, and 10.67 percent, respectively.

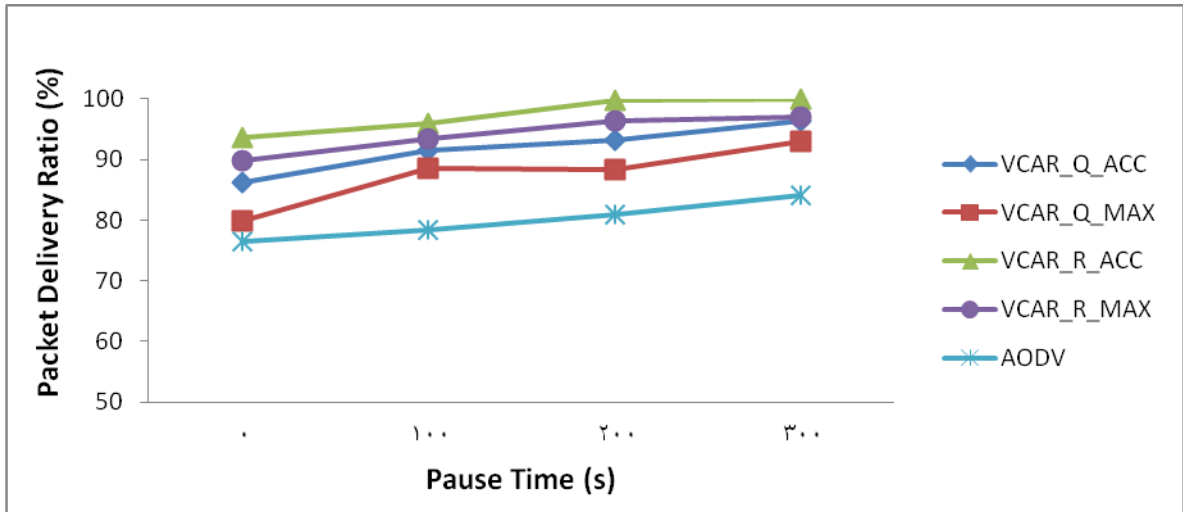


Figure 5-9: Packet Delivery Ratio of 15 sources each one sends 1 packet/s

Figure 5-10 shows the packet delivery ratio for all protocols when the number of sources is fifteen and each source sends two packets per second. The figure shows that all VCAR protocols outperform AODV for all pause times. When the pause time is equal to zero (high mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_MAX, and VCAR_Q_ACC outperform AODV by 33.71, 18.95, 18.29, and 2.80 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 26.57, 18.43, 15.54, and 10.87 percent, respectively.

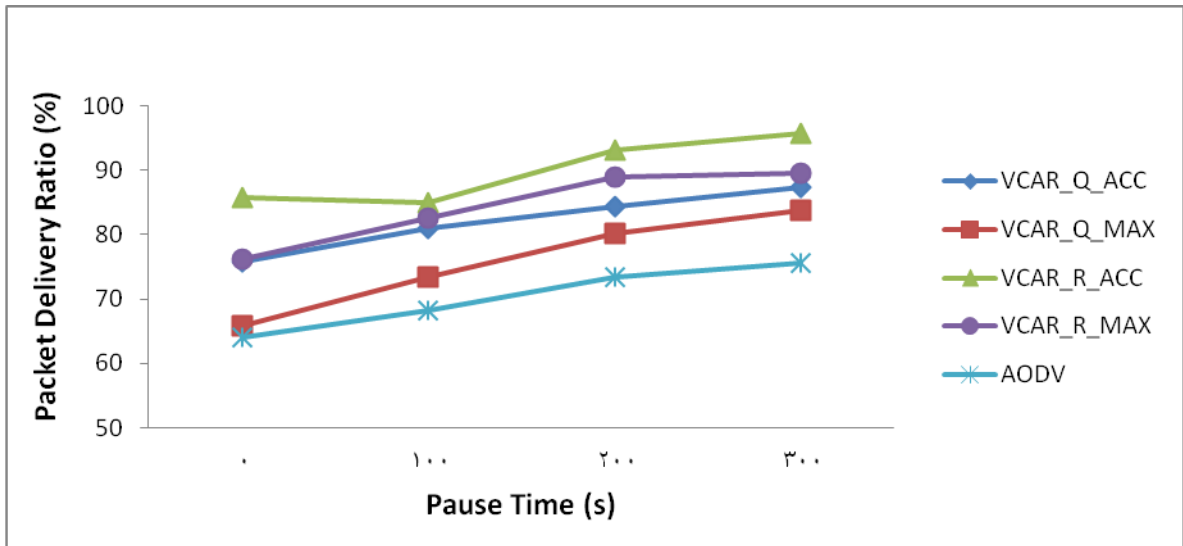


Figure 5-10: Packet Delivery Ratio of 15 sources each one sends 2 packets/s

Figure 5-11 shows the packet delivery ratio for all protocols when the number of sources is fifteen and each source sends four packets per second. The figure shows that all VCAR protocols outperform AODV for all pause times. When the pause time is equal to zero (high mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 30.90, 20.68, 19.73, and 5.25 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 26.72, 19.64, 16.35, and 11.07 percent, respectively.

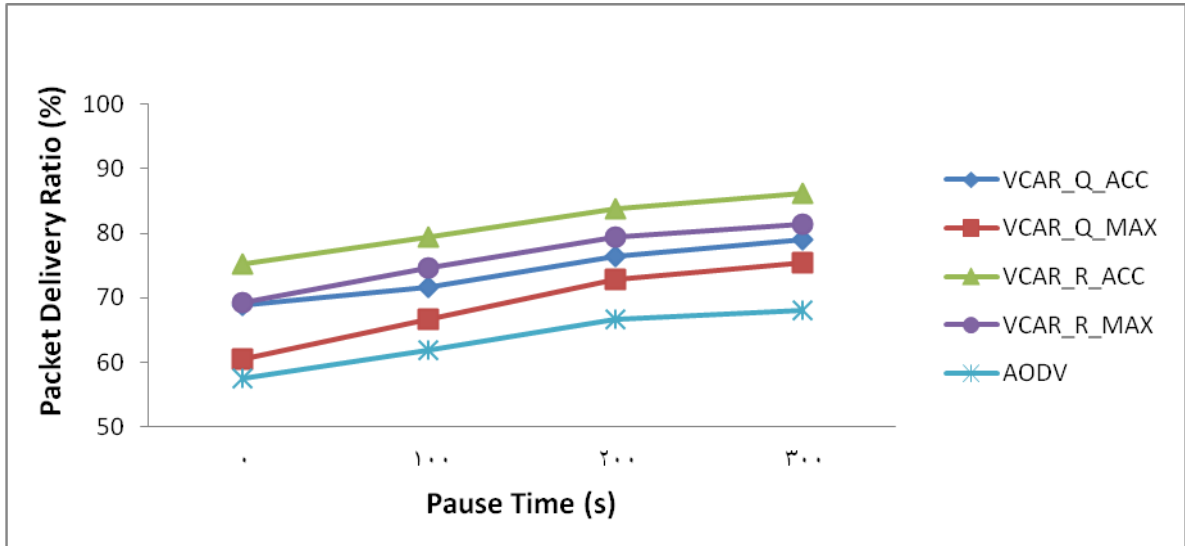


Figure 5-11: Packet Delivery Ratio of 15 sources each one sends 4 packets/s

Figure 5-12 shows the packet delivery ratio for all protocols when the number of sources is fifteen and each source sends six packets per second. The figure shows that all VCAR protocols outperform AODV for all pause times. When the pause time is equal to zero (high mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 27.00, 21.34, 18.24, and 6.80 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 25.73, 19.79, 16.12, and 10.27 percent, respectively.

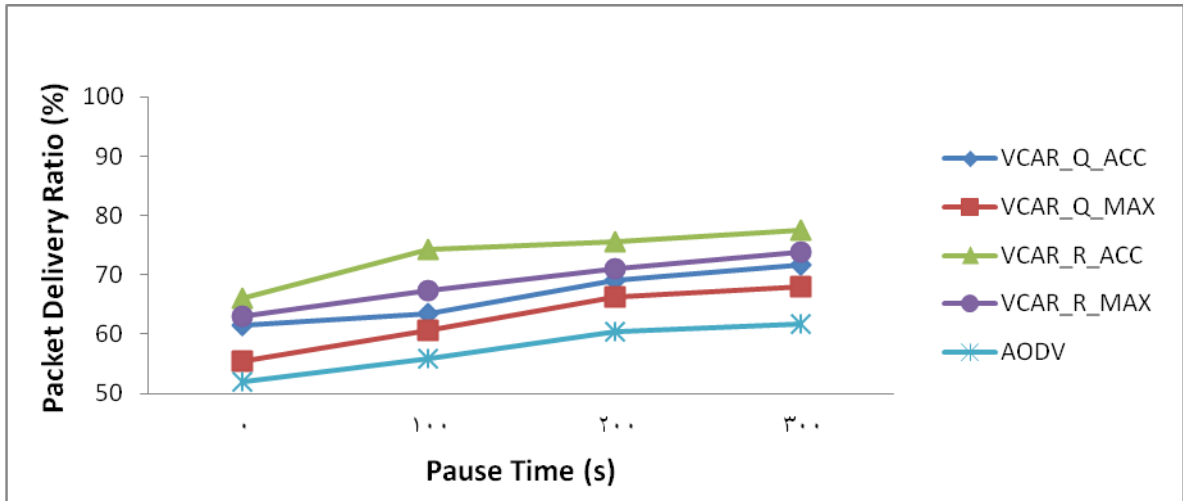


Figure 5-12: Packet Delivery Ratio of 15 sources each one sends 6 packets/s

5-2 Average End-to-End Delay.

Figures 5-13~5-24 show the average end-to-end delay for different transmission rate values and a different number of sources. The results also show that all VCAR protocols have a lower average end to end delay than the AODV protocol for all mobility levels, all number of sources, and all transmission rates. This is because AODV always uses the shortest route even if it is highly congested. While VCAR protocols use less congested routes. Another reason is that VCAR protocols use the most stable route, while AODV does not, so that the route selected by AODV is more likely to break than the route selected by VCAR protocols.

VCAR_R_ACC outperforms all other VCAR protocols in all the cases we have studied. On average, VCAR_R_ACC outperforms AODV by 28.86 percent for low mobility networks, and by 32.77 for high mobility networks.

Figure 5-13 shows the average end-to-end delay for all protocols when the number of sources is five and each source sends one packet per second. The figure shows that all VCAR protocols outperform AODV for all pause times. When the pause time is equal to zero (high mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 32.25, 26.97, 18.11, and 11.03 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 31.04, 26.85, 26.34, and 21.55 percent, respectively.

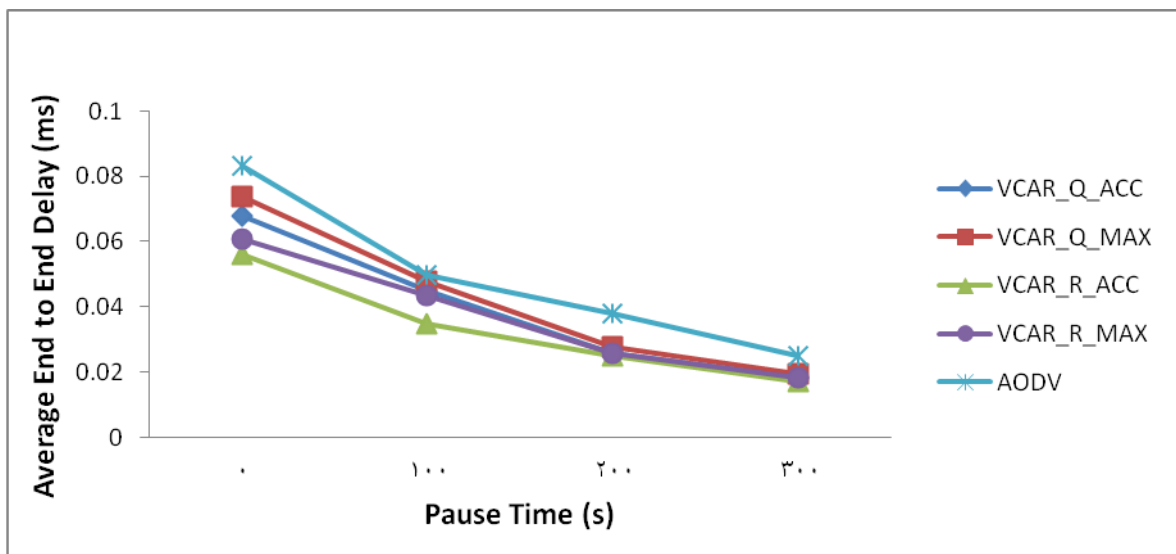


Figure 5-13: Average End to End Delay of 5 sources each one sends 1 packet/s

Figure 5-14 shows the average end-to-end delay for all protocols when the number of sources is five and each source sends two packets per second. The figure shows that all VCAR protocols outperform AODV for all pause times. When the pause time is equal to zero (high mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 28.87, 23.31, 14.01, and 6.59 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), VCAR_R_ACC,

VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 29.93, 25.67, 25.15, and 20.29 percent, respectively.

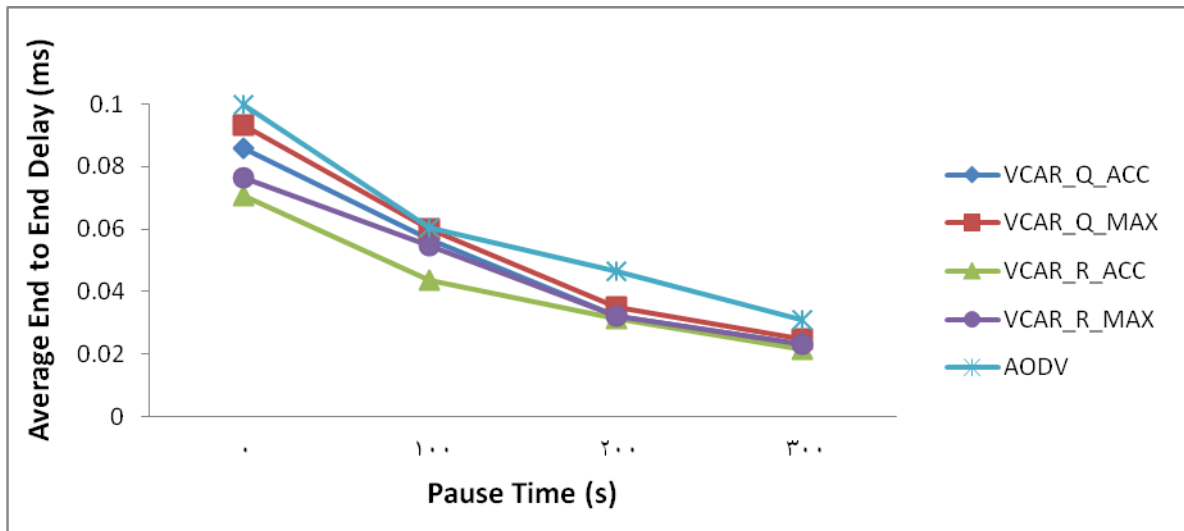


Figure 5-14: Average End to End Delay of 5 sources each one sends 2 packets/s

Figure 5-15 shows the average end-to-end delay for all protocols when the number of sources is five and each source sends four packets per second. The figure shows that all VCAR protocols outperform AODV for all pause times. When the pause time is equal to zero (high mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 28.36, 22.77, 13.40, and 5.93 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 27.78, 23.39, 22.86, and 17.84 percent, respectively.

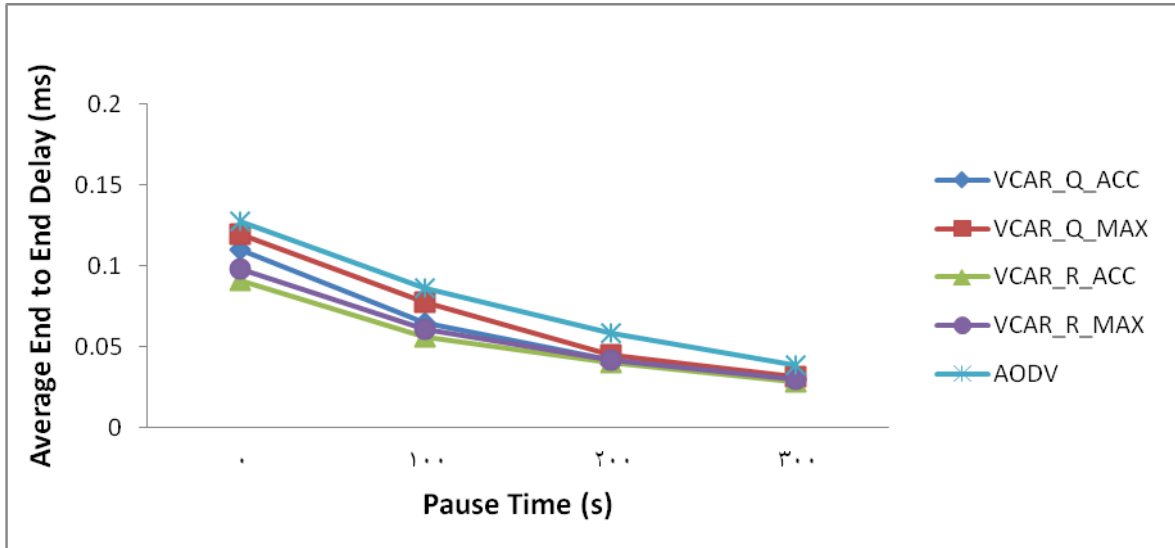


Figure 5-15: Average End to End Delay of 5 sources each one sends 4 packets/s

Figure 5-16 shows the average end-to-end delay for all protocols when the number of sources is five and each source sends six packets per second. The figure shows that all VCAR protocols outperform AODV for all pause times. When the pause time is equal to zero (high mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 27.80, 22.17, 12.73, and 5.20 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 27.14, 22.71, 22.18, and 17.11 percent, respectively.

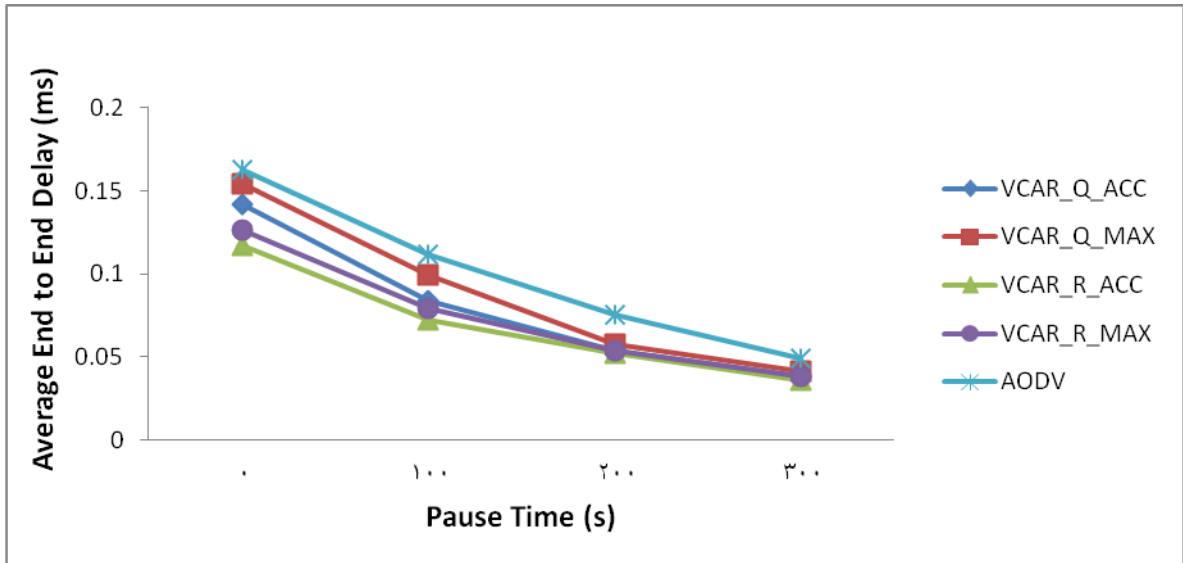


Figure 5-16: Average End to End Delay of 5 sources each one sends 6 packets/s

Figure 5-17 shows the average end-to-end delay for all protocols when the number of sources is ten and each source sends one packet per second. The figure shows that all VCAR protocols outperform AODV for all pause times. When the pause time is equal to zero (high mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 30.84, 25.44, 16.40, and 9.18 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 33.91, 29.90, 29.41, and 24.82 percent, respectively.

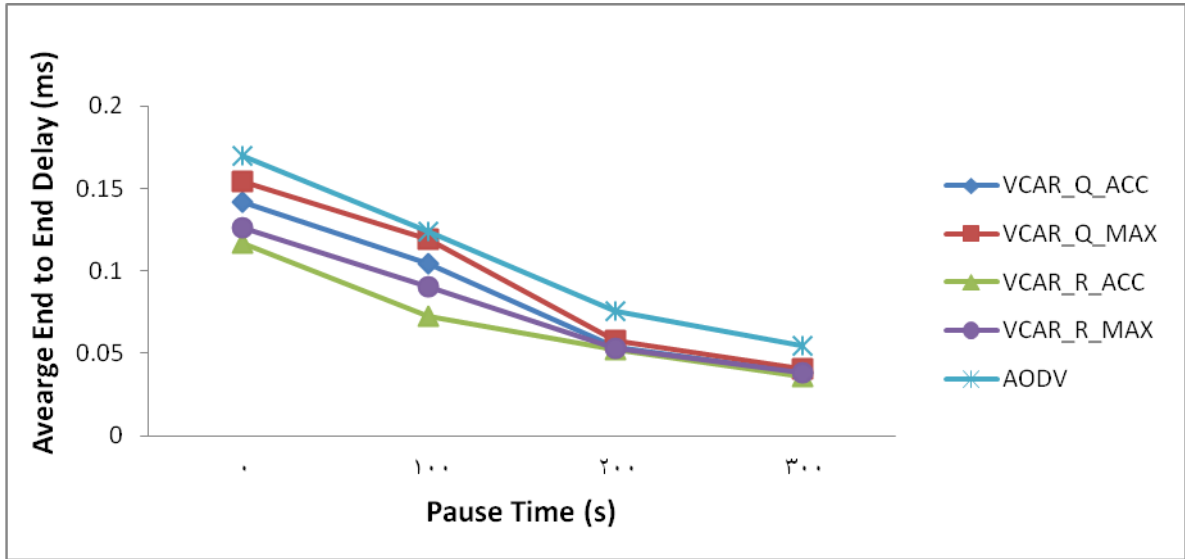


Figure 5-17: Average End to End Delay of 10 sources each one sends 1 packet/s

Figure 5-18 shows the average end-to-end delay for all protocols when the number of sources is ten and each source sends two packets per second. The figure shows that all VCAR protocols outperform AODV for all pause times. When the pause time is equal to zero (high mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 26.52, 20.78, 11.17, and 3.51 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 22.90, 18.21, 17.64, and 12.29 percent, respectively.

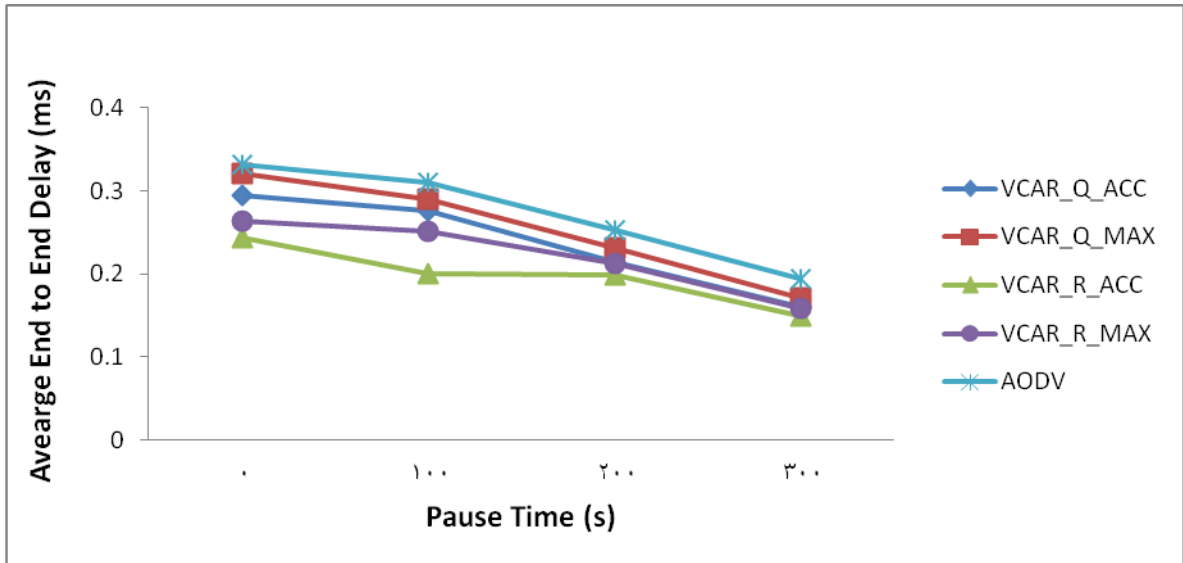


Figure 5-18: Average End to End Delay of 10 sources each one sends 2 packets/s

Figure 5-19 shows the average end-to-end delay for all protocols when the number of sources is ten and each source sends four packets per second. The figure shows that all VCAR protocols outperform AODV for all pause times. When the pause time is equal to zero (high mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 27.10, 23.45, 15.27, and 6.76 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 28.51, 22.60, 20.33, and 14.22 percent, respectively.

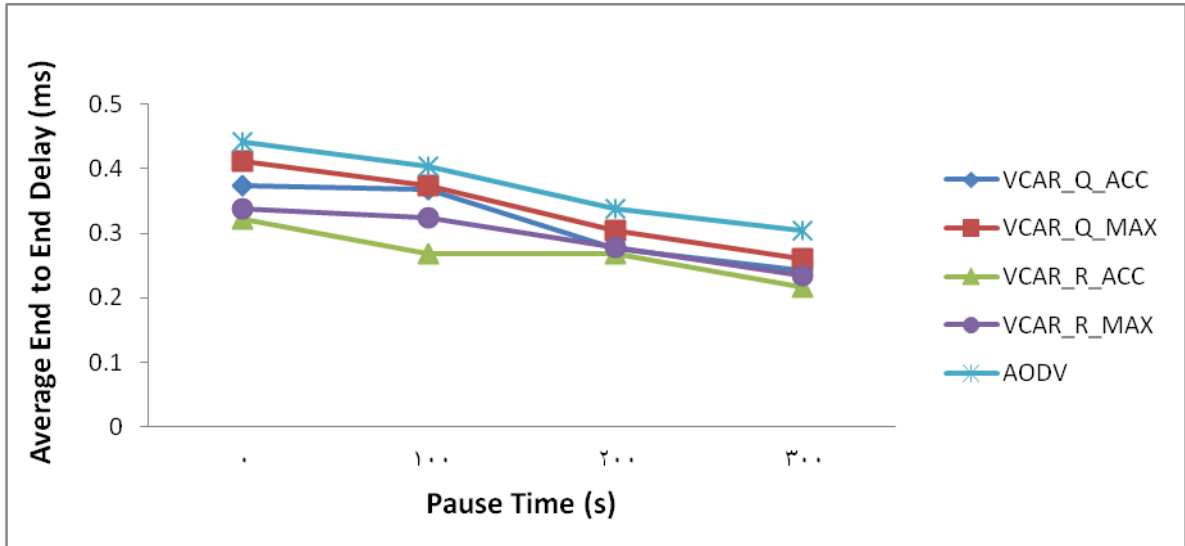


Figure 5-19: Average End to End Delay of 10 sources each one sends 4 packets/s

Figure 5-20 shows the average end-to-end delay for all protocols when the number of sources is ten and each source sends six packets per second. The figure shows that all VCAR protocols outperform AODV for all pause times. When the pause time is equal to zero (high mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 35.03, 32.79, 25.73, and 16.30 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 33.29, 25.88, 21.55, and 14.12 percent, respectively.

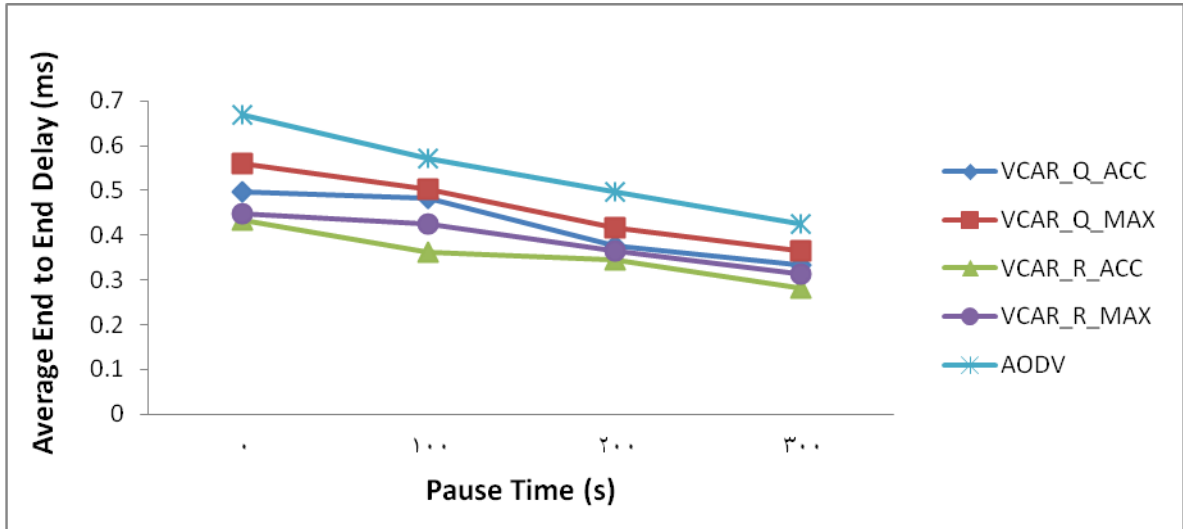


Figure 5-20: Average End to End Delay of 10 sources each one sends 6 packets/s

Figure 5-21 shows the average end-to-end delay for all protocols when the number of sources is fifteen and each source sends one packet per second. The figure shows that all VCAR protocols outperform AODV for all pause times. When the pause time is equal to zero (high mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 34.37, 29.25, 20.66, and 13.81 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 23.43, 18.77, 18.21, and 12.89 percent, respectively.

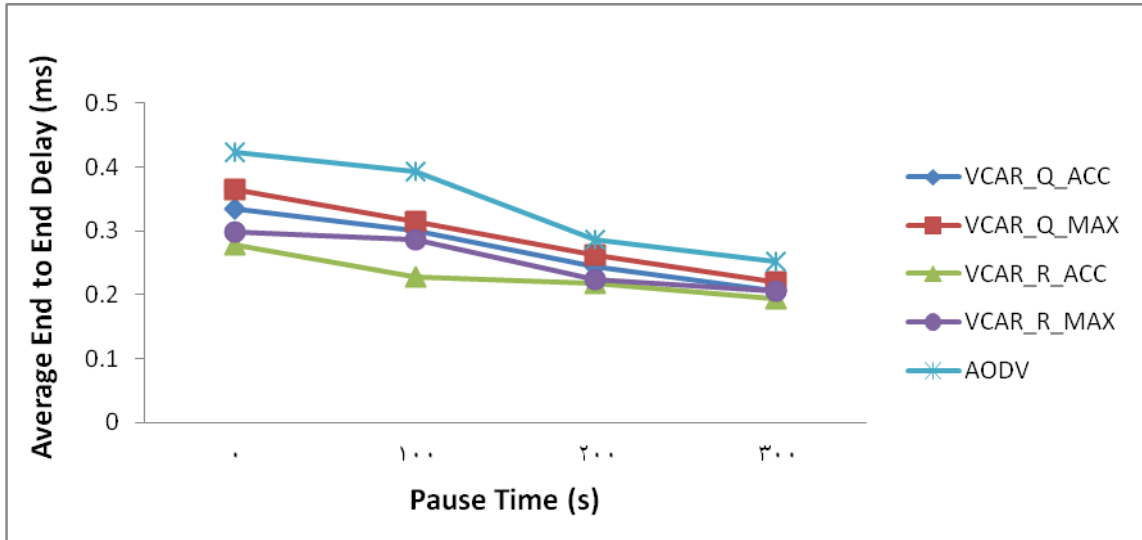


Figure 5-21: Average End to End Delay of 15 sources each one sends 1 packet/s

Figure 5-22 shows the average end-to-end delay for all protocols when the number of sources is fifteen and each source sends two packets per second. The figure shows that all VCAR protocols outperform AODV for all pause times. When the pause time is equal to zero (high mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 32.97, 29.61, 22.09, and 14.26 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 27.69, 21.41, 19.41, and 13.23 percent, respectively.

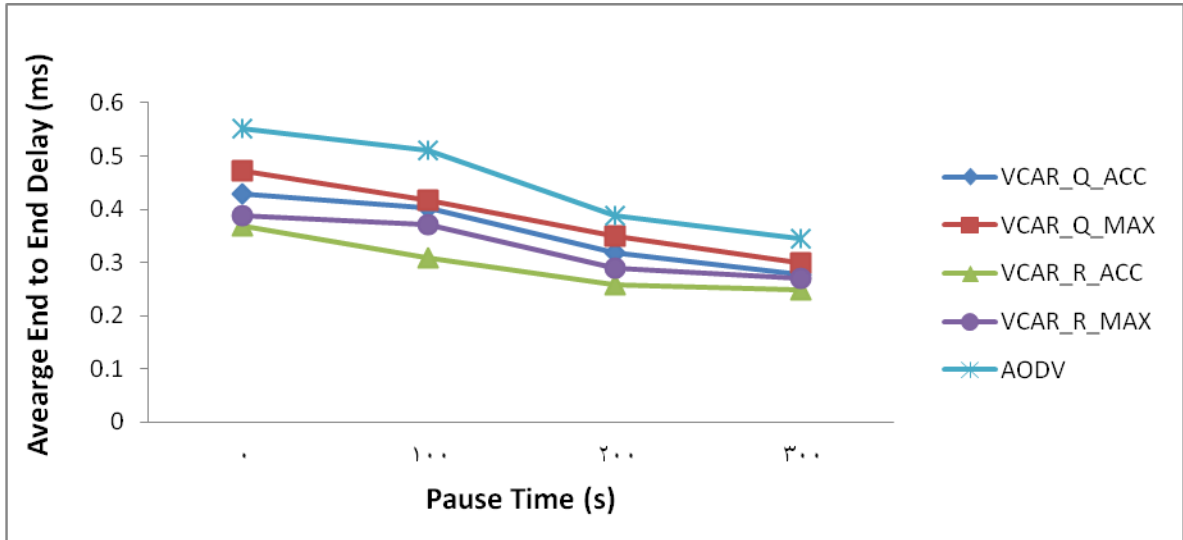


Figure 5-22: Average End to End Delay of 15 sources each one sends 2 packets/s

Figure 5-23 shows the average end-to-end delay for all protocols when the number of sources is fifteen and each source sends four packets per second. The figure shows that all VCAR protocols outperform AODV for all pause times. When the pause time is equal to zero (high mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 34.14, 30.53, 24.08, and 15.38 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 31.77, 24.62, 20.67, 13.64 percent, respectively.

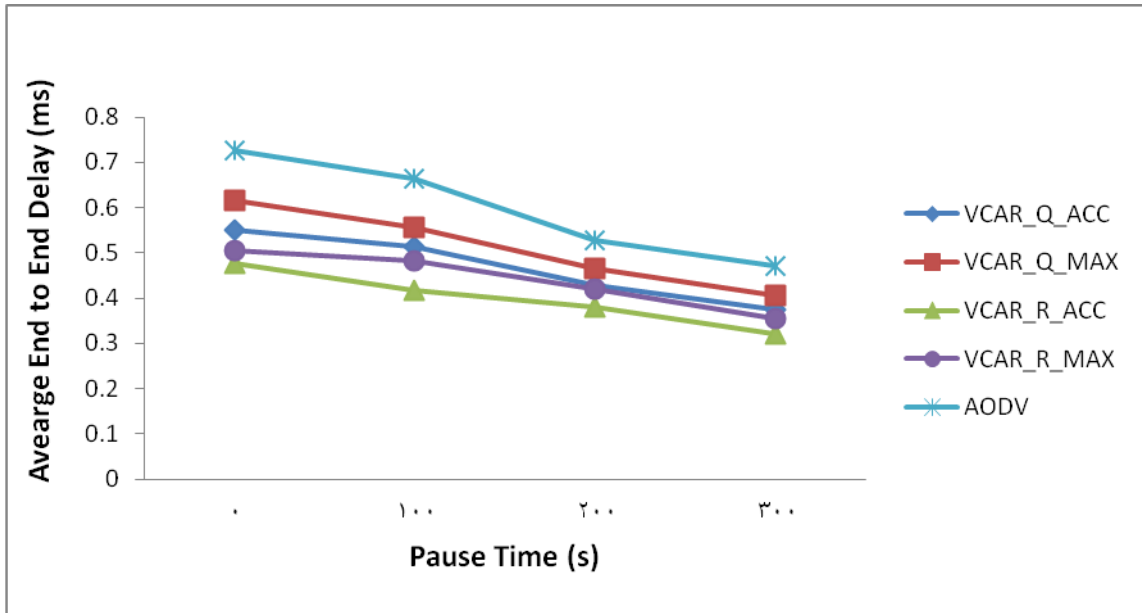


Figure 5-23: Average End to End Delay of 15 sources each one sends 4 packets/s

Figure 5-24 shows the average end-to-end delay for all protocols when the number of sources is fifteen and each source sends six packets per second. The figure shows that all VCAR protocols outperform AODV for all pause times. When the pause time is equal to zero (high mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 36.32, 32.33, 27.00, and 17.57 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 28.24, 23.68, 19.80, and 13.91 percent, respectively.

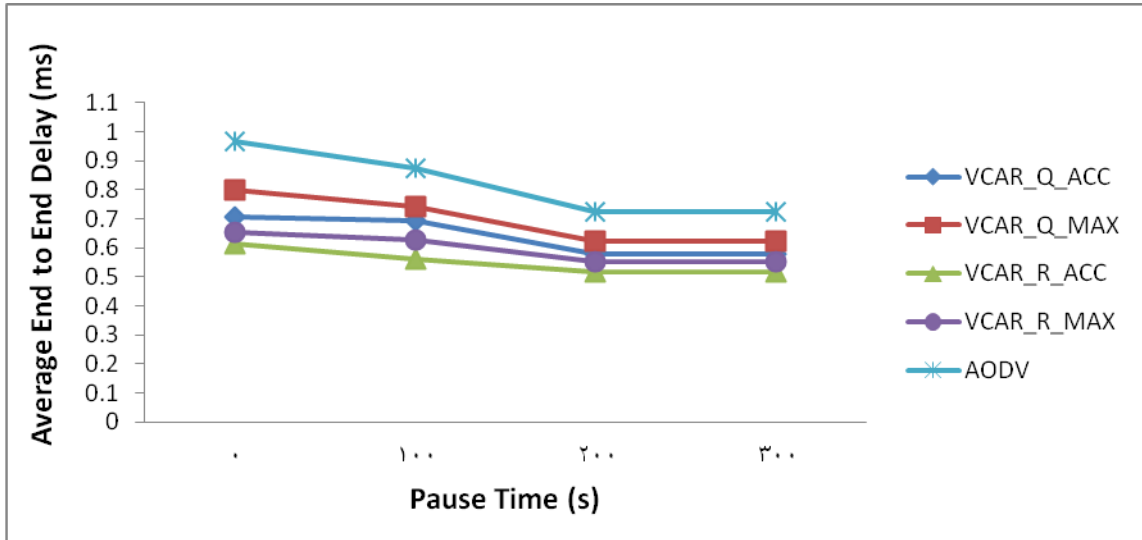


Figure 5-24: Average End to End Delay of 15 sources each one sends 6 packets/s

5-3 Routing Overhead

Figures 5-25~5-36 show the routing overhead for all the simulation scenarios. The results show that all VCAR protocols have a higher routing overhead than AODV especially in high mobility scenarios. In VCAR, the destination may send multiple replies, because it always searches for the best path among all available paths. This behavior increases the number of RREP packets, and consequently increases the routing overhead.

Figure 5-25 shows the routing overhead for all protocols when the number of sources is five and each source sends one packet per second. The figure shows that AODV outperforms all VCAR protocols for all pause times. When the pause time is equal to zero (high mobility network), AODV outperforms VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX by 2.70, 3.64, 6.00, 6.43 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), AODV outperforms

VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX by 4.28, 5.17, 8.60, 11.20 percent, respectively.

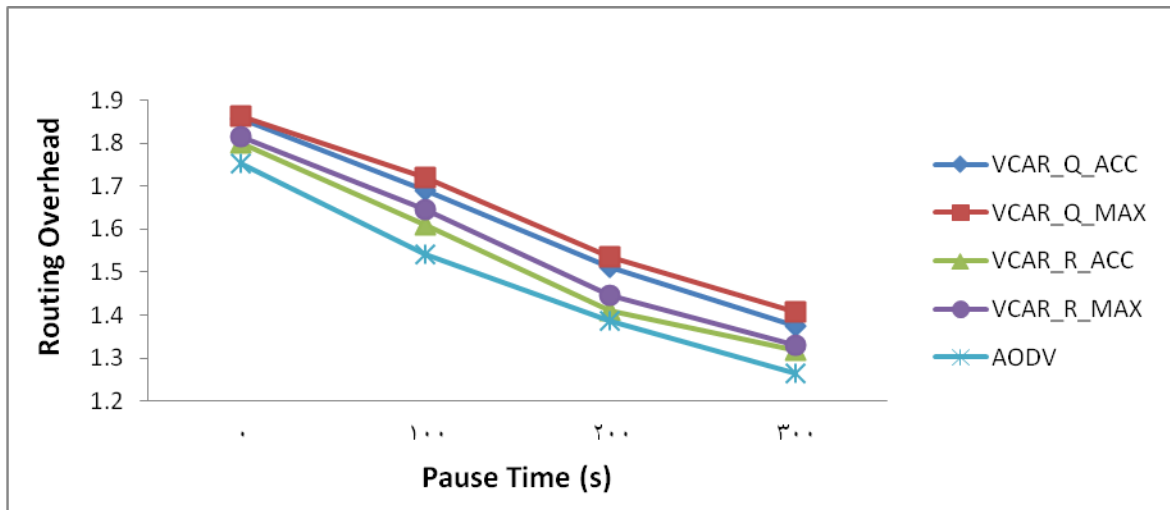


Figure 5-25: Routing Overhead of 5 sources each one sends 1 packet/s

Figure 5-26 shows the routing overhead for all protocols when the number of sources is five and each source sends two packets per second. The figure shows that AODV outperforms all VCAR protocols for all pause times. When the pause time is equal to zero (high mobility network), AODV outperforms VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX by 2.89, 3.83, 6.20, and 6.62 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), AODV outperforms VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX by 4.38, 5.28, 8.71, and 11.31 percent, respectively.

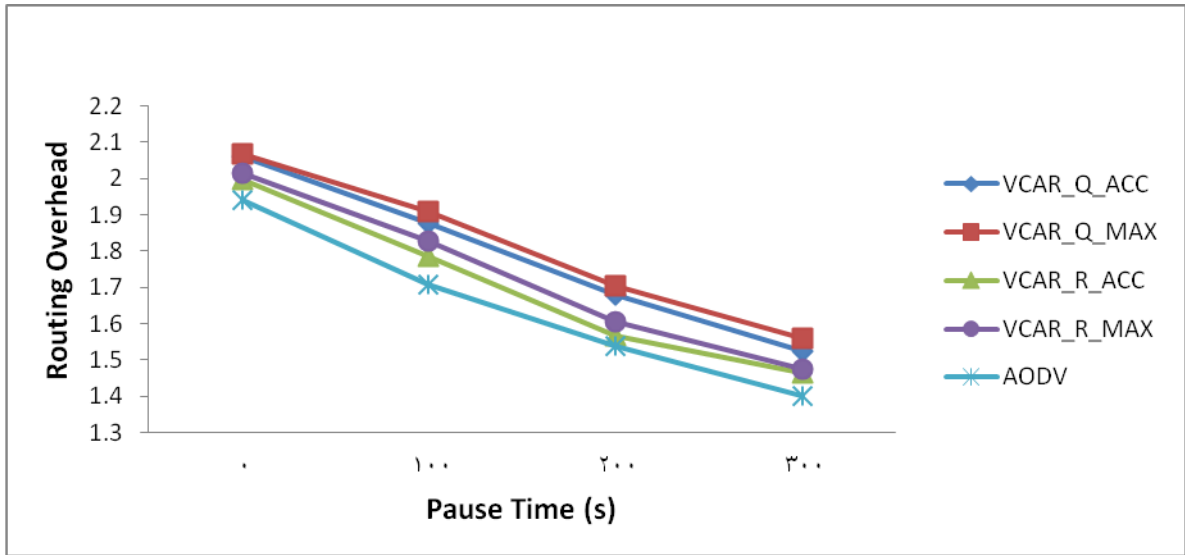


Figure 5-26: Routing Overhead of 5 sources each one sends 2 packets/s

Figure 5-27 shows the routing overhead for all protocols when the number of sources is five and each source sends four packets per second. The figure shows that AODV outperforms all VCAR protocols for all pause times. When the pause time is equal to zero (high mobility network), AODV outperforms VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX by 3.00, 3.95, 6.31, and 6.74 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), AODV outperforms VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX by 3.30, 4.18, 7.58, and 10.16 percent, respectively.

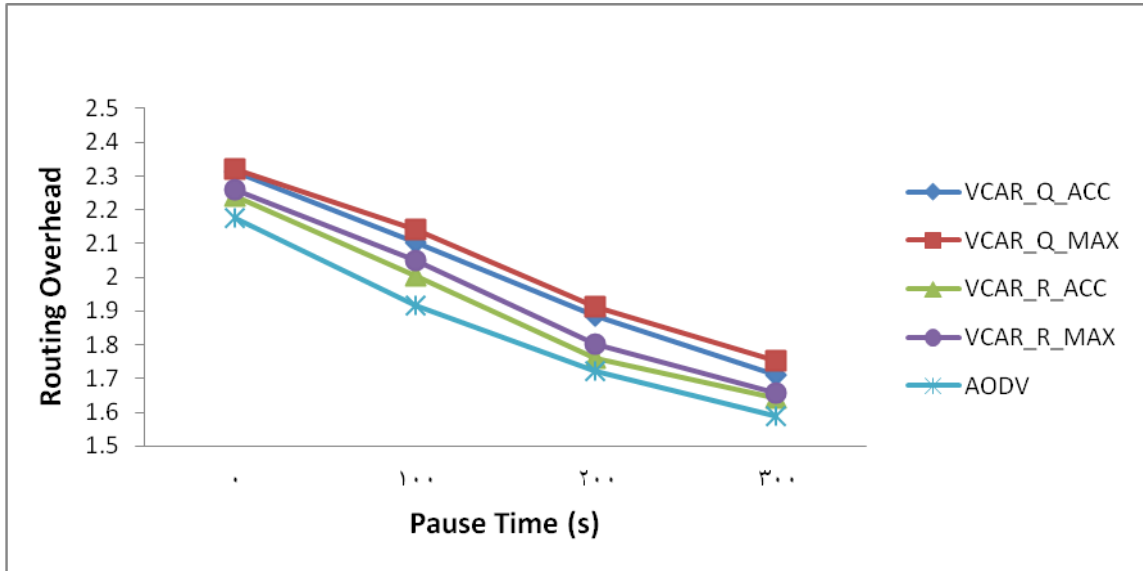


Figure 5-27: Routing Overhead of 5 sources each one sends 4 packets/s

Figure 5-28 shows the routing overhead for all protocols when the number of sources is five and each source sends six packets per second. The figure shows that AODV outperforms all VCAR protocols for all pause times. When the pause time is equal to zero (high mobility network), AODV outperforms VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX by 3.12, 4.07, 6.44, and 6.86 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), AODV outperforms VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX by 3.04, 3.92, 7.32, and 9.89 percent, respectively.

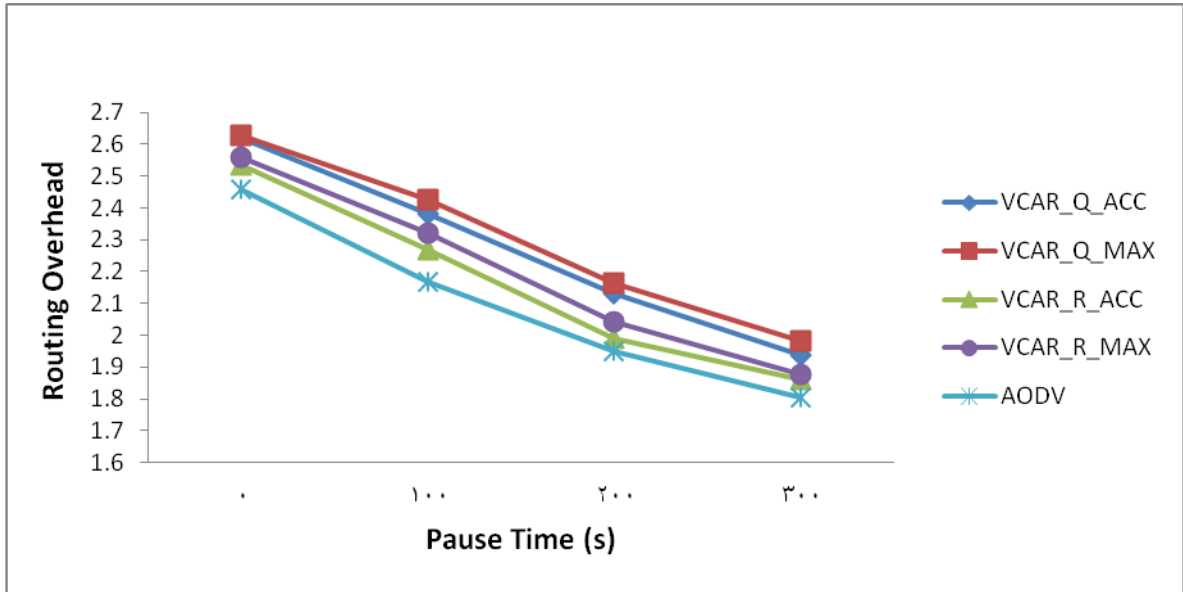


Figure 5-28: Routing Overhead of 5 sources each one sends 6 packets/s

Figure 5-29 shows the routing overhead for all protocols when the number of sources is ten and each source sends one packet per second. The figure shows that AODV outperforms all VCAR protocols for all pause times. When the pause time is equal to zero (high mobility network), AODV outperforms VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX by 11.00, 12.02, 14.28, and 15.03 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), AODV outperforms VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX by 3.56, 4.47, 7.88, and 10.46 percent, respectively.

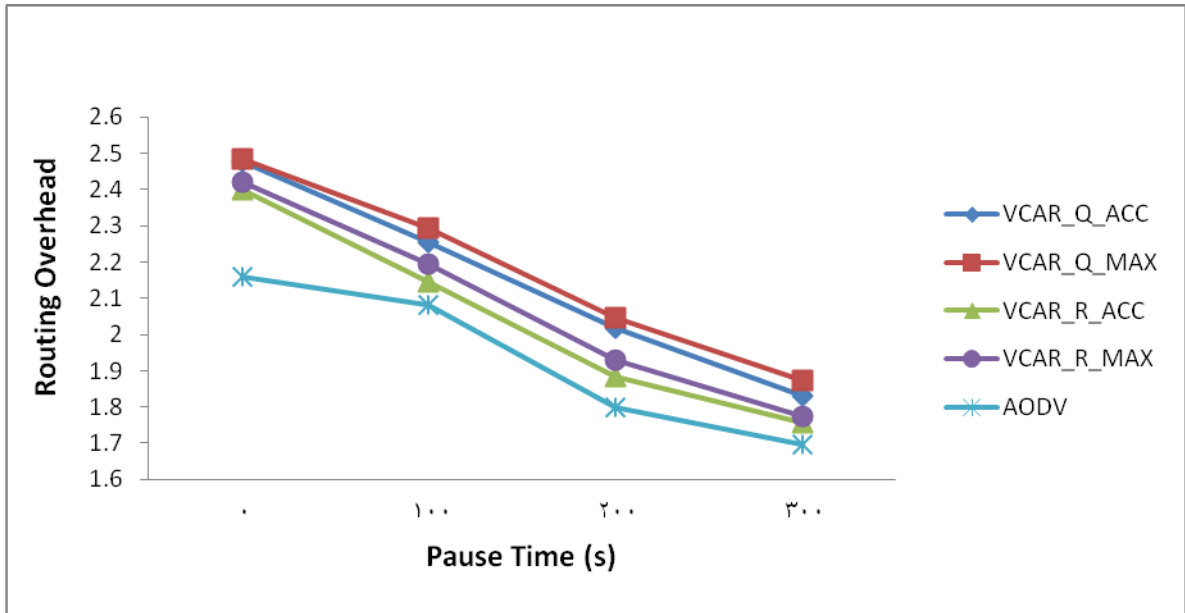


Figure 5-29: Routing Overhead of 10 sources each one sends 1 packet/s

Figure 5-30 shows the routing overhead for all protocols when the number of sources is ten and each source sends two packets per second. The figure shows that AODV outperforms all VCAR protocols for all pause times. When the pause time is equal to zero (high mobility network), AODV outperforms VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX by 13.56, 14.59, 17.21, and 17.67 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), AODV outperforms VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX by 1.06, 1.92, 5.27, and 7.81 percent, respectively.

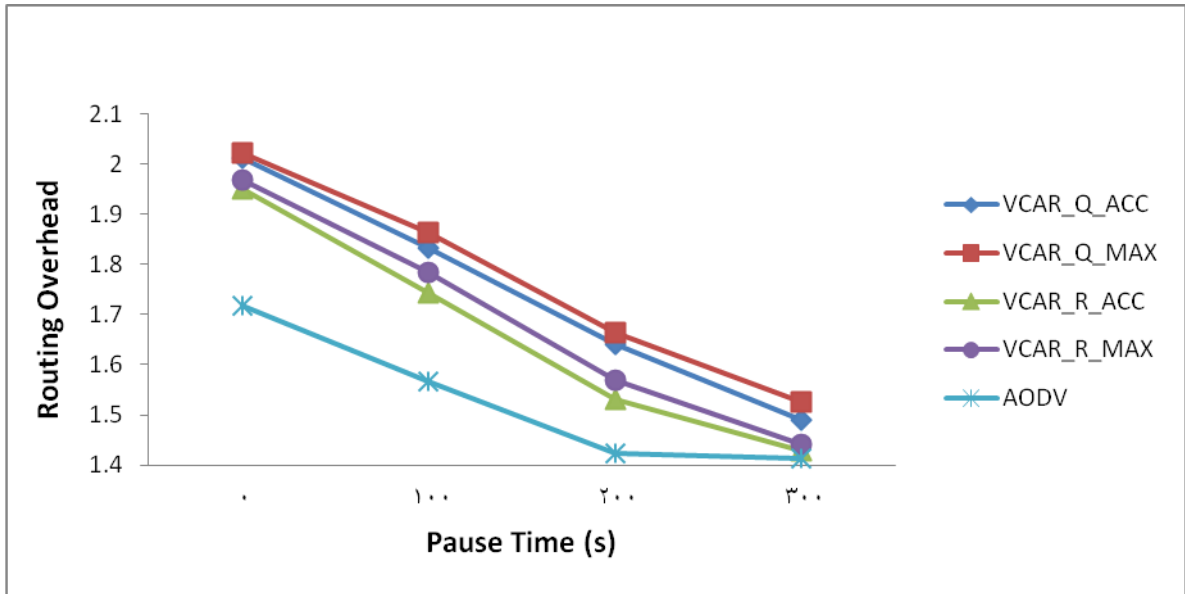


Figure 5-30: Routing Overhead of 10 sources each one sends 2 packets/s

Figure 5-31 shows the routing overhead for all protocols when the number of sources is ten and each source sends four packets per second. The figure shows that AODV outperforms all VCAR protocols for all pause times. When the pause time is equal to zero (high mobility network), AODV outperforms VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX by 1.17, 2.11, 4.47, and 4.90 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), AODV outperforms VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX by 1.97, 3.53, 6.84, and 9.36 percent, respectively.

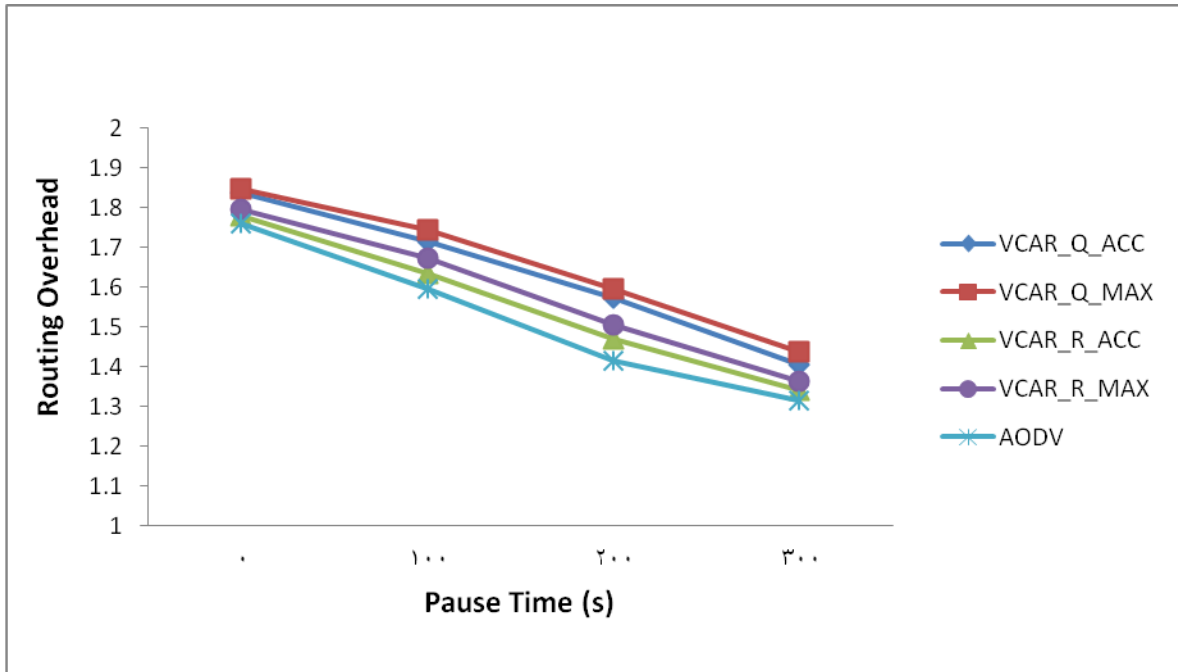


Figure 5-31: Routing Overhead of 10 sources each one sends 4 packets/s

Figure 5-32 shows the routing overhead for all protocols when the number of sources is ten and each source sends six packets per second. The figure shows that AODV outperforms all VCAR protocols for all pause times. When the pause time is equal to zero (high mobility network), AODV outperforms VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX by 3.87, 6.52, 12.84, and 14.00 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), AODV outperforms VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX by 1.96, 10.79, 19.92, and 26.92 percent, respectively.

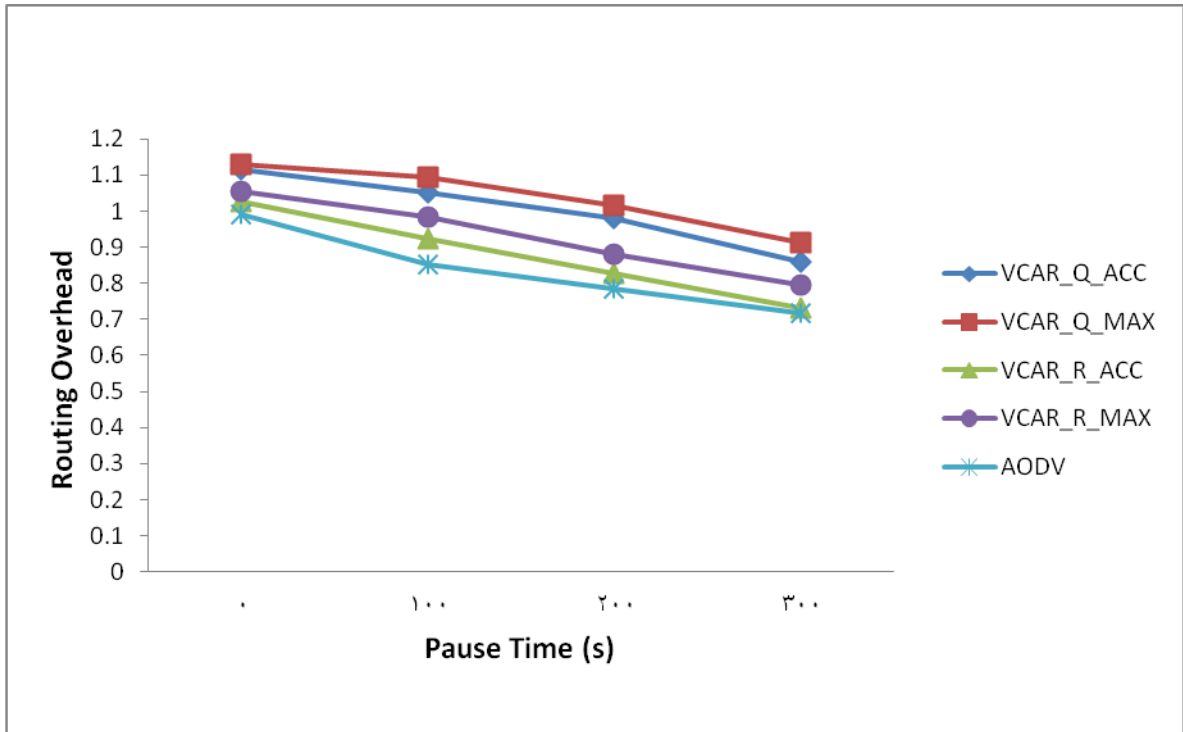


Figure 5-32: Routing Overhead of 10 sources each one sends 6 packets/s

Figure 5-33 shows the routing overhead for all protocols when the number of sources is fifteen and each source sends one packet per second. The figure shows that AODV outperforms all VCAR protocols for all pause times. When the pause time is equal to zero (high mobility network), AODV outperforms VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX by 1.60, 2.53, 4.87, and 5.29 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), AODV outperforms VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX by 0.35, 1.92, 5.25, and 7.77 percent, respectively.

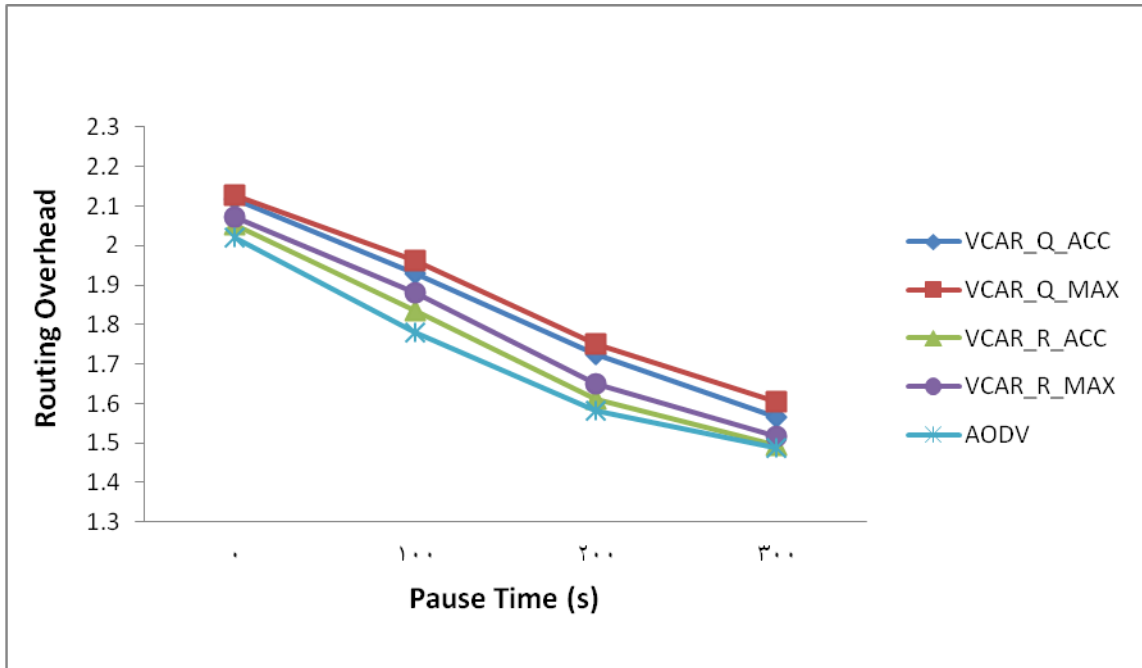


Figure 5-33: Routing Overhead of 15 sources each one sends 1 packet/s

Figure 5-34 shows the routing overhead for all protocols when the number of sources is fifteen and each source sends two packets per second. The figure shows that AODV outperforms all VCAR protocols for all pause times. When the pause time is equal to zero (high mobility network), AODV outperforms VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX by 1.63, 2.57, 4.94, and 5.37 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), AODV outperforms VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX by 0.34, 1.88, 5.14, and 7.62 percent, respectively.

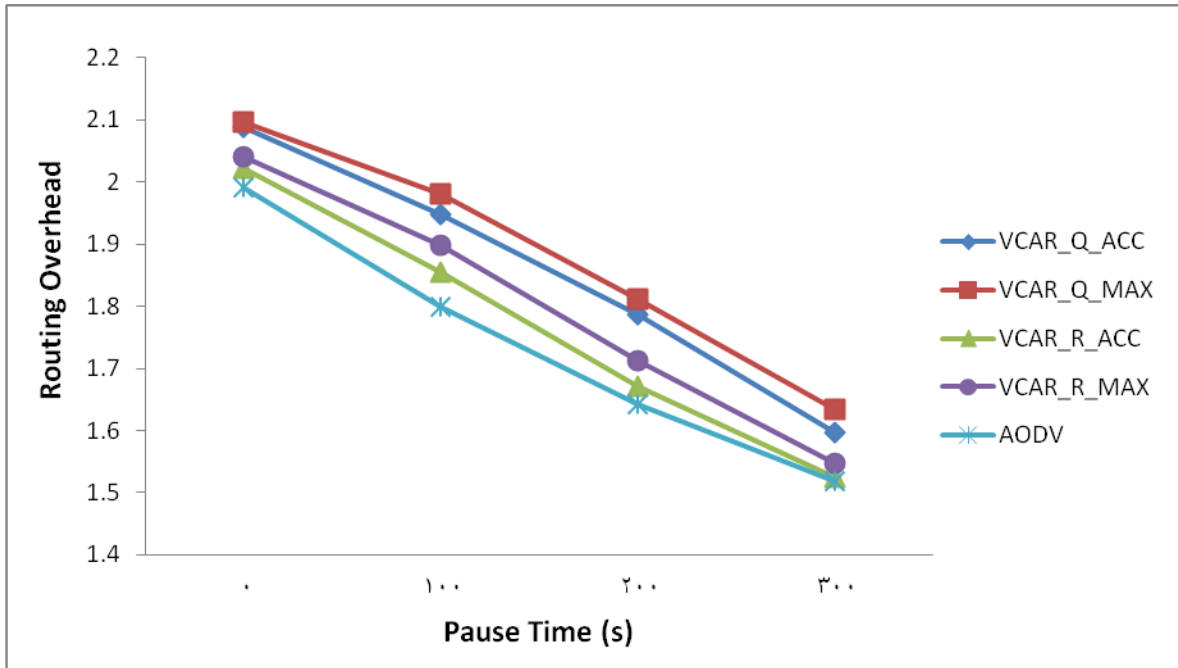


Figure 5-34: Routing Overhead of 15 sources each one sends 2 packets/s

Figure 5-35 shows the routing overhead for all protocols when the number of sources is fifteen and each source sends four packets per second. The figure shows that AODV outperforms all VCAR protocols for all pause times. When the pause time is equal to zero (high mobility network), AODV outperforms VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX by 3.4, 6.05, 12.38, and 13.66 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), AODV outperforms VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX by 1.81, 10.63, 19.80, and 26.92 percent, respectively.

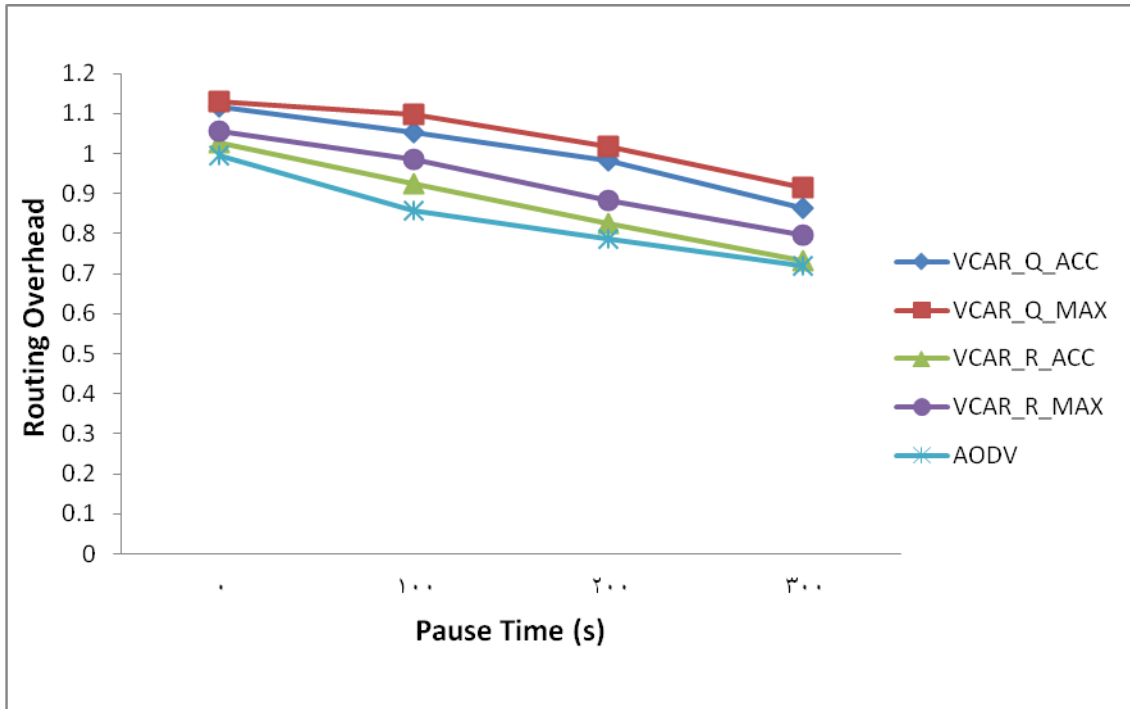


Figure 5-35: Routing Overhead of 15 sources each one sends 4 packets/s

Figure 5-36 shows the routing overhead for all protocols when the number of sources is fifteen and each source sends six packets per second. The figure shows that AODV outperforms all VCAR protocols for all pause times. When the pause time is equal to zero (high mobility network), AODV outperforms VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX by 2.16, 4.92, 11.19, and 12.60 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), AODV outperforms VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX by 4.39, 13.58, 22.99, and 30.48 percent, respectively.

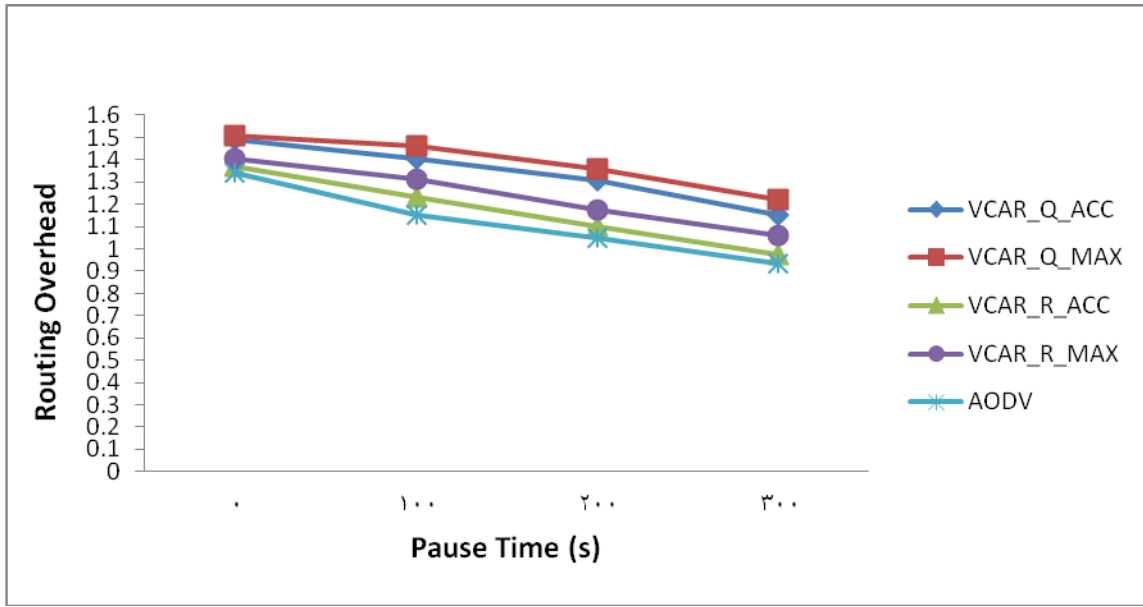


Figure 5-36: Routing Overhead of 15 sources each one sends 6 packets/s

5-4 Energy Consumption Percentage

Figures 5-37~5-48 show the energy percentage consumption for all simulation scenarios.

We can see from these figures that all VCAR protocols outperform AODV in terms of power consumption percentage for all mobility levels, all number of sources, and all transmission rates. This is because VCAR maintains the nodes' energy as much as possible by distributing the load among the intermediate nodes. It reduces the probability of the intermediate node to be exhausted, and subsequently reduces the link failures. This mechanism used by VCAR reduces the power consumption needed to do route maintenance.

Figure 5-37 shows the energy consumption percentage for all protocols when the number of sources is five and each source sends one packet per second. The figure shows that all VCAR protocols outperform AODV for all pause times. When the pause time is

equal to zero (high mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 29.37, 27.24, 21.75, and 19.06 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 30.81, 28.33, 24.07, and 18.48 percent, respectively.

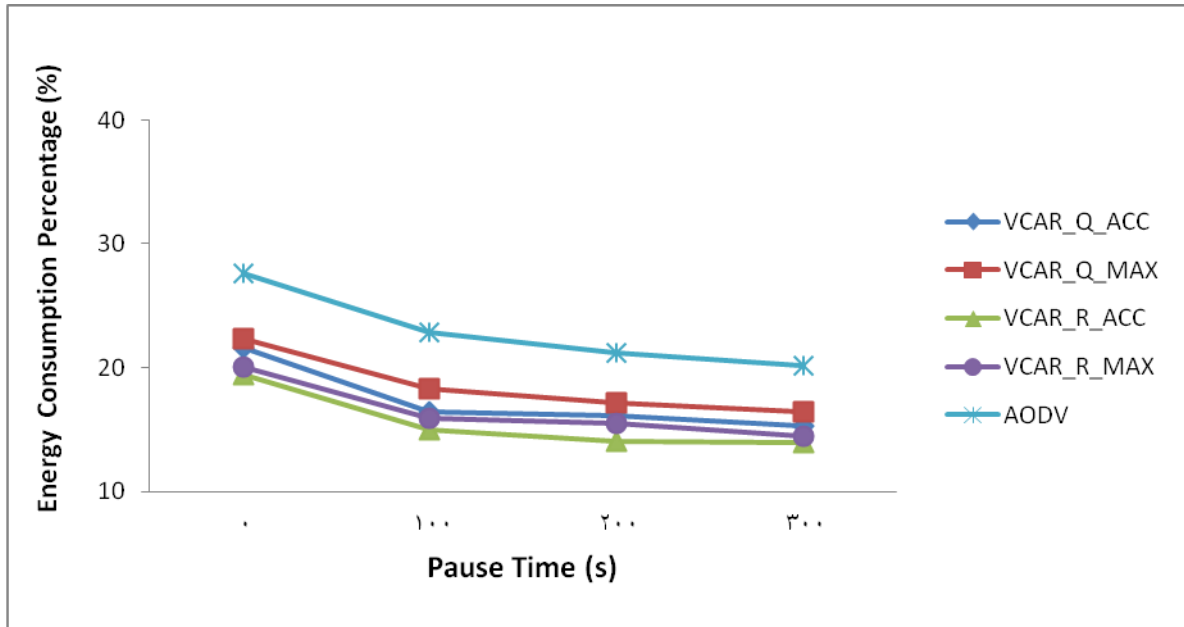


Figure 5-37: Energy Consumption Percentage of 5 sources each one sends 1 packet/s

Figure 5-38 shows the energy consumption percentage for all protocols when the number of sources is five and each source sends two packets per second. The figure shows that all VCAR protocols outperform AODV for all pause times. When the pause time is equal to zero (high mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 27.38, 25.40, 23.66, and 17.77 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 28.03, 25.78, 21.90, and 16.82 percent, respectively.

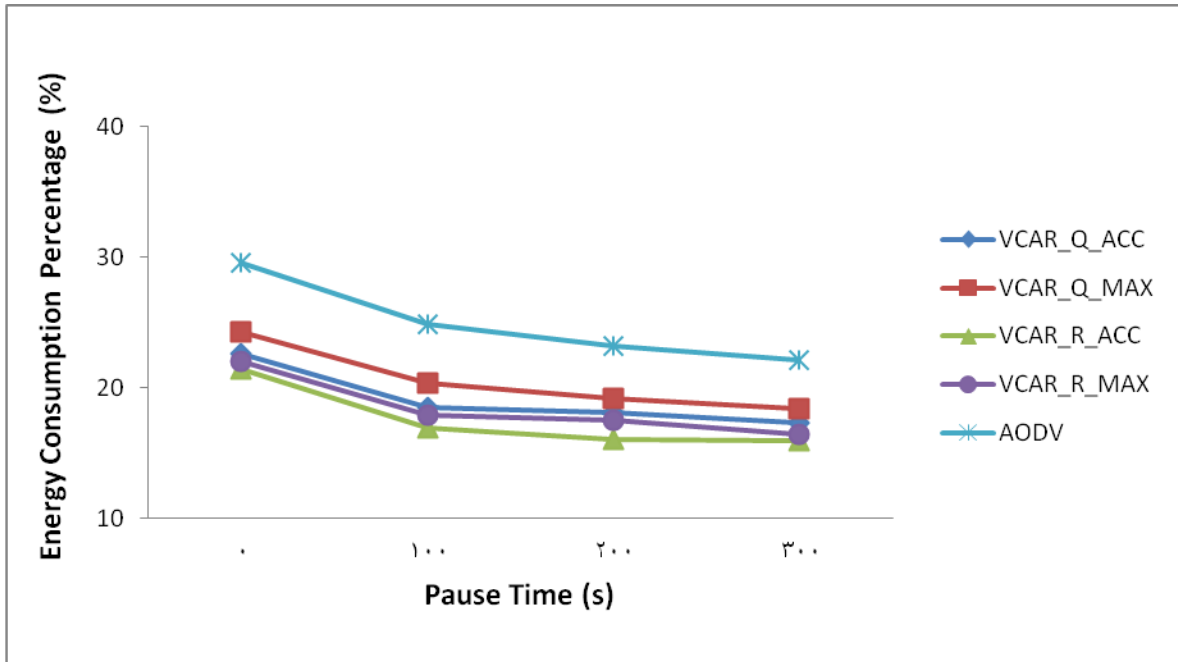


Figure 5-38: Energy Consumption Percentage of 5 sources each one sends 2 packets/s

Figure 5-39 shows the energy consumption percentage for all protocols when the number of sources is five and each source sends four packets per second. The figure shows that all VCAR protocols outperform AODV for all pause times. When the pause time is equal to zero (high mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 24.55, 22.28, 16.41, and 13.54 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 27.97, 25.40, 20.96, and 15.14 percent, respectively.

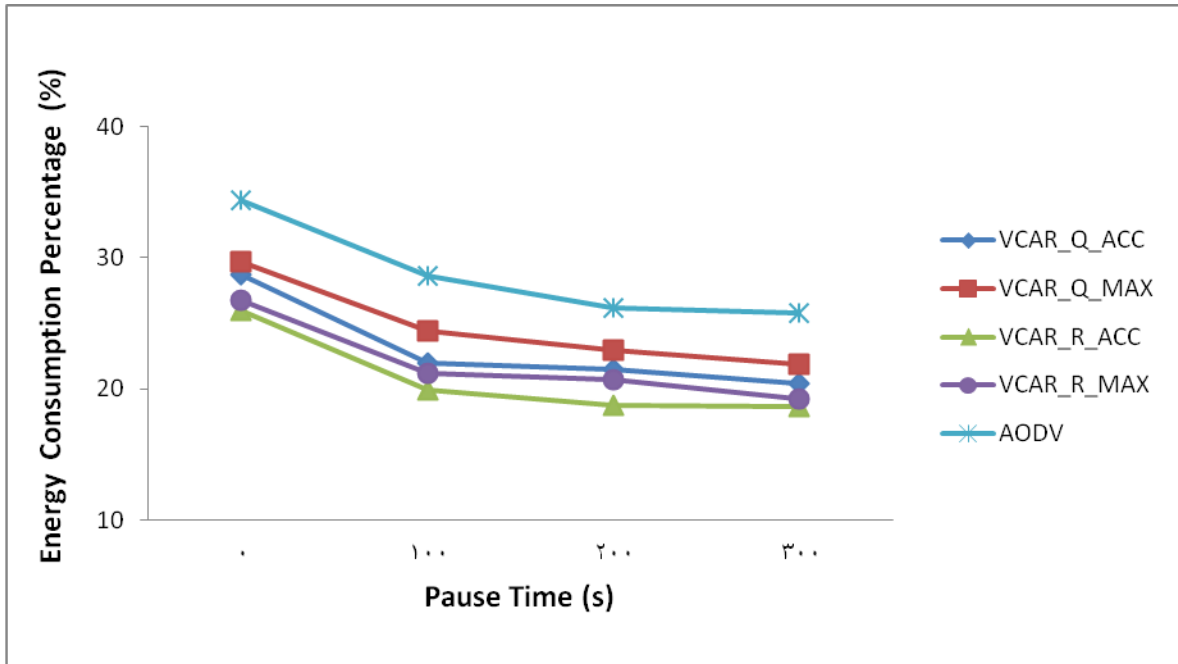


Figure 5-39: Energy Consumption Percentage of 5 sources each one sends 4 packets/s

Figure 5-40 shows the energy consumption percentage for all protocols when the number of sources is five and each source sends six packets per second. The figure shows that all VCAR protocols outperform AODV for all pause times. When the pause time is equal to zero (high mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 25.80, 23.56, 15.72, and 14.97 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 30.15, 27.64, 21.41, and 17.70 percent, respectively.

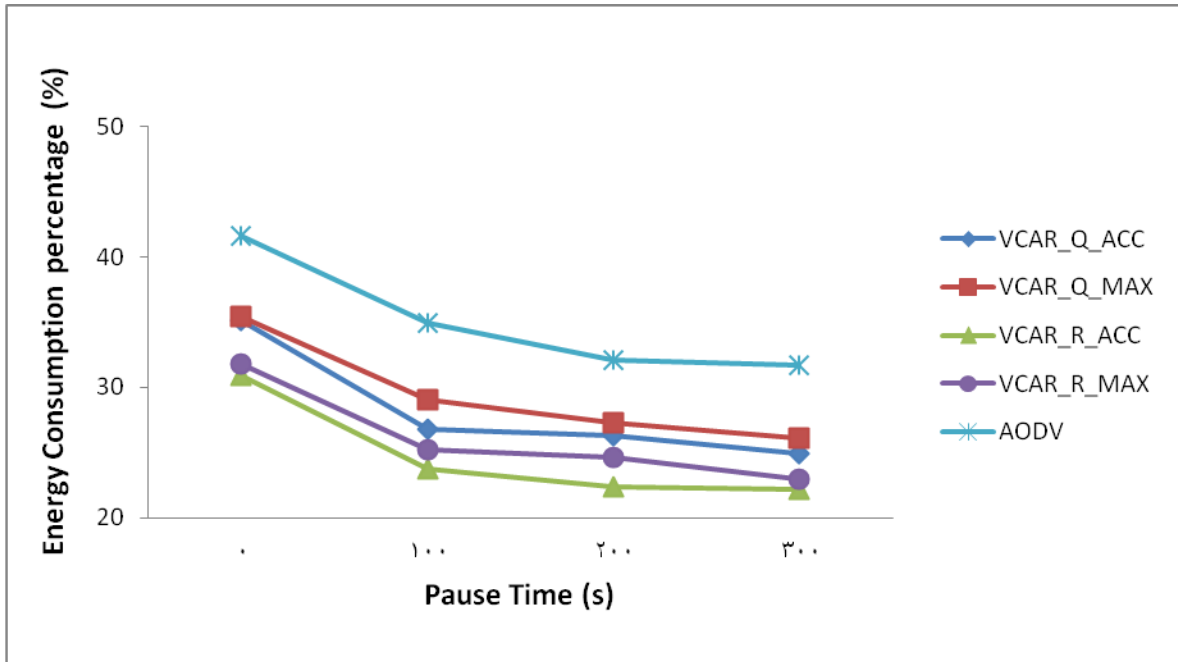


Figure 5-40: Energy Consumption Percentage of 5 sources each one sends 6 packets/s

Figure 5-41 shows the energy consumption percentage for all protocols when the number of sources is ten and each source sends one packet per second. The figure shows that all VCAR protocols outperform AODV for all pause times. When the pause time is equal to zero (high mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 27.56, 25.38, 19.75, and 16.99 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 35.25, 32.93, 28.94, and 23.71 percent, respectively.

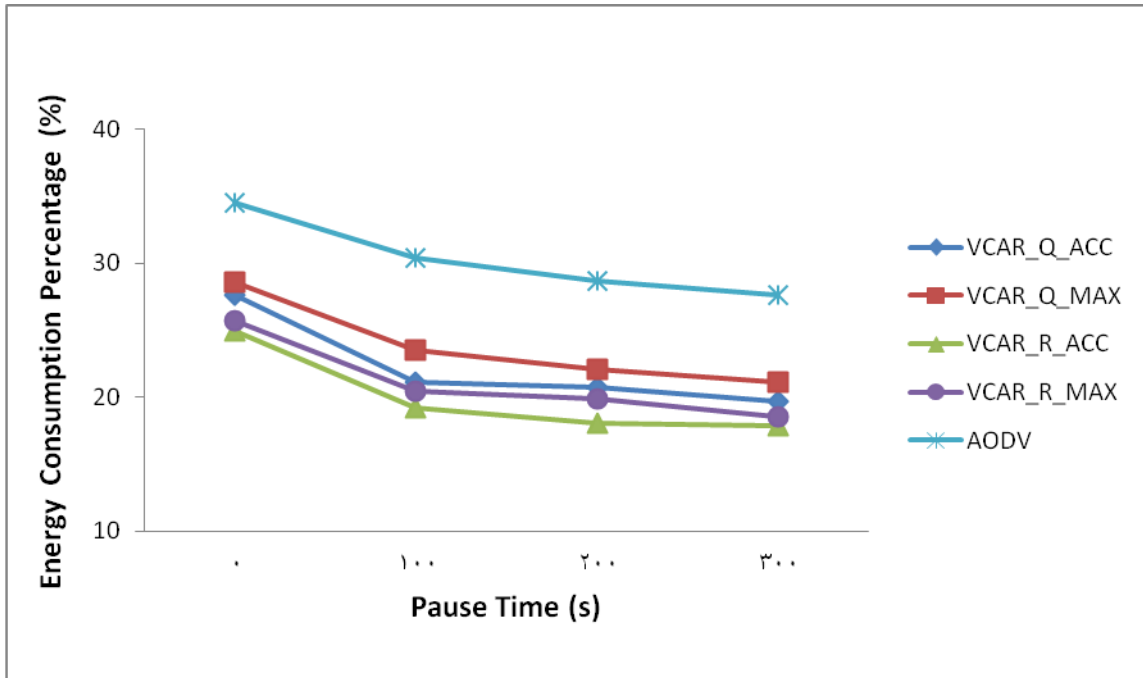


Figure 5-41: Energy Consumption Percentage of 10 sources each one sends 1 packet/s

Figure 5-42 shows the energy consumption percentage for all protocols when the number of sources is ten and each source sends two packets per second. The figure shows that all VCAR protocols outperform AODV for all pause times. When the pause time is equal to zero (high mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 24.91, 24.89, 21.50, 21.02 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 32.92, 32.53, 30.53, and 27.46 percent, respectively.

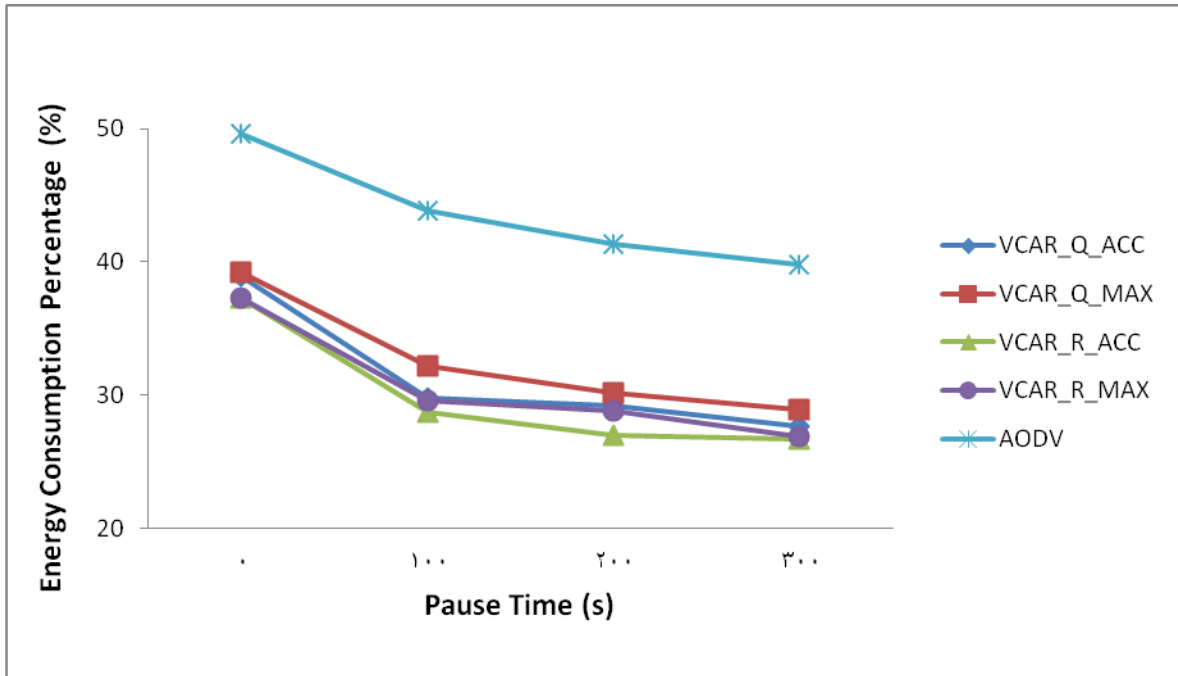


Figure 5-42:Energy Consumption Percentage of 10 sources each one sends 2 packets/s

Figure 5-43 shows the energy consumption percentage for all protocols when the number of sources is ten and each source sends four packets per second. The figure shows that all VCAR protocols outperform AODV for all pause times. When the pause time is equal to zero (high mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 16.80, 14.30, 12.20, and 10.89 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 12.18, 7.89, 5.86, and 5.12 percent, respectively.

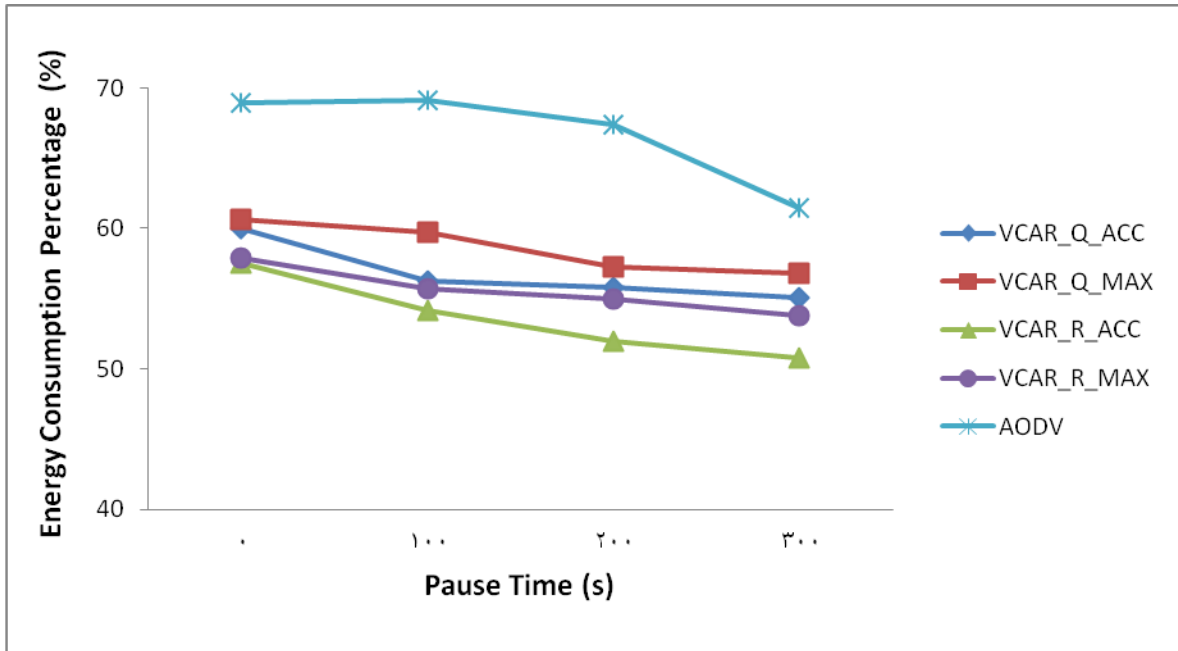


Figure 5-43:Energy Consumption Percentage of 10 sources each one sends 4 packets/s

Figure 5-44 shows the energy consumption percentage for all protocols when the number of sources is ten and each source sends six packets per second. The figure shows that all VCAR protocols outperform AODV for all pause times. When the pause time is equal to zero (high mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 16.83, 15.41, 12.38, and 10.49 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 20.52, 20.06, 18.63, and 17.13 percent, respectively.

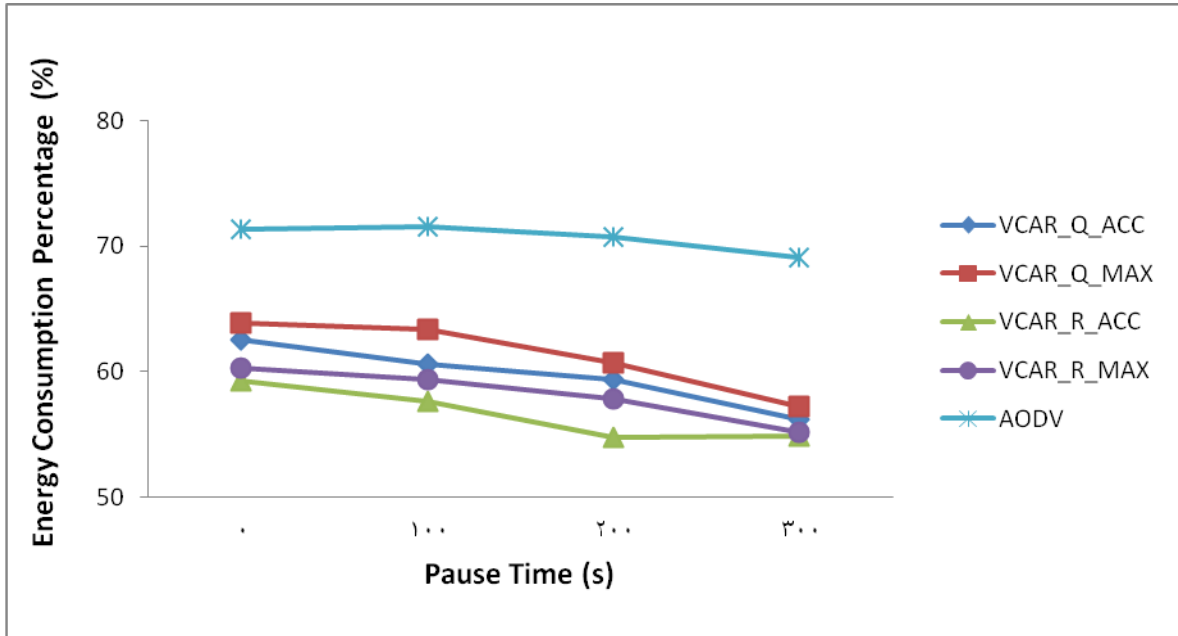


Figure 5-44:Energy Consumption Percentage of 10 sources each one sends 6 packets/s

Figure 5-45 shows the energy consumption percentage for all protocols when the number of sources is fifteen and each source sends one packet per second. The figure shows that all VCAR protocols outperform AODV for all pause times. When the pause time is equal to zero (high mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 25.16, 21.80, 21.53, and 16.22 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 20.23, 18.08, 17.19, and 13.29 percent, respectively.

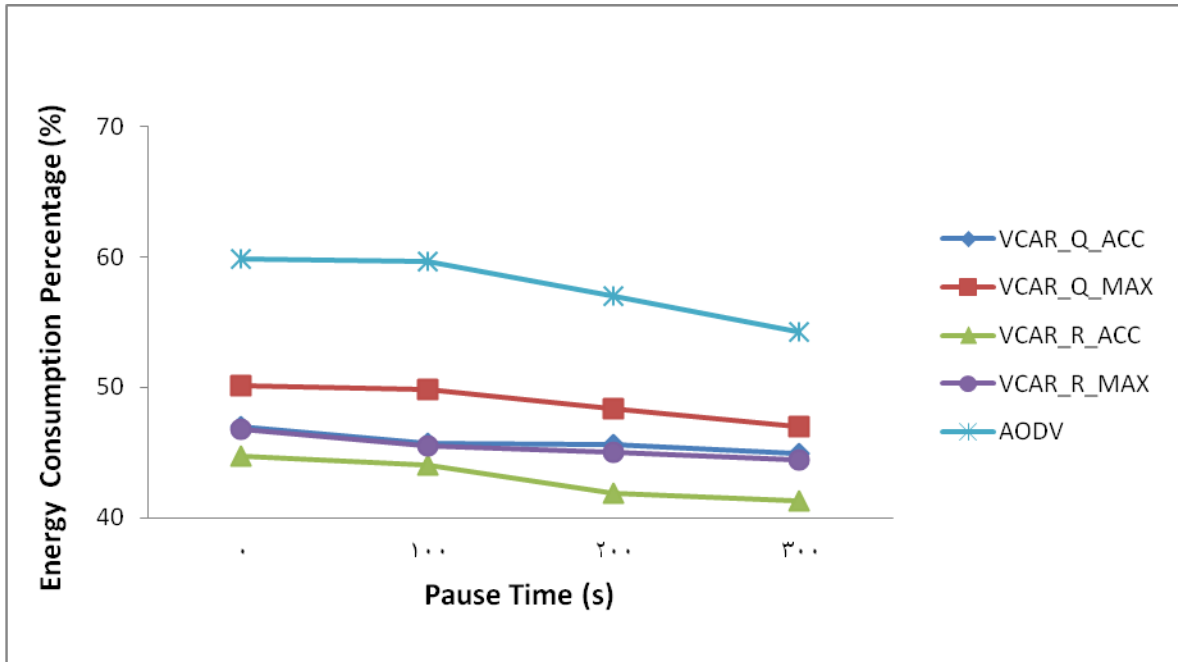


Figure 5-45: Energy Consumption Percentage of 15 sources each one sends 1 packet/s

Figure 5-46 shows the energy consumption percentage for all protocols when the number of sources is fifteen and each source sends two packets per second. The figure shows that all VCAR protocols outperform AODV for all pause times. When the pause time is equal to zero (high mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 20.98, 18.92, 17.67, and 13.46 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 20.36, 17.58, 15.62, and 12.98 percent, respectively.

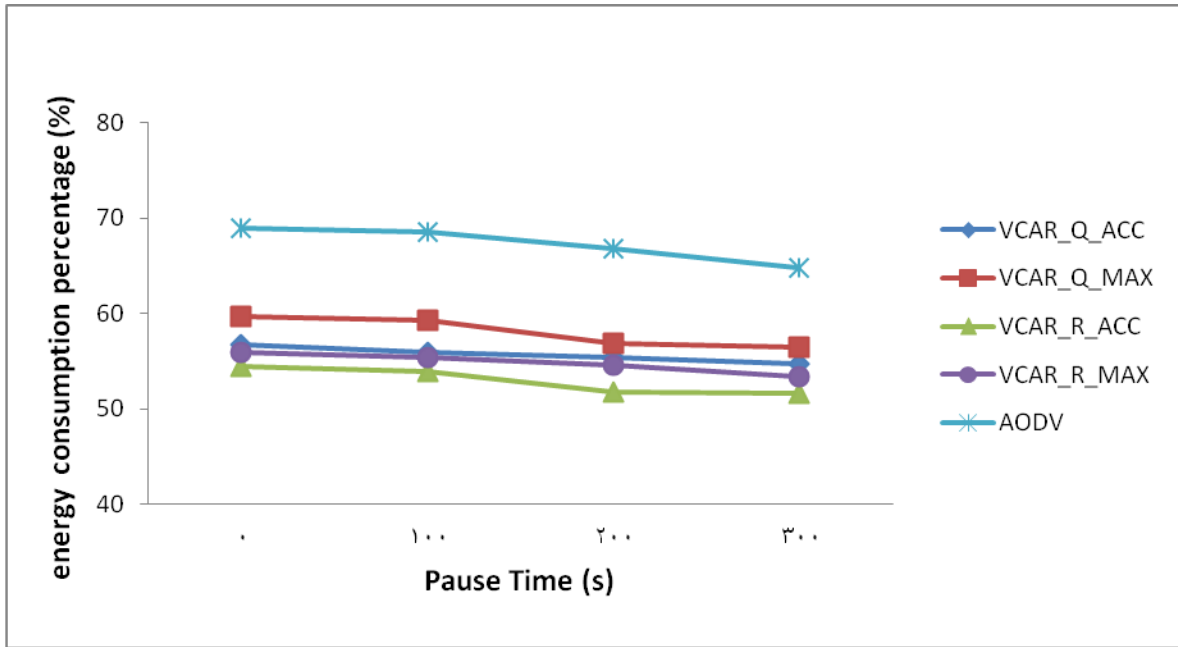


Figure 5-46:Energy Consumption Percentage of 15 sources each one sends 2 packets/s

Figure 5-47 shows the energy consumption percentage for all protocols when the number of sources is fifteen and each source sends four packets per second. The figure shows that all VCAR protocols outperform AODV for all pause times. When the pause time is equal to zero (high mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 17.19, 14.75, 14.02, and 8.70 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 20.17, 16.99, 15.50, and 13.94 percent, respectively.

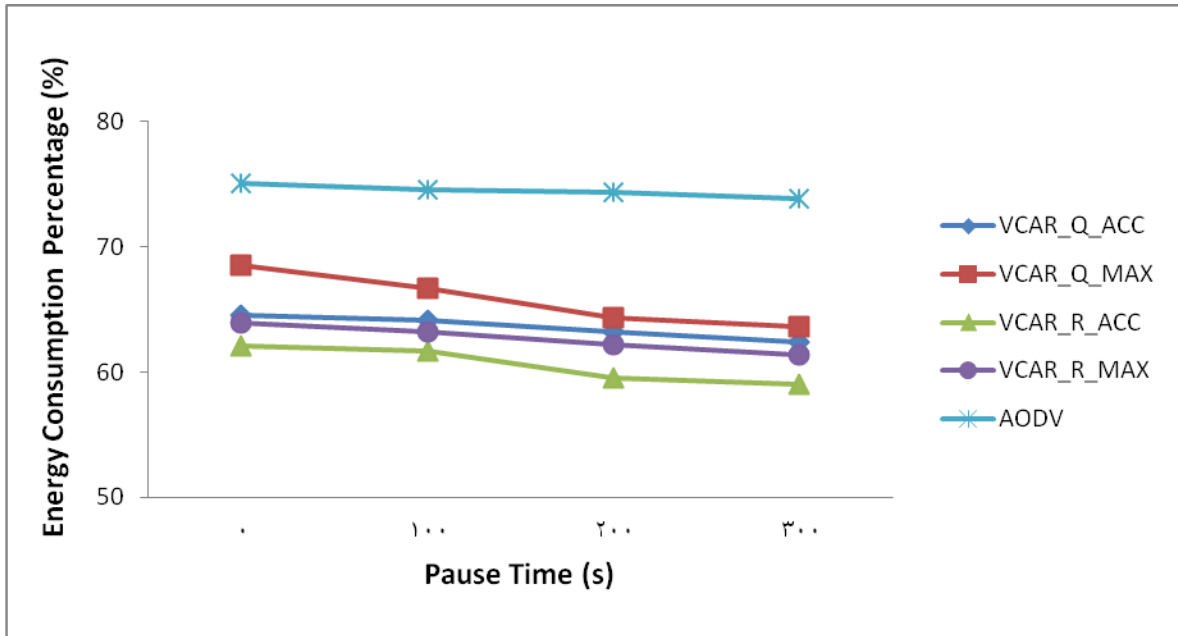


Figure 5-47:Energy Consumption Percentage of 15 sources each one sends 4 packets/s

Figure 5-48 shows the energy consumption percentage for all protocols when the number of sources is fifteen and each source sends six packets per second. The figure shows that all VCAR protocols outperform AODV for all pause times. When the pause time is equal to zero (high mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 14.94, 13.44, 11.98, and 10.23 percent, respectively. When the pause time is equal to 300 seconds (low mobility network), VCAR_R_ACC, VCAR_R_MAX, VCAR_Q_ACC, and VCAR_Q_MAX outperform AODV by 17.30, 13.76, 12.71, and 12.20 percent, respectively.

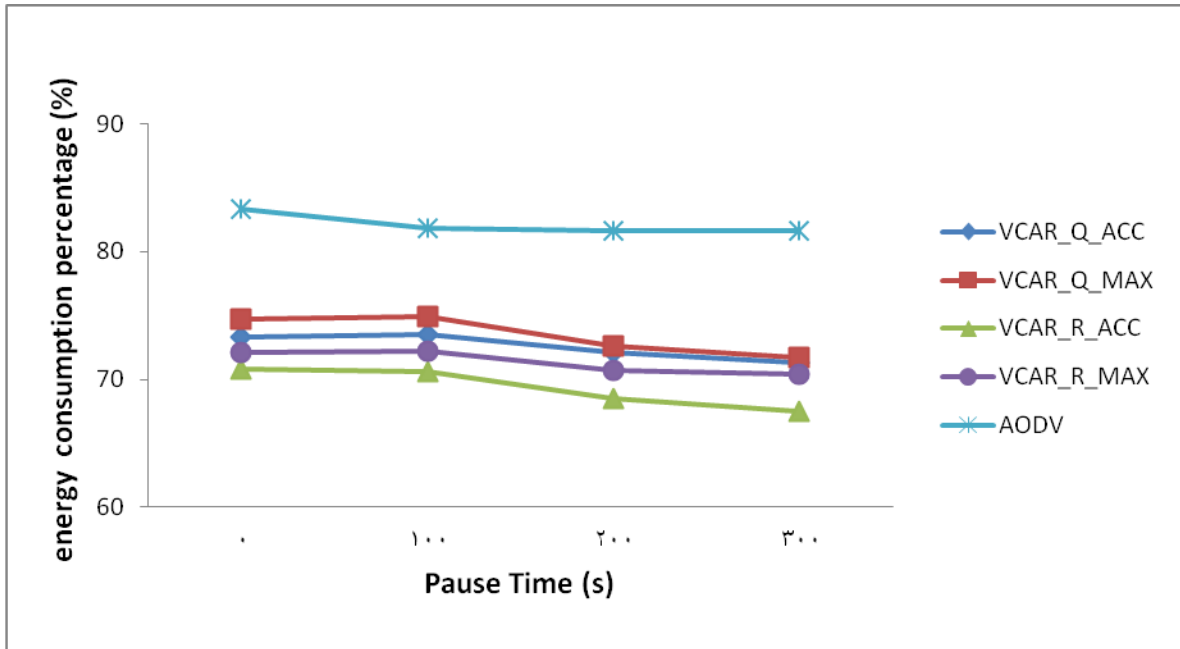


Figure 5-48:Energy Consumption Percentage of 15 sources each one sends 6 packets/s

CHAPTER SIX: CONCLUSIONS AND FUTURE WORK

6-1 Conclusions

In this study, we have implemented a new routing protocol for MANETs called VCAR, which selects the route according to both the congestion and the velocity of the intermediate nodes in the network.

VCAR was evaluated against AODV by extensive simulation using NS-2 simulation environment. The performance was evaluated using four metrics: packets delivery ratio, average end-to-end delay, routing overhead, and energy consumption percentage.

We have used different scenarios in the simulation. The nodes move according to the random-waypoint model, with pause times 0, 100, 200, 300 in each experiment. CBR traffic was generated using the transmission rate for 1, 2, 4 and 6 packets per second, repeated for 5, 10, and 15 sources.

The results collected from the simulation show that VCAR outperforms AODV by a significant value in terms of packet delivery ratio, because it selects the more stable route available from source to destination.

VCAR also outperforms AODV by a significant value in terms of average end-to-end delay, because it selects the least congested route from source to destination, so that the packets need not to wait too long in the intermediate nodes.

VCAR also outperforms AODV in terms of energy consumption percentage, because it distributes the load among intermediate nodes, and reduces the energy consumed for route maintenance.

In the other hand, the VCAR routing overhead is higher than the AODV routing overhead, this is because destination in VCAR may sends many route replies for the same route request to select the best route.

6-2 Future Work

Future work may include:

- The performance of VCAR may be enhanced by tuning the weight factor (α) dynamically according to the state of the network.
- More studies are needed in the route reply mechanism used in VCAR in order to enhance the routing overhead.

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المُلخَص

الشبكات المتحركة الخاصة هي شبكات لاسلكية تعمل بدون وجود بنية تحتية وتتكون من مجموعة من العُقد المتحركة التي تحتاج إلى بناء مسارات بين بعضها عند الحاجة إلى إرسال البيانات. تستخدم هذه العُقد خوارزميات التمرير لبناء وصيانة المسارات بين المصدر والهدف. إن معظم خوارزميات التمرير المطورة تختار مسارات التمرير دون أن تأخذ بعين الاعتبار كلاً من الازدحام في العُقد الوسطية وحركة هذه العُقد. بعض خوارزميات التمرير التي تم تطويرها تأخذ بعين الاعتبار الازدحام فقط في اختيار المسارات، مما يؤدي إلى اختيار مسارات غير مستقرة مما يزيد من احتمال حدوث أعطال متكررة في المسارات وبالتالي حدوث فقدان متكرر للرزق المرسل. في حين تعمل خوارزميات أخرى على اختيار المسارات اعتماداً على سرعة العقد الوسطية مما يؤدي إلى اختيار مسارات مزدحمة، وبالتالي حدوث بطء في وصول الرزم من المصدر إلى الهدف.

في هذا البحث تم اقتراح خوارزمية تمرير اعتماداً على السرعة والازدحام في العُقد التي يتكون منها المسار. وتقوم هذه الخوارزمية باختيار المسار الأفضل باستخدام معيارين هما مستوى الازدحام في العُقد الوسطية وسرعة حركة هذه العُقد.

لقياس مستوى الازدحام في العُقد الوسطية تستخدم الخوارزمية المقترحة طريقتين. الطريقة الأولى تتم باستخدام عدد الرزم التي تنتظر في طابور الانتظار في العُقد والطرقة الثانية تتم باستخدام عدد المسارات التي تمر من خلال العُقد.

تعمل الخوارزمية المقترحة على اختيار المسار الأفضل بأسلوبين. الأسلوب الأول يختار المسار الأفضل بناءً على القيمة التراكمية لجميع العُقد الوسطية والأسلوب الثاني يختار المسار الأفضل بناءً على قيمة النهاية العظمى من بين جميع العُقد الوسطية.

تم تنفيذ أربع إصدارات من الخوارزمية المقترحة. الإصدار الأول يعتبر المسار الأفضل هو المسار ذو القيمة التراكمية الأقل من حيث عدد الرزم التي تنتظر في طابور الانتظار وسرعة حركة العُقد الوسطية. الإصدار الثاني يعتبر المسار الأفضل هو المسار ذو القيمة العظمى الأقل من حيث عدد الرزم التي تنتظر في طابور الانتظار وسرعة حركة العُقد الوسطية. الإصدار الثالث يعتبر المسار الأفضل هو المسار ذو القيمة التراكمية الأقل من حيث عدد المسارات التي

تمر من خلال العُقد الوسطية وسرعة حركة العُقد الوسطية. أما الإصدار الرابع فيعتبر المسار الأفضل هو المسار ذو القيمة العظمى الأقل من حيث عدد المسارات التي تمر من خلال العُقد الوسطية وسرعة حركة العُقد الوسطية.

تم إجراء محاكاة مكثفة للخوارزميات المقترحة باستخدام محاكي الشبكة (NS2) وتم استخراج نتائج الخوارزميات المقترحة ومقارنتها مع أداء خوارزمية (AODV) وأظهرت نتائج المحاكاة تحسينات ملحوظة بالنسبة لنسبة تسليم الرزم ومعدل زمن التأخير ونسبة استهلاك الطاقة. ولكن النتائج أظهرت أيضا أن أداء الخوارزميات المقترحة كان أسوأ بالنسبة لمعيار كلفة التمرير.

APPENDIX A: RESULTS WHEN $\alpha=0.25$

Table A-1: Packet Delivery Ratio when $\alpha=0.25$

Five sources each one sends one packet/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	86.36	85	92.49	91.25
	100	94.38	93.61	98.13	97.04
	200	96.2	95.74	99.12	98.23
	300	97.24	96.03	100	100
Five sources each one sends two packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	89.64	88.73	94.51	93.05
	100	94.5	94.11	97.24	96.84
	200	95.67	95.21	97.91	98.01
	300	96.99	96.47	100	100
Five sources each one sends four packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	90.33	89.91	91.36	90.91
	100	94.22	93.42	95.84	94.35
	200	93.19	92.77	97.41	95.84
	300	95.48	95.19	98.13	97.55
Five sources each one sends six packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	88.72	88.31	89.35	89.01
	100	90.52	89.86	91.14	90.69
	200	92.39	91.68	94.52	93.41
	300	95.54	95.09	97.77	96.25
Ten sources each one sends one packet/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	89.05	88.74	90.28	89.91
	100	92.11	91.58	93.33	92.76
	200	94.54	93.83	96.36	95.81
	300	95.51	94.47	97.7	96.59
Ten sources each one sends two packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	89.34	89.05	90.98	90.24
	100	92.87	92.33	93.91	93.25
	200	94.35	93.87	96.65	95.42
	300	95.48	94.13	97.05	96.48

Ten sources each one sends four packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	90.95	90.27	91.54	91.05
	100	81.84	81.29	82.94	82.15
	200	93.88	93.62	95.98	95.64
	300	95.62	94.41	97.23	96.87
Ten sources each one sends six packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	83.62	83.54	84.16	84.08
	100	67.62	64.84	69.25	68.11
	200	72.12	71.91	73.84	72.65
	300	89.91	89.17	91.52	90.51
Fifteen sources each one sends one packet/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	90.92	90.55	91.55	91.45
	100	94.13	93.98	95.88	95.14
	200	95.82	95.15	97.16	96.74
	300	98.14	98.04	98.91	98.85
Fifteen sources each one sends two packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	89.49	88.66	91.02	90.54
	100	91.51	90.16	93.67	92.81
	200	93.34	92.05	95.23	94.11
	300	94.43	94.12	96.91	95.74
Fifteen sources each one sends four packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	78.56	77.22	80.65	79.12
	100	68.12	66.94	70.05	69.51
	200	74.35	74.05	76.12	75.3
	300	81.95	81.55	82.61	82.05
Fifteen sources each one sends six packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	61.51	61.08	63.02	62.84
	100	50.13	49.85	51.81	50.44
	200	55.64	54.71	57.14	56.41
	300	64.02	63.77	65.31	64.51

Table A-2: Average End-to-End Delay when $\alpha=0.25$

Five sources each one sends one packet/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	0.4178	0.4419	0.3244	0.3502
	100	0.3114	0.3952	0.2418	0.2711
	200	0.0667	0.0702	0.0502	0.0592
	300	0.0228	0.0294	0.0134	0.0203
Five sources each one sends two packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	0.1728	0.1763	0.1662	0.1705
	100	0.1864	0.1926	0.1601	0.1788
	200	0.085	0.0912	0.0642	0.0778
	300	0.0239	0.0252	0.0197	0.0215
Five sources each one sends four packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	0.3601	0.3696	0.3455	0.3549
	100	0.1744	0.1752	0.1502	0.1629
	200	0.0602	0.0667	0.0512	0.0561
	300	0.0274	0.0281	0.0212	0.0229
Five sources each one sends six packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	0.2025	0.2305	0.1692	0.1837
	100	0.2448	0.249	0.2194	0.2208
	200	0.0639	0.0651	0.0606	0.0624
	300	0.0274	0.0281	0.0211	0.0262
Ten sources each one sends one packet/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	0.1779	0.1806	0.1662	0.1724
	100	0.2405	0.2487	0.2344	0.2391
	200	0.0602	0.0619	0.0584	0.0591
	300	0.0306	0.0311	0.0282	0.0294
Ten sources each one sends two packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	0.0971	0.0986	0.0914	0.0952
	100	0.2091	0.2112	0.1968	0.2007
	200	0.0514	0.0526	0.0488	0.0509
	300	0.0423	0.0435	0.0394	0.041

Ten sources each one sends four packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	0.1369	0.1394	0.1305	0.1334
	100	0.3611	0.3678	0.3507	0.3594
	200	0.0904	0.0913	0.0891	0.0899
	300	0.1368	0.1397	0.1256	0.1311
Ten sources each one sends six packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	0.1994	0.2016	0.1928	0.1957
	100	0.7711	0.7809	0.7501	0.7607
	200	0.4601	0.4715	0.4109	0.4315
	300	0.2218	0.2307	0.2005	0.2109
Fifteen sources each one sends one packet/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	0.1102	0.1118	0.1048	0.1094
	100	0.1347	0.1376	0.1314	0.1335
	200	0.0552	0.0571	0.0512	0.0534
	300	0.0468	0.0491	0.0448	0.0455
Fifteen sources each one sends two packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	0.0922	0.0934	0.0905	0.0912
	100	0.1638	0.1645	0.1611	0.1627
	200	0.0579	0.0582	0.0554	0.0568
	300	0.0519	0.0526	0.0492	0.0504
Fifteen sources each one sends four packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	0.2335	0.2391	0.2294	0.2314
	100	0.6208	0.6233	0.6159	0.6165
	200	0.5218	0.5266	0.5088	0.5142
	300	0.4508	0.4552	0.4415	0.4463
Fifteen sources each one sends six packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	0.4483	0.4491	0.4426	0.4468
	100	1.1308	1.1416	1.1056	1.1125
	200	1.2271	1.2391	1.2041	1.2151
	300	1.1294	1.1402	1.1042	1.1119

Table A-3: Routing Overhead when $\alpha=0.25$

Five sources each one sends one packet/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	2.95	3.08	2.64	2.77
	100	2.47	2.84	2.35	2.51
	200	2.42	2.66	2.04	2.29
	300	1.88	2.09	1.67	1.51
Five sources each one sends two packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	1.95	1.99	1.84	1.91
	100	1.61	1.69	1.47	1.55
	200	1.42	1.46	1.31	1.35
	300	1.34	1.42	1.15	1.27
Five sources each one sends four packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	4.12	4.33	3.42	3.91
	100	2.93	2.98	2.74	2.89
	200	3.04	3.19	2.56	2.87
	300	2.19	2.25	2.07	2.12
Five sources each one sends six packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	3.56	3.67	3.33	3.41
	100	2.73	2.91	2.51	2.69
	200	2.09	2.24	1.74	1.9
	300	1.85	1.93	1.41	1.62
Ten sources each one sends one packet/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	3.81	3.88	3.66	3.71
	100	3.15	3.31	2.97	3.04
	200	2.5	2.68	2.35	2.41
	300	2.13	2.19	1.91	2.05
Ten sources each one sends two packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	2.89	2.94	2.77	2.84
	100	2.81	2.97	2.68	2.74
	200	1.99	2.01	1.94	1.97
	300	1.83	1.94	1.58	1.74
Ten sources each one	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX

sends four packets/second	0	2.88	2.93	2.75	2.81
	100	2.76	2.84	2.54	2.63
	200	2.04	2.15	1.91	1.99
	300	1.84	1.95	1.62	1.77
Ten sources each one sends six packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	2.69	2.74	2.48	2.56
	100	2.25	2.34	2.04	2.15
	200	1.95	2.04	1.7	1.85
300	1.58	1.64	1.31	1.49	
Fifteen sources each one sends one packet/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	4.05	4.12	3.82	3.91
	100	3.71	3.92	3.52	3.64
	200	3.27	3.35	3.09	3.16
300	3.14	3.21	2.99	3.07	
Fifteen sources each one sends two packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	3.05	3.13	2.81	2.92
	100	2.63	2.66	2.54	2.59
	200	2.06	2.14	1.92	1.99
300	1.88	1.95	1.52	1.67	
Fifteen sources each one sends four packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	2.58	2.64	2.33	2.41
	100	2.11	2.25	1.85	1.99
	200	1.56	1.69	1.35	1.47
300	1.62	1.84	1.15	1.24	
Fifteen sources each one sends six packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	1.86	1.94	1.55	1.67
	100	1.57	1.6	1.34	1.42
	200	1.36	1.49	1.15	1.24
300	1.56	1.62	1.31	1.47	

Table A-4: Energy Consumption Percentage when $\alpha=0.25$

Five sources each one sends one packet/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	22.6214	23.0527	22.0112	22.2511
	100	24.5188	25.5611	22.1924	22.145
	200	20.1025	21.281	19.7225	19.9499
	300	19.9447	20.9049	18.3255	19.1496
Five sources each one sends two packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	24.5199	25.1096	23.4136	23.9048
	100	27.0051	29.3681	25.441	26.7513
	200	24.6625	25.2276	23.9354	24.8152
	300	23.4082	24.9157	22.1289	24.2099
Five sources each one sends four packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	44.6914	46.0105	39.7124	41.5823
	100	47.3815	49.9127	42.1189	45.1806
	200	38.6944	38.107	36.2881	37.4166
	300	35.6488	36.1492	34.1264	34.9005
Five sources each one sends six packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	57.3631	59.0058	53.1684	55.9871
	100	62.6712	64.1823	57.9118	60.0048
	200	51.3514	51.9803	50.1182	50.9488
	300	49.1825	49.8522	44.9258	47.3522
Ten sources each one sends one packet/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	31.0558	31.9024	28.1425	30.4418
	100	35.1863	36.2088	33.1204	34.0255
	200	31.3947	32.1558	28.3035	30.9044
	300	29.9076	30.1184	26.4712	29.115
Ten sources each one sends two packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	42.0778	43.1725	39.2114	40.8256
	100	48.1185	49.4618	46.2914	47.5182
	200	42.0582	42.8273	40.9112	41.5722
	300	36.7582	37.7522	36.312	37.8436
Ten sources each one	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX

sends four packets/second	0	61.6184	62.8724	59.2841	60.357
	100	76.4429	76.9831	73.0059	74.5896
	200	68.0593	69.9211	65.1479	66.6971
	300	60.3641	61.2776	58.0022	59.4682
Ten sources each one sends six packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	82.5978	83.0478	79.2584	80.3647
	100	84.1814	84.7691	83.4935	83.978
	200	84.9823	85.306	82.2831	83.9921
	300	73.9498	74.7055	72.1628	72.8671
Fifteen sources each one sends one packet/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	35.7204	36.5478	33.5198	34.2487
	100	38.5611	38.9155	36.8221	37.0489
	200	38.5697	39.0047	33.4189	37.2941
	300	32.3778	34.8256	30.3051	31.8452
Fifteen sources each one sends two packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	51.9544	52.0047	50.1584	51.3641
	100	57.3814	58.6284	55.8513	56.1482
	200	52.5825	53.3716	50.8521	51.4497
	300	50.6721	51.0852	48.3252	49.0026
Fifteen sources each one sends four packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	82.2792	83.1047	80.4459	81.367
	100	83.6472	84.4426	81.5712	82.9147
	200	82.9258	83.5528	80.2591	81.0159
	300	76.508	77.1184	74.6462	75.4738
Fifteen sources each one sends six packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	89.0059	89.1578	88.0157	88.6347
	100	88.9528	89.4426	87.9501	88.1056
	200	86.9561	87.2914	85.9411	86.047
	300	81.1834	81.7623	80.0451	80.6834

APPENDIX B: RESULTS WHEN $\alpha=0.75$

Table B-1: Packet Delivery Ratio when $\alpha=0.75$

Five sources each one sends one packet/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	90.32	89.98	92.25	91.08
	100	97.02	96.81	98.13	97.52
	200	98.05	97.91	99.25	98.94
	300	100	100	100	100
Five sources each one sends two packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	92.15	91.95	94.25	93.64
	100	96.95	96.66	97.85	97.05
	200	98.06	97.82	99.61	98.82
	300	100	100	100	100
Five sources each one sends four packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	93.25	92.94	95.82	95.12
	100	97.52	96.91	98.85	98.63
	200	98.15	97.56	98.99	98.62
	300	100	100	100	100
Five sources each one sends six packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	94.13	93.68	95.44	94.95
	100	95.84	95.11	96.91	96.06
	200	97.82	97.75	98.95	98.15
	300	100	100	100	100
Ten sources each one sends one packet/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	93.81	93.11	94.95	94.05
	100	96.11	95.96	97.62	96.98
	200	97.51	97.32	98.85	98.05
	300	100	100	100	100
Ten sources each one sends two packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	93.12	92.84	94.63	93.91
	100	95.85	95.11	97.41	96.92
	200	98.95	98.81	99.95	99.12
	300	100	100	100	100

Ten sources each one sends four packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	93.54	92.98	96.24	94.12
	100	85.36	84.81	87.2	86.15
	200	97.81	97.42	98.56	98.05
	300	99.21	99.09	99.56	99.44
Ten sources each one sends six packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	89.54	88.09	91.62	90.15
	100	73.88	72.91	75.92	74.25
	200	75.55	74.98	77.63	76.53
	300	94.06	93.53	96.42	95.15
Fifteen sources each one sends one packet/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	94.14	93.62	97.41	96.62
	100	97.42	96.98	98.04	97.62
	200	98.75	98.65	99.52	99.11
	300	100	100	100	100
Fifteen sources each one sends two packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	94.51	94.66	96.47	95.74
	100	95.99	95.71	97.88	96.95
	200	98.42	98.05	98.95	98.66
	300	100	100	100	100
Fifteen sources each one sends four packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	84.35	83.33	86.35	84.62
	100	73.15	72.14	75.82	74.09
	200	81.26	80.22	82.19	81.85
	300	84.16	83.55	87.15	85.84
Fifteen sources each one sends six packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	67.15	66.6	70.26	68.25
	100	55.09	54.27	57.41	56.33
	200	61.72	60.79	63.19	62.94
	300	68.52	67.53	70.14	69.52

Table B-2: Average End-to-End Delay when $\alpha=0.75$

Five sources each one sends one packet/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	0.3101	0.3209	0.3012	0.3094
	100	0.2391	0.2451	0.2214	0.2359
	200	0.0523	0.0542	0.0499	0.0511
	300	0.014	0.0142	0.0133	0.0138
Five sources each one sends two packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	0.1423	0.1475	0.1341	0.1391
	100	0.1392	0.1412	0.1306	0.1371
	200	0.0522	0.0541	0.0491	0.0509
	300	0.0146	0.015	0.0133	0.0142
Five sources each one sends four packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	0.2911	0.3056	0.2715	0.2803
	100	0.0915	0.0923	0.0808	0.0866
	200	0.0352	0.0364	0.0322	0.0334
	300	0.0149	0.0153	0.0134	0.0145
Five sources each one sends six packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	0.1405	0.1423	0.1322	0.1362
	100	0.2011	0.2025	0.1925	0.1995
	200	0.0536	0.0551	0.0512	0.0524
	300	0.0155	0.0159	0.0149	0.0151
Ten sources each one sends one packet/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	0.1472	0.151	0.1391	0.1421
	100	0.2107	0.2116	0.2022	0.2094
	200	0.0523	0.0544	0.0492	0.0509
	300	0.0239	0.0245	0.0221	0.0235
Ten sources each one sends two packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	0.0795	0.0806	0.0766	0.0782
	100	0.1823	0.1882	0.1711	0.1762
	200	0.0425	0.0429	0.0392	0.0411
	300	0.0345	0.0351	0.0325	0.0335

Ten sources each one sends four packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	0.1205	0.1255	0.1145	0.1192
	100	0.2715	0.2805	0.262	0.2688
	200	0.0811	0.0852	0.0772	0.0798
	300	0.1342	0.1367	0.1255	0.1312
Ten sources each one sends six packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	0.1719	0.1786	0.1625	0.1675
	100	0.641	0.6558	0.605	0.6229
	200	0.4493	0.4527	0.4262	0.4381
	300	0.2315	0.2388	0.2155	0.2205
Fifteen sources each one sends one packet/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	0.0942	0.0962	0.0885	0.0911
	100	0.114	0.118	0.107	0.111
	200	0.0461	0.0473	0.0435	0.0452
	300	0.0321	0.0341	0.0305	0.0315
Fifteen sources each one sends two packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	0.0761	0.0782	0.0741	0.0752
	100	0.1512	0.1523	0.1442	0.1492
	200	0.0459	0.0471	0.0431	0.0445
	300	0.0425	0.0443	0.0394	0.0412
Fifteen sources each one sends four packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	0.2113	0.2151	0.2012	0.2071
	100	0.5617	0.572	0.5514	0.5534
	200	0.4512	0.4579	0.4382	0.4476
	300	0.4052	0.4192	0.3801	0.3941
Fifteen sources each one sends six packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	0.4182	0.42005	0.3951	0.4021
	100	0.9971	1.005	0.9541	0.9735
	200	1.0995	1.1009	1.0251	1.0741
	300	0.9811	0.9905	0.9422	0.9624

Table B-3: Routing Overhead when $\alpha=0.75$

Five sources each one sends one packet/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	2.86	2.99	2.45	2.79
	100	2.16	2.29	1.98	2.05
	200	2.09	2.15	1.99	2.03
	300	1.69	1.9	1.47	1.56
Five sources each one sends two packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	1.95	1.99	1.85	1.9
	100	1.56	1.6	1.47	1.5
	200	1.39	1.45	1.27	1.31
	300	1.26	1.3	1.09	1.19
Five sources each one sends four packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	3.37	3.42	3.29	3.31
	100	2.86	2.9	2.71	2.79
	200	2.7	2.74	2.49	2.61
	300	2.27	2.36	1.99	2.09
Five sources each one sends six packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	3.36	3.49	3.08	3.19
	100	2.57	2.6	2.39	2.49
	200	1.79	1.83	1.69	1.72
	300	1.42	1.49	1.32	1.36
Ten sources each one sends one packet/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	3.8	3.96	3.57	3.69
	100	3.17	3.3	2.99	3.05
	200	2.48	2.56	2.29	2.39
	300	1.96	2.01	1.79	1.86
Ten sources each one sends two packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	2.79	2.8	2.69	2.75
	100	2.68	2.71	2.56	2.63
	200	1.96	2.04	1.86	1.91
	300	1.8	1.85	1.6	1.67

Ten sources each one sends four packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	2.75	2.79	2.68	2.71
	100	2.59	2.63	2.51	2.56
	200	1.98	2.02	1.88	1.93
	300	1.7	1.74	1.6	1.64
Ten sources each one sends six packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	2.61	2.67	2.37	2.52
	100	2.28	2.33	2.03	2.09
	200	1.58	1.6	1.46	1.51
	300	1.37	1.41	1.25	1.31
Fifteen sources each one sends one packet/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	3.96	4.02	3.69	3.84
	100	3.59	3.68	3.39	3.51
	200	3.19	3.31	3.05	3.09
	300	3.02	3.11	2.89	2.94
Fifteen sources each one sends two packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	2.88	2.91	2.77	2.81
	100	2.58	2.6	2.44	2.53
	200	1.9	1.95	1.79	1.84
	300	1.49	1.52	1.42	1.47
Fifteen sources each one sends four packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	2.37	2.39	2.28	2.31
	100	1.92	1.95	1.81	1.86
	200	1.39	1.42	1.31	1.37
	300	1.2	1.25	1.09	1.14
Fifteen sources each one sends six packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	1.56	1.65	1.47	1.51
	100	1.34	1.41	1.26	1.29
	200	1.23	1.26	1.1	1.19
	300	1.2	1.24	1.09	1.16

Table B-4: Energy Consumption Percentage when $\alpha=0.75$

Five sources each one sends one packet/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	20.8423	21.0062	19.2571	20.2051
	100	20.9145	21.6402	19.5247	20.5061
	200	18.8531	19.2012	17.3145	18.251
	300	17.2105	17.7533	16.1682	16.9875
Five sources each one sends two packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	21.1358	21.9251	20.5815	20.9487
	100	22.8257	23.1594	21.9825	22.1056
	200	20.4253	21.5591	19.1982	19.9823
	300	20.4235	20.9915	18.5249	19.1503
Five sources each one sends four packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	35.0056	36.4891	32.651	34.9213
	100	37.1594	37.9926	35.1258	36.6482
	200	33.9158	34.0052	30.3054	32.1159
	300	30.3052	31.8251	27.5681	29.0062
Five sources each one sends six packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	42.6381	43.4369	40.1592	41.5921
	100	49.5266	51.5126	46.3105	48.5261
	200	41.4925	42.5812	38.1256	39.2251
	300	37.8625	39.0062	34.3526	36.5261
Ten sources each one sends one packet/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	24.5923	25.7435	22.4512	23.6215
	100	26.6041	26.9005	24.5162	25.9871
	200	25.2581	26.0059	22.8403	24.5162
	300	23.6591	24.2489	21.5941	22.6931
Ten sources each one sends two packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	34.0059	35.1189	31.625	33.4259
	100	39.3901	41.5582	37.4982	38.4921
	200	36.2259	38.0047	33.4925	35.4469
	300	34.9562	35.0048	31.5271	33.456

Ten sources each one sends four packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	54.2591	55.3698	50.5741	52.3655
	100	67.2515	69.0014	63.0059	65.2689
	200	59.2681	61.5294	55.2256	57.261
	300	54.2192	54.9821	52.2561	53.3156
Ten sources each one sends six packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	73.6259	75.0045	70.4256	72.1519
	100	74.6982	75.4135	72.1956	73.6984
	200	74.2351	75.2142	71.9852	73.4025
	300	65.2081	66.6288	62.524	64.5291
Fifteen sources each one sends one packet/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	31.9158	32.0058	29.1315	31.1584
	100	32.6175	33.4911	30.2193	31.5802
	200	30.3052	31.9806	28.014	29.3614
	300	28.5611	28.9628	26.2734	27.881
Fifteen sources each one sends two packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	46.2294	47.0058	42.3632	44.9152
	100	49.1058	50.4822	47.5274	48.1266
	200	43.805	44.1142	41.4109	42.2825
	300	42.0052	44.4809	40.5281	41.8536
Fifteen sources each one sends four packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	73.6822	75.118	70.2595	72.3569
	100	74.9184	76.0051	71.5591	73.3056
	200	73.8824	76.1542	70.3521	72.1718
	300	71.9522	72.5496	69.1482	71.0082
Fifteen sources each one sends six packets/second	Pause Time	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX
	0	81.4592	82.5561	78.4159	79.2045
	100	82.5694	83.5529	78.1215	80.806
	200	78.6173	79.8259	76.6482	77.9205
	300	74.5917	75.2845	72.2458	73.3589

APPENDIX C: 95% CONFIDENCE INTERVAL OF PACKET

DELIVERY RATION WHEN $\alpha=0.50$

Table C-1: Confidence Interval of Packet Delivery Ratio with Pause Time=0

Five sources each one sends one packet/second	Confidence Interval	VCAR_Q_A CC	VCAR_Q_M AX	VCAR_R_A CC	VCAR_R_M AX	AODV
	Min	88.157	85.427	96.040	91.616	84.131
	Max	89.237	86.800	96.658	93.191	84.990
Five sources each one sends two packets/second	Confidence Interval	VCAR_Q_A CC	VCAR_Q_M AX	VCAR_R_A CC	VCAR_R_M AX	AODV
	Min	84.631	82.010	92.198	87.952	79.924
	Max	85.667	83.328	92.792	89.464	80.741
Five sources each one sends four packets/second	Confidence Interval	VCAR_Q_A CC	VCAR_Q_M AX	VCAR_R_A CC	VCAR_R_M AX	AODV
	Min	80.399	77.909	87.589	83.554	75.808
	Max	81.384	79.162	88.152	84.991	76.583
Five sources each one sends six packets/second	Confidence Interval	VCAR_Q_A CC	VCAR_Q_M AX	VCAR_R_A CC	VCAR_R_M AX	AODV
	Min	75.977	73.624	82.771	78.958	71.903
	Max	76.907	74.807	83.304	80.316	72.638

Ten sources each one sends one packet/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	87.284	84.581	95.089	90.709	82.724
	Max	88.353	85.941	95.701	92.269	83.570
Ten sources each one sends two packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	86.079	79.468	93.776	89.457	76.425
	Max	87.133	80.810	94.380	90.995	77.256
Ten sources each one sends four packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	76.378	68.541	85.713	79.323	66.656
	Max	77.362	69.827	86.267	80.785	67.432
Ten sources each one sends six packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	68.929	60.366	75.660	69.296	57.444
	Max	69.774	61.499	76.127	70.574	58.112
Fifteen sources each one sends	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV

one packets/second	Min	85.736	79.152	93.403	89.100	76.166
	Max					
Fifteen sources each one sends two packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	75.350	65.277	85.441	75.546	63.725
	Max	76.288	66.502	85.968	76.938	64.467
	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
Fifteen sources each one sends four packets/second	Min	68.348	59.887	74.948	68.678	57.101
	Max	69.186	61.011	75.411	69.944	57.766
Fifteen sources each one sends six packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	61.025	54.942	65.744	62.434	51.628
	Max	61.77	55.973	66.149	63.585	52.229

Table C-2: Confidence Interval of Packet Delivery Ratio with Pause Time=300

Five sources each one sends one packet/seco nd	Confiden ce Interval	VCAR_Q_A CC	VCAR_Q_M AX	VCAR_R_A CC	VCAR_R_M AX	AODV
	Min	98.001	95.331	99.946	99.859	95.250
	Max	98.769	96.954	99.966	99.759	95.875
Five sources each one sends two packets/seco	Confiden ce Interval	VCAR_Q_A CC	VCAR_Q_M AX	VCAR_R_A CC	VCAR_R_M AX	AODV
	Min	94.081	91.518	97.387	95.322	88.160

nd	Max	94.819	93.076	98.393	96.313	88.738
Five sources each one sends four packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_M AX	VCAR_R_A CC	VCAR_R_M AX	AODV
	Min	89.377	86.942	92.518	90.556	85.244
	Max	90.078	88.422	93.474	91.497	85.777
Five sources each one sends six packet/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_M AX	VCAR_R_A CC	VCAR_R_M AX	AODV
	Min	84.461	82.160	87.429	85.576	79.059
	Max	85.124	83.559	88.333	86.465	79.544
Ten sources each one sends one packet/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_M AX	VCAR_R_A CC	VCAR_R_M AX	AODV
	Min	94.532	91.927	98.794	95.894	90.986
	Max	95.892	93.156	99.205	96.645	91.854
Ten sources each one sends two packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_M AX	VCAR_R_A CC	VCAR_R_M AX	AODV
	Min	95.222	92.627	97.702	96.478	88.444
	Max	95.968	94.204	97.926	97.480	89.014

Ten sources each one sends four packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	91.345	87.256	96.992	93.787	78.926
	Max	92.052	88.741	98.120	94.202	79.461
Ten sources each one sends six packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	79.403	75.451	86.280	81.783	68.255
	Max	80.018	76.735	87.255	82.145	68.718
Fifteen sources each one sends one packet/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	96.077	92.258	99.802	98.584	83.793
	Max	96.821	93.829	100.026	99.582	84.361
Fifteen sources each one sends two packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	86.995	83.101	95.135	89.320	75.333
	Max	87.668	84.515	96.209	89.716	75.844

Fifteen sources each one sends four packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_M_AX	VCAR_R_ACC	VCAR_R_M_AX	AODV
	Min	78.771	74.852	85.638	81.134	67.733
	Max	79.381	76.126	86.605	81.493	68.192
Fifteen sources each one sends six packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_M_AX	VCAR_R_ACC	VCAR_R_M_AX	AODV
	Min	71.325	67.422	77.089	73.698	61.453
	Max	71.877	68.570	77.959	74.024	61.869

APPENDIX D: 95% CONFIDENCE INTERVAL OF END-TO-END DELAY WHEN $\alpha=0.50$

Table D-1: Confidence Interval of End-to-End Delay with Pause Time=0

Five sources each one sends one packet/second	Confidence Interval	VCAR_Q_A CC	VCAR_Q_M AX	VCAR_R_A CC	VCAR_R_M AX	AODV
	Min	0.0672	0.0733	0.0555	0.0601	0.0817
	Max	0.0687	0.0744	0.0570	0.0612	0.0843
Five sources each one sends two packets/second	Confidence Interval	VCAR_Q_A CC	VCAR_Q_M AX	VCAR_R_A CC	VCAR_R_M AX	AODV
	Min	0.0847	0.0923	0.0699	0.0757	0.0981
	Max	0.0866	0.0938	0.0718	0.0771	0.1012

Five sources each one sends four packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	0.1090	0.1187	0.0899	0.0974	0.1252
	Max	0.1114	0.1206	0.0924	0.0991	0.1292
Five sources each one sends six packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	0.1402	0.1531	0.1162	0.1253	0.1603
	Max	0.1443	0.155	0.1192	0.1284	0.1652
Ten sources each one sends one packet/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	0.1401	0.1527	0.1156	0.1252	0.1668
	Max	0.1432	0.1551	0.1188	0.1275	0.1721
Ten sources each one sends two packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	0.2919	0.3181	0.2408	0.2608	0.3270
	Max	0.2983	0.3231	0.2474	0.2656	0.3374
Ten sources each one sends four	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV

packets/second	Min	0.3699	0.4084	0.3174	0.3348	0.4345
	Max	0.3781	0.4148	0.3262	0.3410	0.4484
Ten sources each one sends six packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	0.4906	0.5548	0.4281	0.4449	0.6575
nd	Max	0.5015	0.5635	0.4399	0.4530	0.6785
	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
Fifteen sources each one sends one packet/second	Min	0.3317	0.3614	0.2736	0.2964	0.4160
	Max	0.3390	0.3671	0.2812	0.3018	0.4293
Fifteen sources each one sends two packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	0.4252	0.4694	0.3648	0.3849	0.5431
nd	Max	0.4346	0.4768	0.3749	0.3919	0.5605
	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
Fifteen sources each one sends four	Min	0.5452	0.6096	0.4715	0.4998	0.7147

packets/second	Max					
		0.5572	0.6192	0.4849	0.5090	0.7375
Fifteen sources each one sends six packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	0.6989	0.7917	0.6081	0.6491	0.9529
	Max	0.7144	0.8042	0.6249	0.6610	0.9833

Table D-2: Confidence Interval of End-to-End Delay with Pause Time=300

Five sources each one sends one packet/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	0.0183	0.0194	0.0170	0.0180	0.0248
	Max	0.0186	0.0198	0.0175	0.0186	0.0252
Five sources each one sends two packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	0.0319	0.0345	0.0281	0.0318	0.0462
	Max	0.0328	0.0354	0.0349	0.0326	0.0470
Five sources each one sends four packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	0.0296	0.0314	0.0276	0.0292	0.0384
	Max	0.0301	0.0321	0.0283	0.0301	0.0390
Five sources each one sends six packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	0.0382	0.041	0.036	0.0382	0.049
	Max	0.0394	0.0419	0.0369	0.0398	0.0508

Ten sources each one sends one packet/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	0.0381	0.0404	0.0355	0.0375	0.0539
	Max	0.0387	0.0413	0.0364	0.0387	0.0548
Ten sources each one sends two packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	0.1586	0.1684	0.1479	0.1562	0.1926
	Max	0.1611	0.1722	0.1515	0.1614	0.1957
Ten sources each one sends four packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	0.2397	0.2572	0.2142	0.2308	0.3007
	Max	0.2434	0.2630	0.2194	0.2385	0.3057
Ten sources each one sends six packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	0.3298	0.3598	0.2793	0.3089	0.4202
	Max	0.3349	0.3679	0.2860	0.3192	0.4271
Fifteen sources each one sends	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV

one packet/second	Min	0.2049	0.2175	0.1910	0.2017	0.2504
	Max	0.2081	0.2223	0.1956	0.2084	0.2545
Fifteen sources each one sends two packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	0.2755	0.2957	0.2462	0.2653	0.3417
	Max	0.2798	0.3023	0.2521	0.2742	0.3474
	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
Fifteen sources each one sends four packets/second	Min	0.3705	0.4020	0.3174	0.3490	0.4669
	Max	0.3763	0.4110	0.3250	0.3607	0.4746
Fifteen sources each one sends six packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	0.4983	0.5465	0.4091	0.4591	0.6466
	Max	0.5061	0.5588	0.4189	0.4744	0.6573

APPENDIX E: 95% CONFIDENCE INTERVAL OF ROUTING

OVERHEAD WHEN $\alpha=0.50$

Table E-1: Confidence Interval of Routing Overhead with Pause Time=0

Five sources each one sends one packet/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_M_AX	VCAR_R_ACC	VCAR_R_M_AX	AODV
	Min	1.8496	1.8582	1.7934	1.8015	1.7387
	Max	1.8639	1.8700	1.8044	1.8291	1.7644
Five sources each one sends two packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_M_AX	VCAR_R_ACC	VCAR_R_M_AX	AODV
	Min	2.0530	2.0626	1.9907	1.9997	1.9265
	Max	2.0689	2.0758	2.0028	2.0303	1.9549
Five sources each one sends four packet/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_M_AX	VCAR_R_ACC	VCAR_R_M_AX	AODV
	Min	2.3037	2.3145	2.2338	2.2439	2.1593
	Max	2.3216	2.3293	2.2474	2.2783	2.1911
Five sources each one sends six packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_M_AX	VCAR_R_ACC	VCAR_R_M_AX	AODV
	Min	2.6081	2.6202	2.5289	2.5403	2.4417
	Max	2.6283	2.6369	2.5443	2.5792	2.4777
Ten sources each one sends one	Confidence Interval	VCAR_Q_ACC	VCAR_Q_M_AX	VCAR_R_ACC	VCAR_R_M_AX	AODV

packet/second	Min	2.4661	2.4776	2.3913	2.4021	2.1456
	Max	2.4852	2.4934	2.4059	2.4389	2.1758
Ten sources each one sends two packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_M_AX	VCAR_R_ACC	VCAR_R_M_AX	AODV
	Min	2.0050	2.0143	1.9441	1.9529	1.7047
	Max	2.0205	2.0272	1.9560	1.9828	1.7298
Ten sources each one sends four packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_M_AX	VCAR_R_ACC	VCAR_R_M_AX	AODV
	Min	1.8309	1.8395	1.7745	1.7826	1.7463
	Max	1.8452	1.8514	1.7855	1.8103	1.7725
Ten sources each one sends six packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_M_AX	VCAR_R_ACC	VCAR_R_M_AX	AODV
	Min	1.1051	1.1186	1.0190	1.0326	0.9693
	Max	1.1269	1.1366	1.0356	1.0745	1.0088
Fifteen sources each one sends one packet/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_M_AX	VCAR_R_ACC	VCAR_R_M_AX	AODV
	Min	2.1105	2.1203	2.0464	2.0557	2.0055
	Max	2.1269	2.1339	2.0589	2.0872	2.0351

Fifteen sources each one sends two packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	2.0805	2.0903	2.0164	2.0257	1.9755
	Max	2.0969	2.1039	2.0289	2.0572	2.0051
Fifteen sources each one sends four packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	1.1074	1.1219	1.0205	1.0342	0.9751
	Max	1.1291	1.1400	1.0372	1.0763	1.0149
Fifteen sources each one sends six packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	1.4765	1.4979	1.3589	1.3790	1.3142
	Max	1.5055	1.5221	1.3811	1.4350	1.3678

Table E-2: Confidence Interval of Routing overhead with Pause Time=300

Five sources each one sends one packet/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_M_AX	VCAR_R_ACC	VCAR_R_M_AX	AODV
	Min	1.3476	1.3822	1.3033	1.3102	1.2522
	Max	1.3987	1.4300	1.3337	1.3494	1.2766
Five sources each one sends two packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_M_AX	VCAR_R_ACC	VCAR_R_M_AX	AODV
	Min	1.4958	1.5342	1.4467	1.4543	1.3886
	Max	1.5526	1.5872	1.4804	1.4978	1.4156
Five sources each one sends four packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_M_AX	VCAR_R_ACC	VCAR_R_M_AX	AODV
	Min	1.6785	1.7216	1.6233	1.6319	1.5745
	Max	1.7422	1.7810	1.6612	1.6807	1.6051
Five sources each one sends six packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_M_AX	VCAR_R_ACC	VCAR_R_M_AX	AODV
	Min	1.9002	1.9490	1.8378	1.8475	1.7869
	Max	1.9723	2.0163	1.8807	1.9027	1.8216
Ten sources each one sends one	Confidence Interval	VCAR_Q_ACC	VCAR_Q_M_AX	VCAR_R_ACC	VCAR_R_M_AX	AODV

packet/second	Min	1.7968	1.8429	1.7378	1.7469	1.6808
	Max	1.8650	1.9066	1.7783	1.7992	1.7135
Ten sources each one sends two packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_M_AX	VCAR_R_ACC	VCAR_R_M_AX	AODV
	Min	1.4608	1.4983	1.4128	1.4203	1.4007
	Max	1.5162	1.5501	1.4458	1.4628	1.4279
Ten sources each one sends four packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_M_AX	VCAR_R_ACC	VCAR_R_M_AX	AODV
	Min	1.3796	1.4143	1.3258	1.3420	1.3028
	Max	1.4309	1.4623	1.3564	1.3814	1.3277
Ten sources each one sends six packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_M_AX	VCAR_R_ACC	VCAR_R_M_AX	AODV
	Min	0.8220	0.8748	0.7083	0.7655	0.6997
	Max	0.8996	0.9473	0.7554	0.8250	0.7359
Fifteen sources each one sends one packet/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_M_AX	VCAR_R_ACC	VCAR_R_M_AX	AODV
	Min	1.5377	1.5772	1.4766	1.4950	1.4744
	Max	1.5960	1.6317	1.5114	1.5397	1.5031

Fifteen sources each one sends two packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	1.5677	1.6072	1.5066	1.5250	1.5044
	Max	1.6260	1.6617	1.5414	1.5697	1.5331
Fifteen sources each one sends four packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	0.8236	0.8774	0.7094	0.7667	0.7018
	Max	0.9014	0.9502	0.7565	0.8263	0.7381
Fifteen sources each one sends six packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	1.0981	1.1714	0.9446	1.0222	0.9114
	Max	1.2019	1.2686	1.0074	1.1018	0.9586

APPENDIX F: 95% CONFIDENCE INTERVAL OF ENERGY

CONSUMPTION PERCENTAGE WHEN $\alpha=0.50$

Table F-1: Confidence Interval of Energy Consumption Percentage with Pause Time=0

Five sources each one sends one packet/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	21.260	22.016	19.120	19.736	27.272
	Max	21.895	22.621	19.835	20.391	27.879
Five sources each one sends two packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	23.811	24.658	21.414	22.105	30.408
	Max	24.522	25.335	22.215	22.838	31.085
Five sources each one sends four packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	28.335	29.343	25.483	26.305	34.027
	Max	29.181	30.149	26.436	27.177	34.784
Five sources each one sends six	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min					
	Max					

packets/second	Min	34.568	34.918	30.324	31.303	41.172
	Max	35.601	35.878	31.458	32.340	42.089
Ten sources each one sends one packet/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_M_AX	VCAR_R_ACC	VCAR_R_M_AX	AODV
	Min	27.256	28.226	24.512	25.303	34.090
	Max	28.070	29.001	25.429	26.142	34.848
Ten sources each one sends two packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_M_AX	VCAR_R_ACC	VCAR_R_M_AX	AODV
	Min	38.388	38.665	36.586	36.671	49.085
	Max	39.536	39.728	37.954	37.887	50.178
Ten sources each one sends four packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_M_AX	VCAR_R_ACC	VCAR_R_M_AX	AODV
	Min	59.099	59.803	56.494	56.968	68.142
	Max	60.866	61.447	58.607	58.857	69.658
Ten sources each one sends six packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_M_AX	VCAR_R_ACC	VCAR_R_M_AX	AODV
	Min	61.584	62.983	58.240	59.357	70.547
	Max	63.424	64.714	60.418	61.324	72.116

Fifteen sources each one sends one packet/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	46.216	49.451	43.905	46.014	59.216
	Max	47.687	50.813	45.659	47.573	60.455
Fifteen sources each one sends two packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	55.927	58.888	53.483	55.006	68.241
	Max	57.599	60.443	55.483	56.793	69.649
Fifteen sources each one sends four packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	63.554	67.598	60.984	62.930	74.221
	Max	65.453	69.373	63.265	64.979	75.815
Fifteen sources each one sends six packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	72.220	73.742	69.537	70.906	82.351
	Max	74.378	75.768	72.138	73.256	84.202

Table F-2: Confidence Interval of Energy Consumption Percentage with Pause Time=300

Five sources each one sends one packet/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	15.135	16.246	13.820	14.315	19.951
	Max	15.508	16.652	14.103	14.607	20.406
Five sources each one sends two packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	16.952	18.195	15.478	16.032	22.644
	Max	17.369	18.650	15.795	16.360	23.161
Five sources each one sends four packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	20.172	21.652	18.419	19.079	25.543
	Max	20.669	22.193	18.796	19.469	26.126
Five sources each one sends six packet/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	24.610	25.766	21.919	22.703	31.341
	Max	25.216	26.410	22.367	23.168	32.056

Ten sources each one sends one packet/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	19.404	20.828	17.718	18.352	27.330
	Max	19.882	21.348	18.080	18.727	27.954
Ten sources each one sends two packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	27.330	28.531	26.445	26.597	39.374
	Max	28.002	29.244	26.985	27.141	40.273
Ten sources each one sends four packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	54.369	56.061	50.244	53.210	60.807
	Max	55.707	57.462	51.272	54.298	62.055
Ten sources each one sends six packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	55.504	56.515	54.324	54.637	68.417
	Max	56.869	57.928	55.435	55.755	69.684
Fifteen sources each one sends	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV

one packet/second	Min	44.325	46.405	42.791	43.956	53.507
	Max	45.489	47.640	43.728	44.898	54.956
Fifteen sources each one sends two packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	54.023	55.699	51.073	52.877	64.099
Fifteen sources each one sends four packets/second	Max	55.352	57.091	52.149	53.958	65.517
	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
Fifteen sources each one sends six packets/second	Min	61.671	62.795	58.360	60.708	73.183
	Max	63.188	64.364	59.594	61.950	74.576
Fifteen sources each one sends six packets/second	Confidence Interval	VCAR_Q_ACC	VCAR_Q_MAX	VCAR_R_ACC	VCAR_R_MAX	AODV
	Min	70.400	70.794	67.033	69.699	82.585
Fifteen sources each one sends six packets/second	Max	72.132	72.564	68.009	71.125	83.961